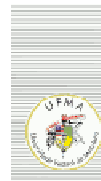


Projeto e Análise de Algoritmos

Aula 1: Algoritmos de Ordenação

Prof. Carlos de Salles
Terças-feiras, 8h20 às 11h10



deINF
departamentodeinformática



Algoritmos de Ordenação

- Insertsort
- Mergesort
- Heapsort
- Quicksort



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Algoritmos de Ordenação

- Dado um vetor **v** com **N** elementos, a ordenação consiste em organizar todos esses **N** elementos em uma ordem (não-crescente, não-decrescente etc)
- Exemplo:
 - $v = \{1, 9, 8, 5, 3, 7, 4\}$
 - Depois de ordenado de forma não-decrescente, v fica:
 - $v = \{1, 3, 4, 5, 7, 8, 9\}$



Insert Sort

- Outros nomes:
 - Bubble sort
 - Selection sort
- Consiste em posicionar o menor ou o maior elemento em seu lugar correto **N** vezes consecutivas

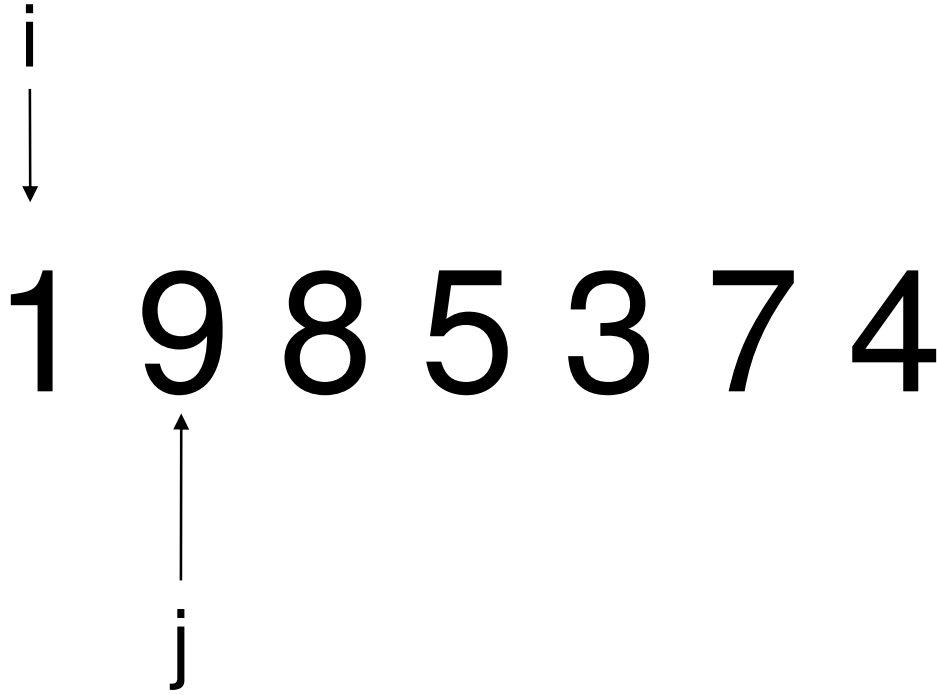


Insert Sort

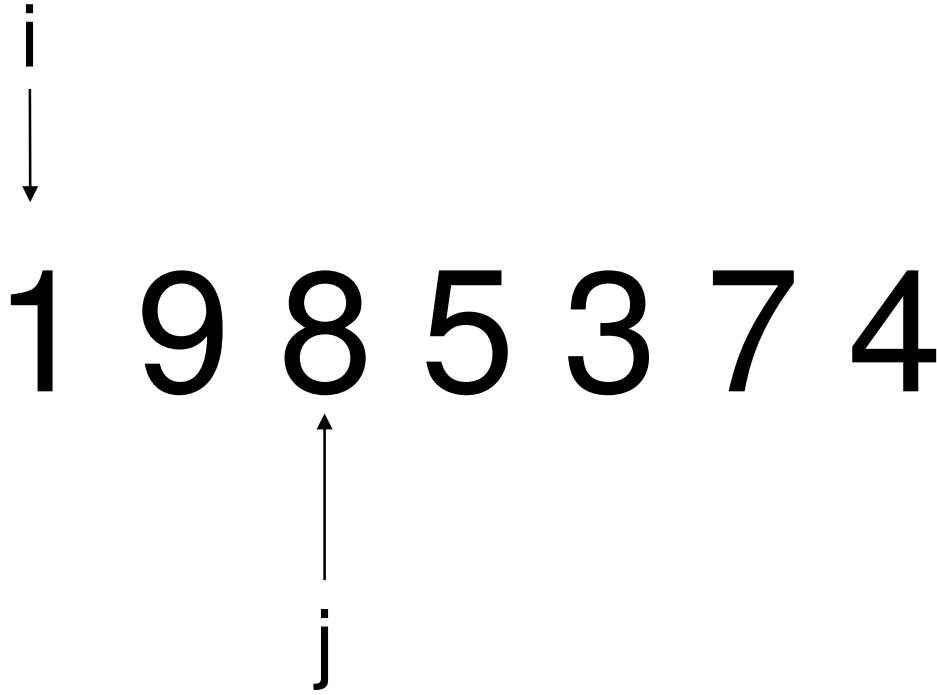
```
for i=1,N do  
    for j=i+1,N do  
        if v[i]>v[j] then  
            v[i],v[j] = v[j],v[i]  
        end  
    end  
end
```



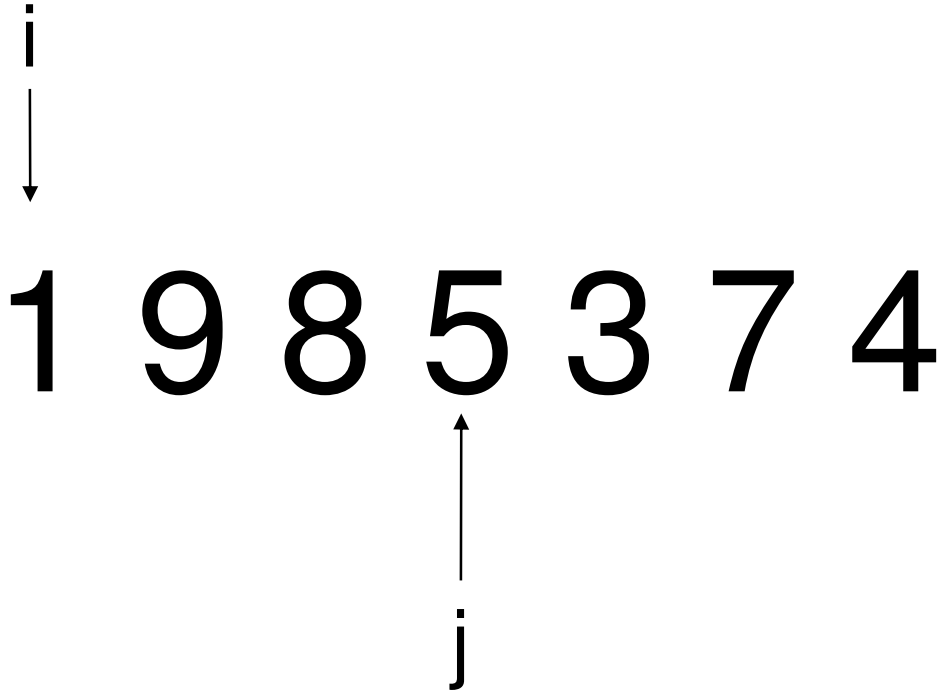
Insert Sort



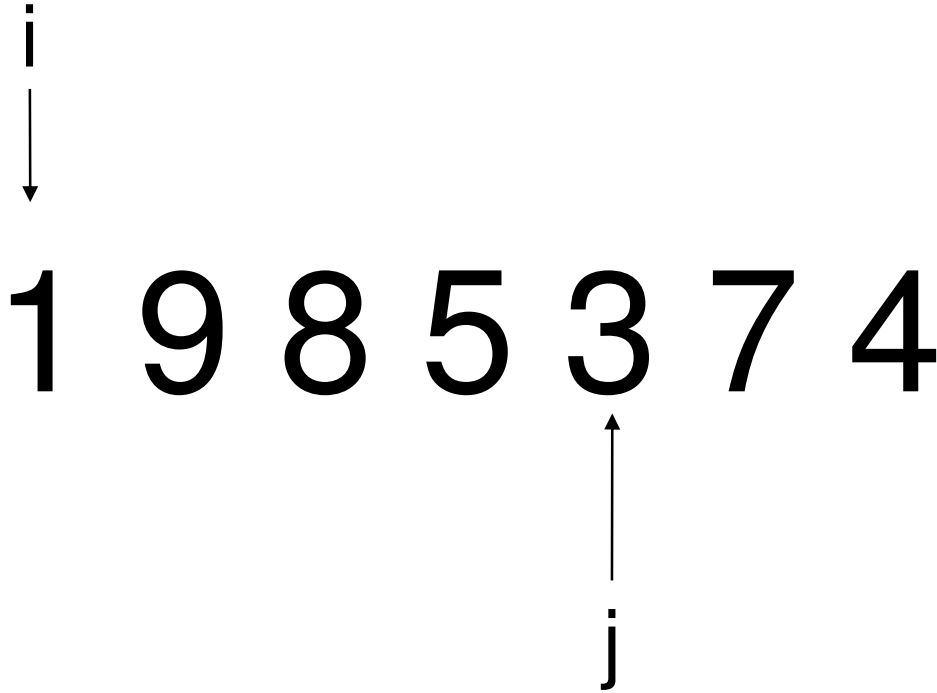
Insert Sort



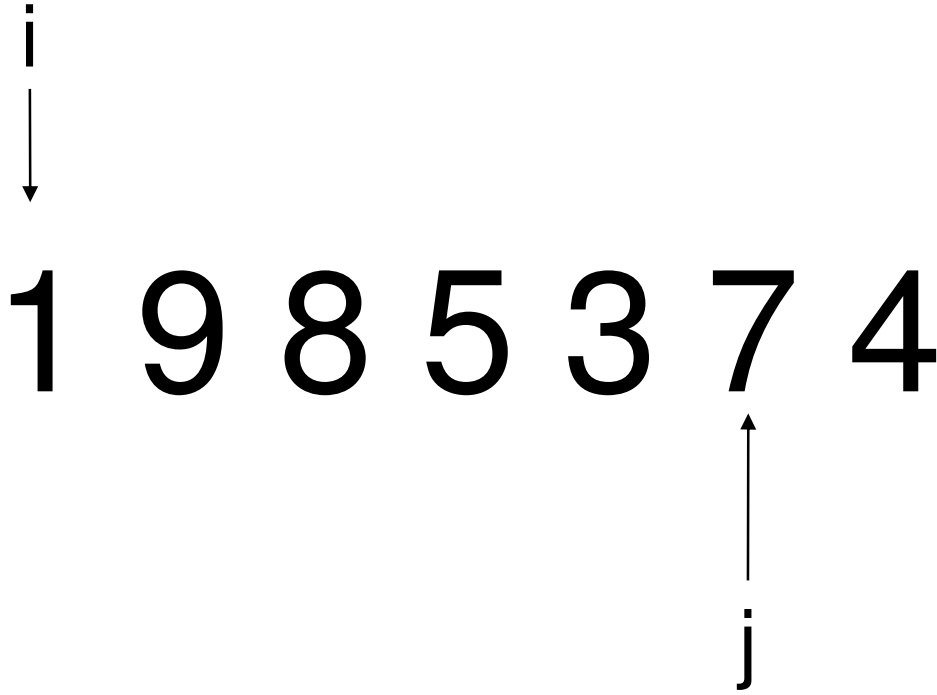
Insert Sort



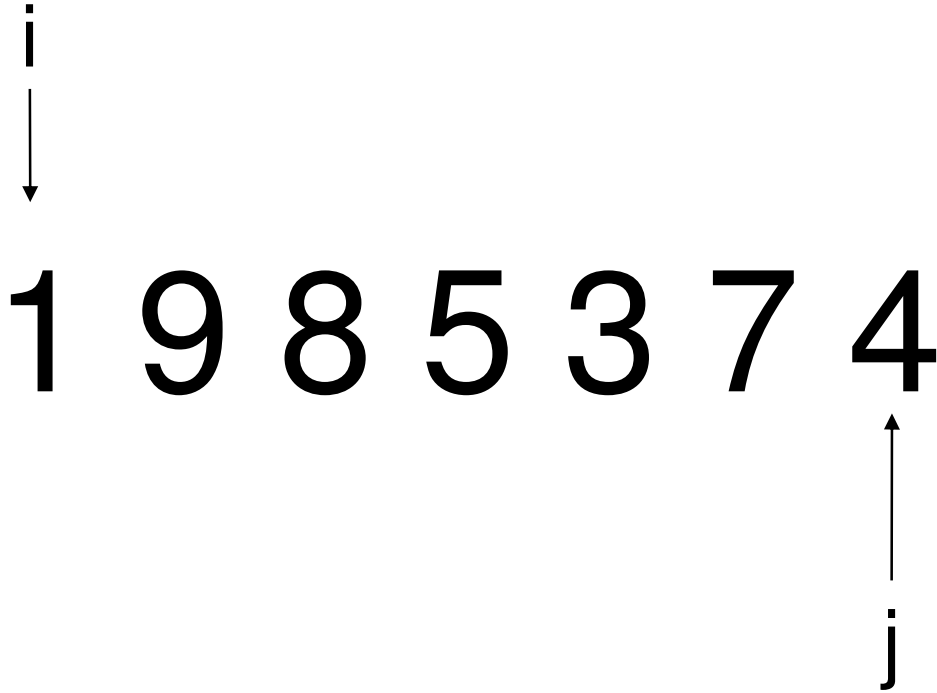
Insert Sort



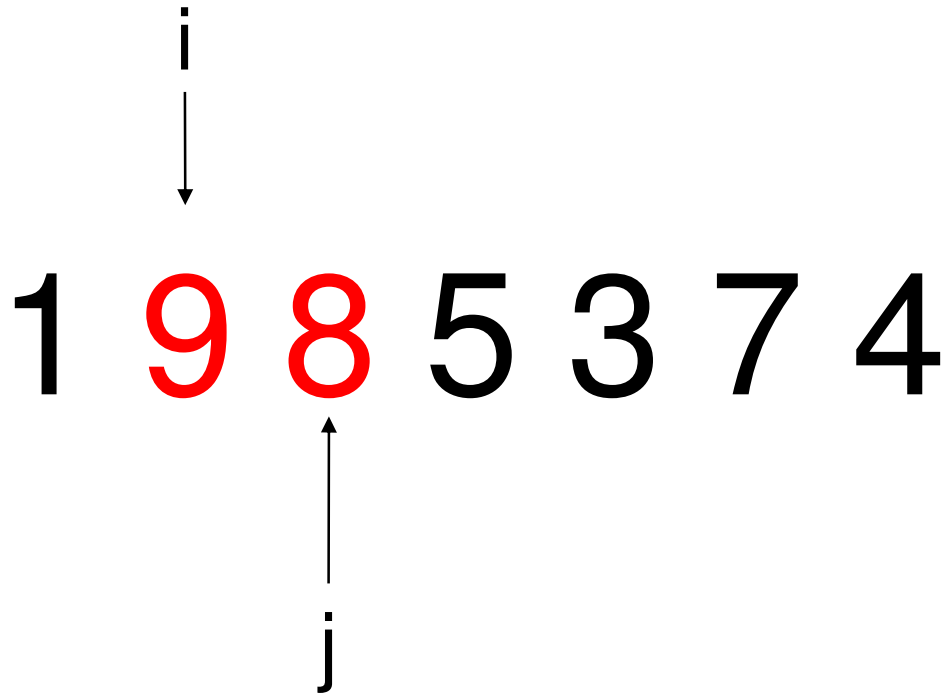
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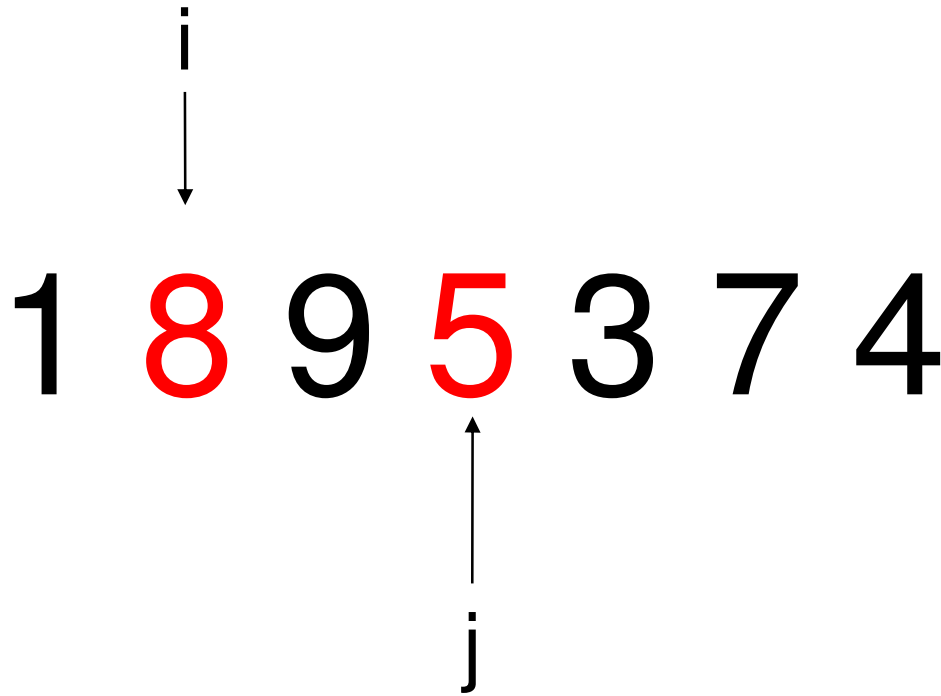
Insert Sort



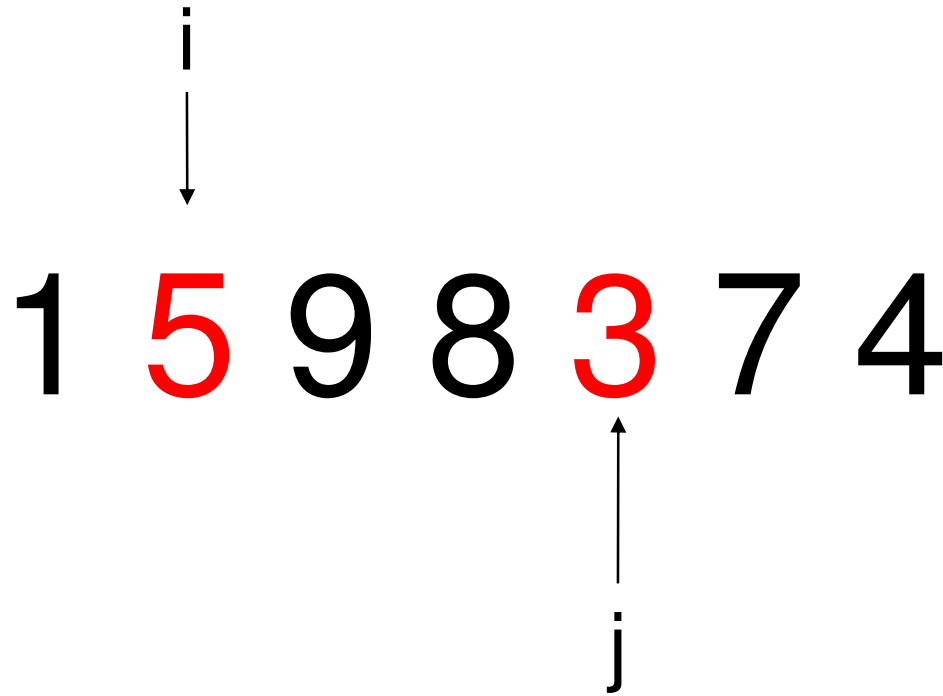
Insert Sort



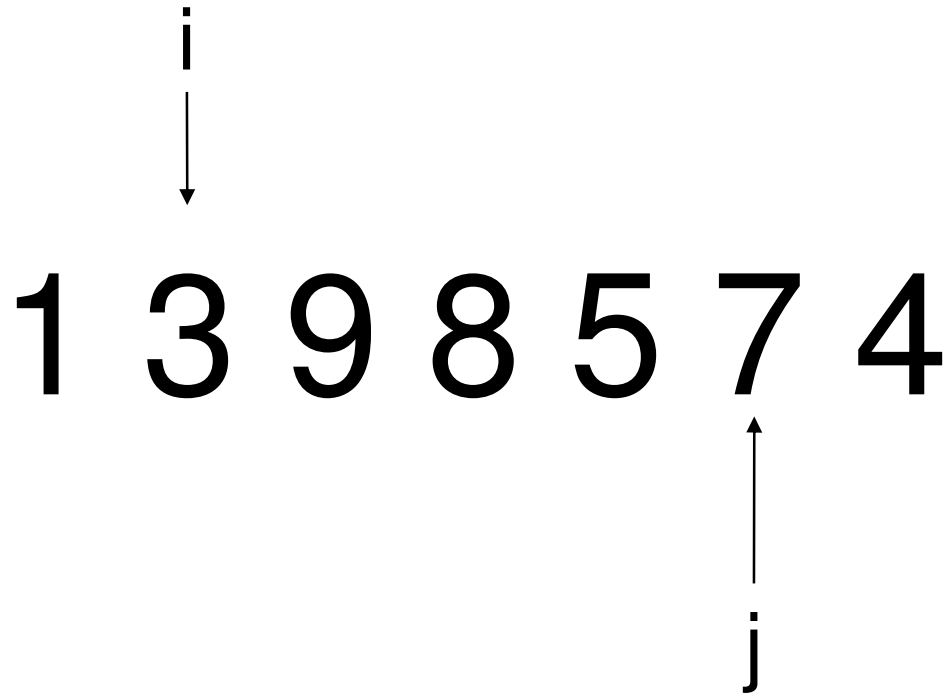
Insert Sort



Insert Sort



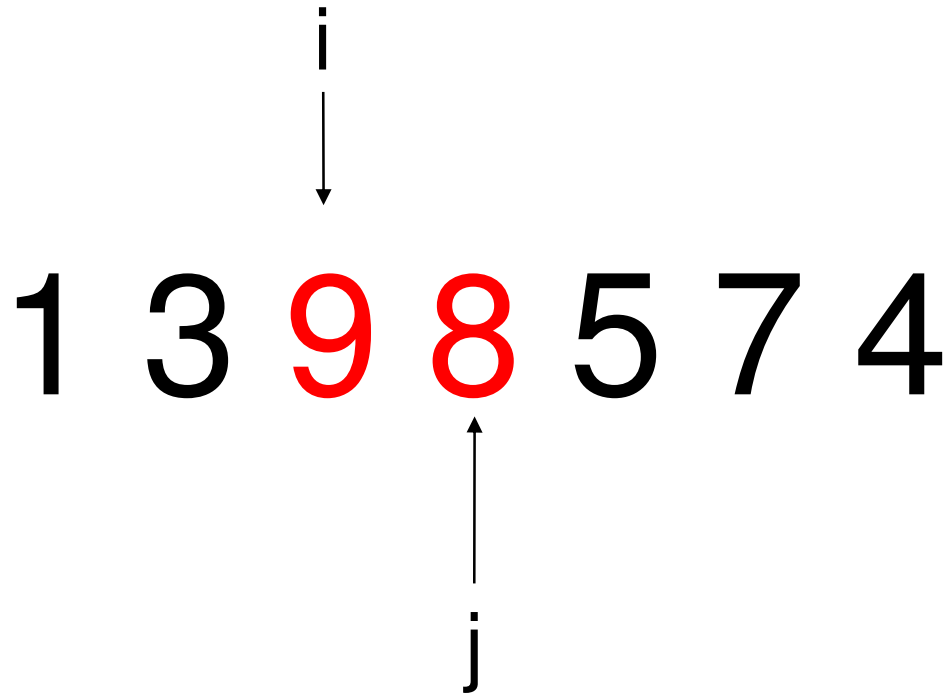
Insert Sort



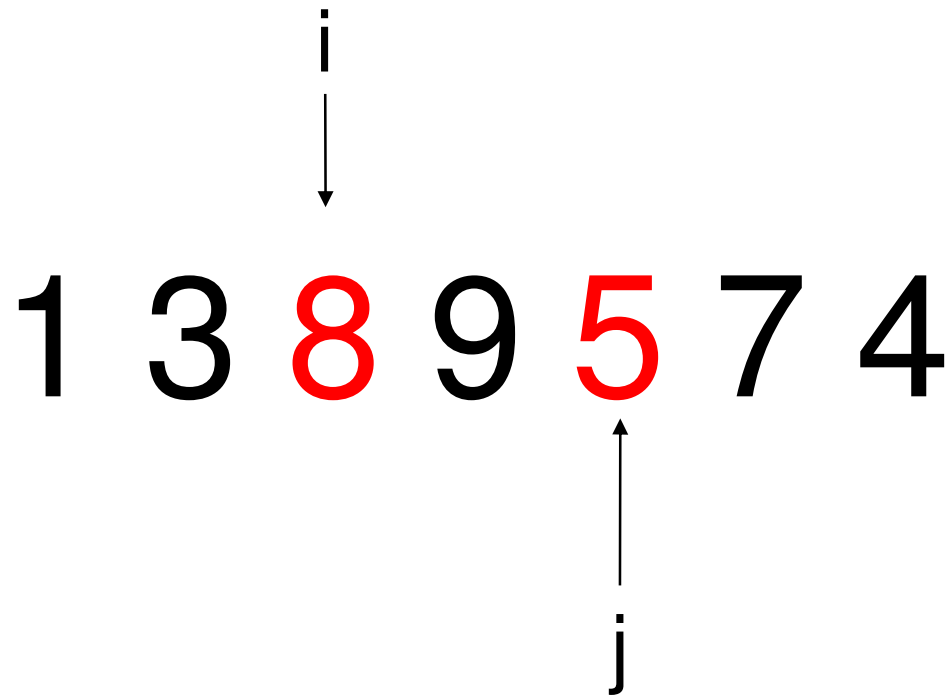
Insert Sort

i
↓
1 3 9 8 5 7 4
↑
j

Insert Sort



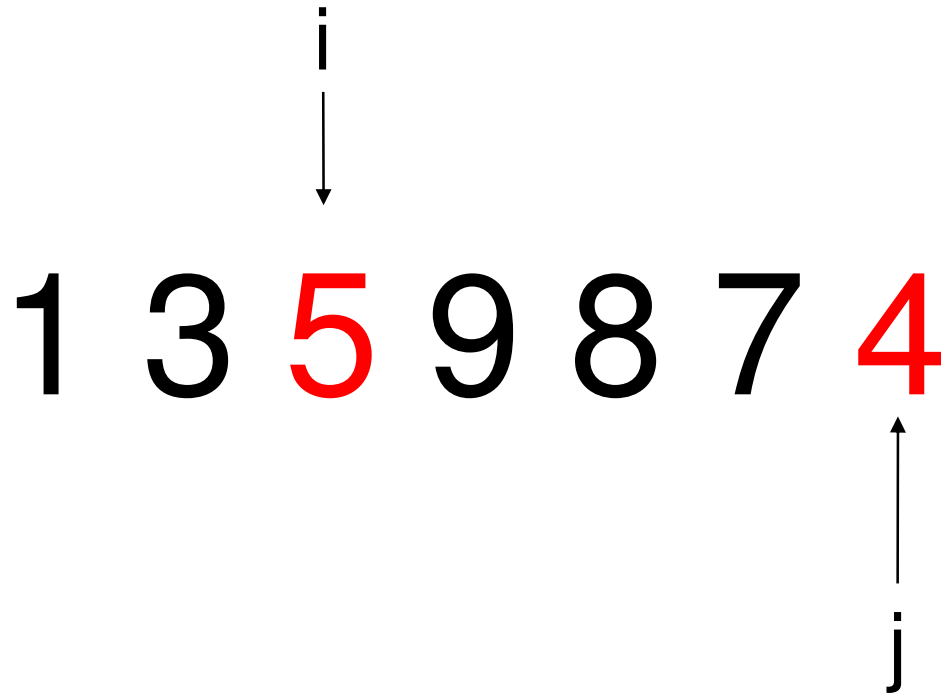
Insert Sort



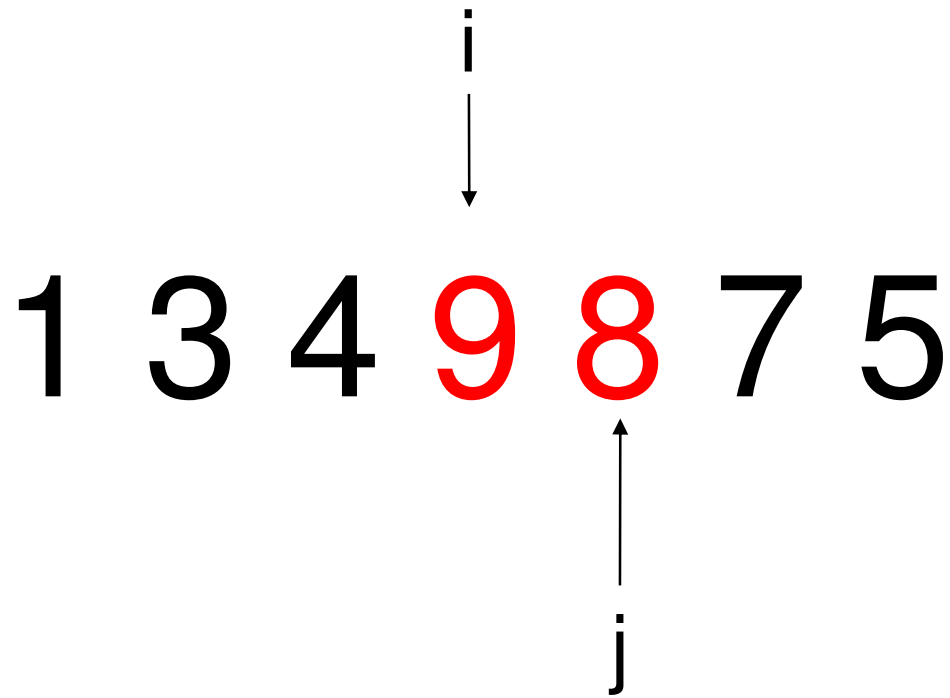
Insert Sort

i
↓
1 3 5 9 8 7 4
↑
j

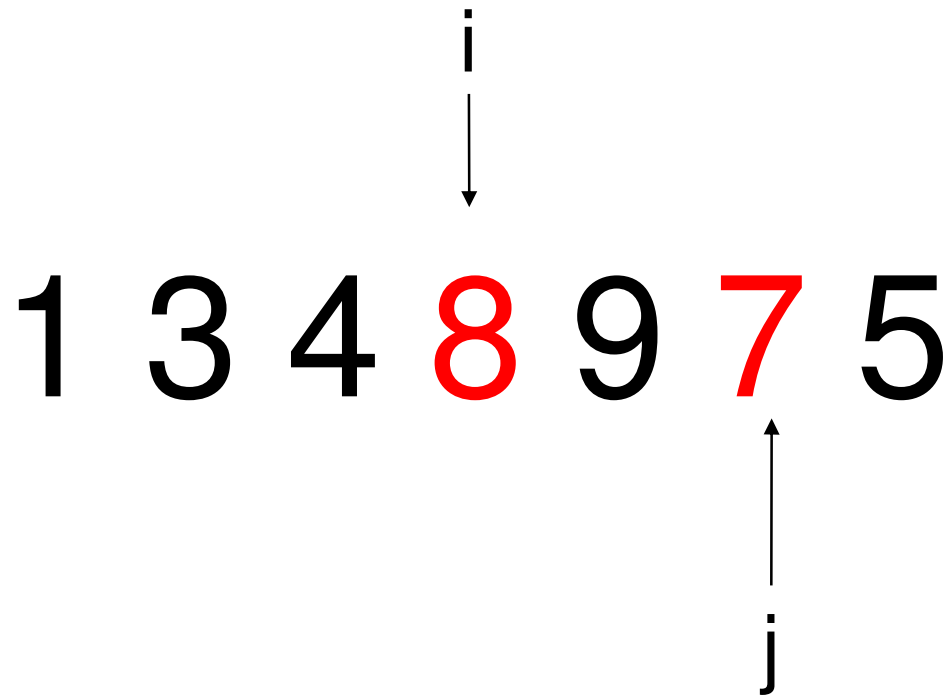
Insert Sort



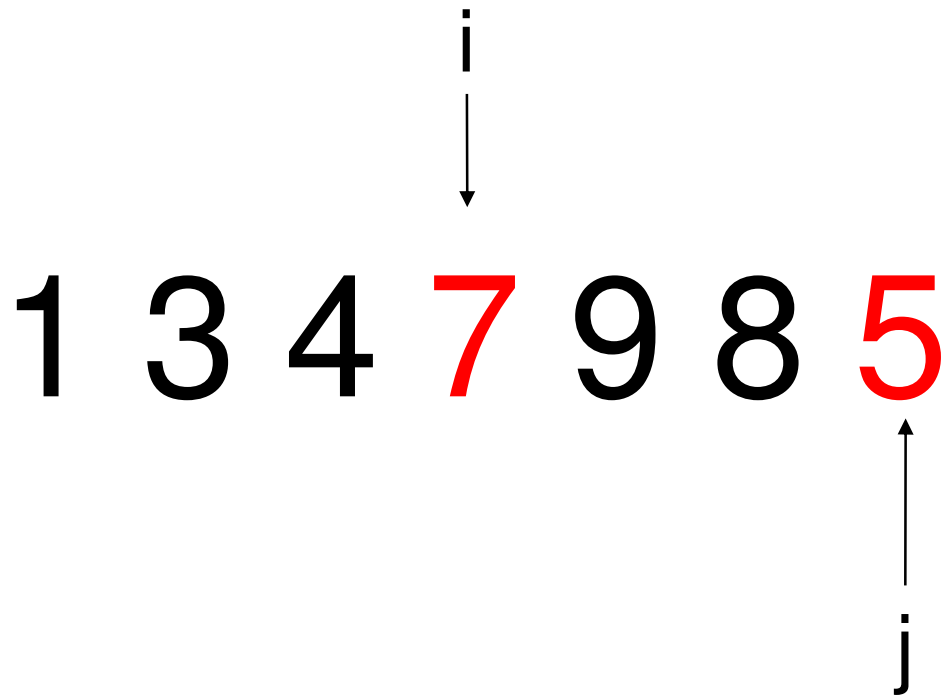
Insert Sort



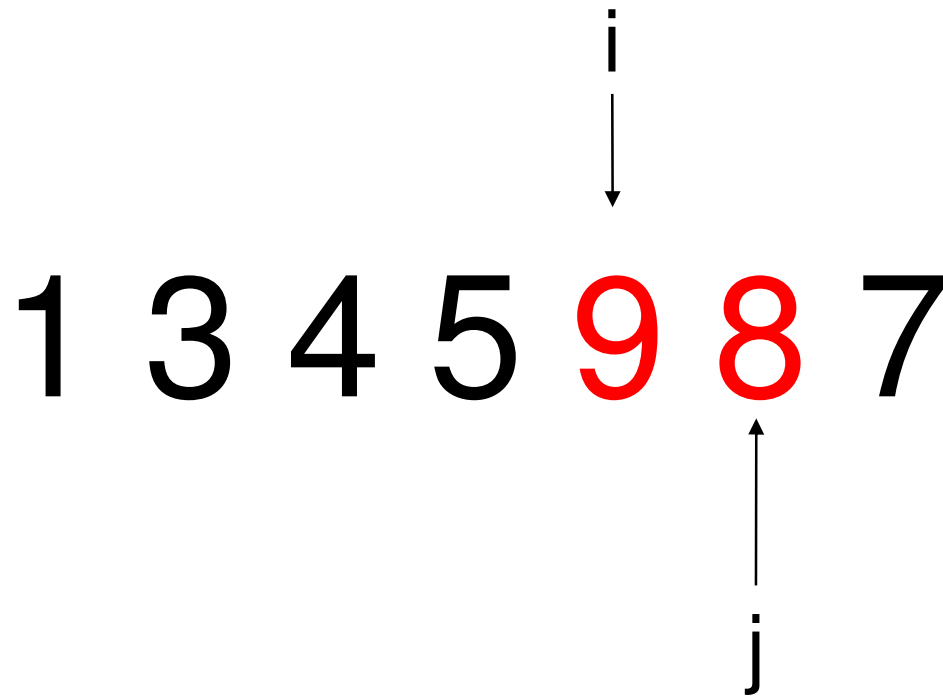
Insert Sort



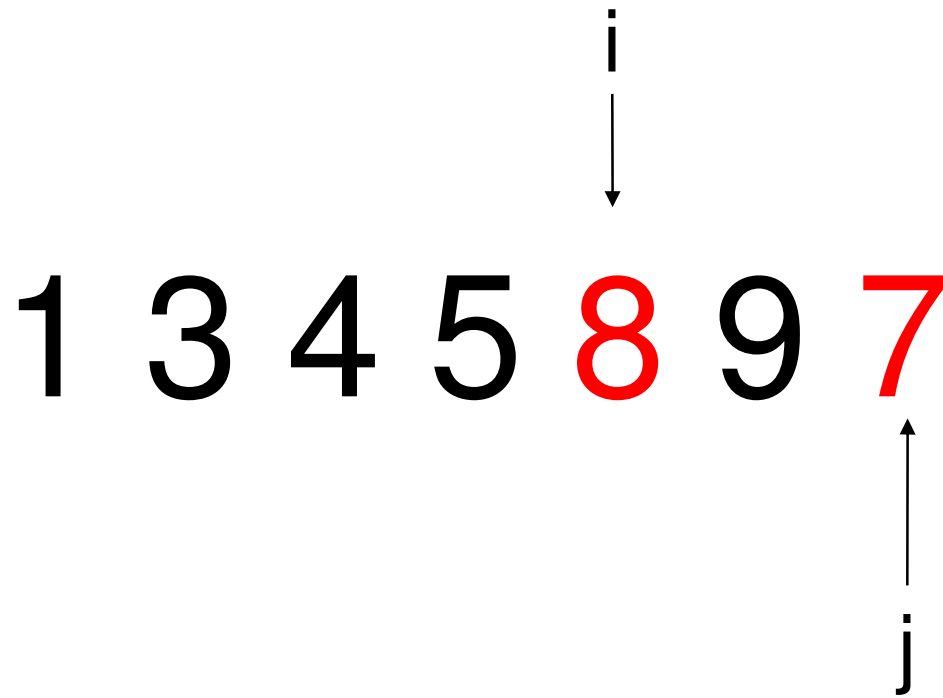
Insert Sort



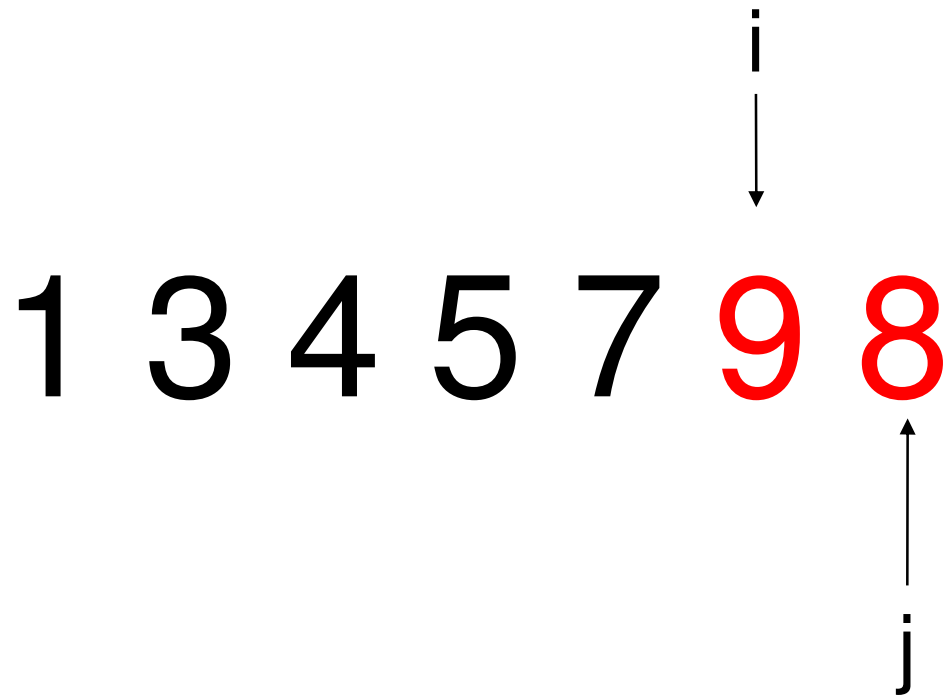
Insert Sort



Insert Sort



Insert Sort



Insert Sort

1 3 4 5 7 8 9

Insert Sort

1 3 4 5 7 8 9

Número de operações:

$(N-1) + (N-2) + (N-3) + \dots + 3 + 2 + 1$

$$\frac{N*(N-1)}{2}$$



Insert Sort

- Ideia:
 - Se o vetor já estiver ordenado em alguma interação, seria interessante detectar isso e parar o algoritmo.
- Como detecto se um vetor está ordenado?

```
ordenado = true
```

```
for i=1,N do
```

```
    if v[i]>v[i+1] then ordenado = false end
```

```
end
```

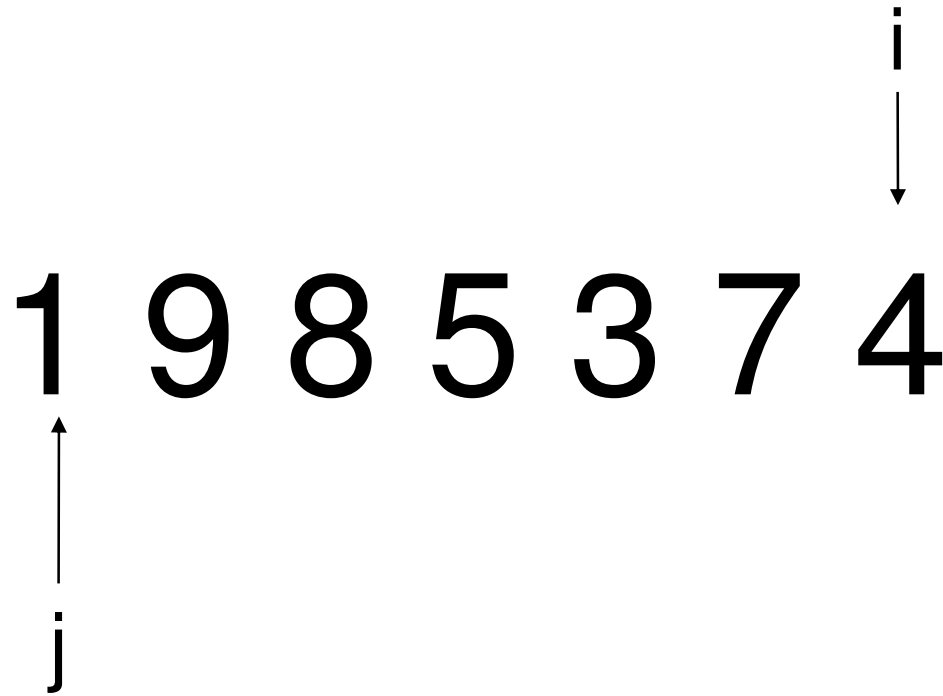


Insert Sort (melhoria)

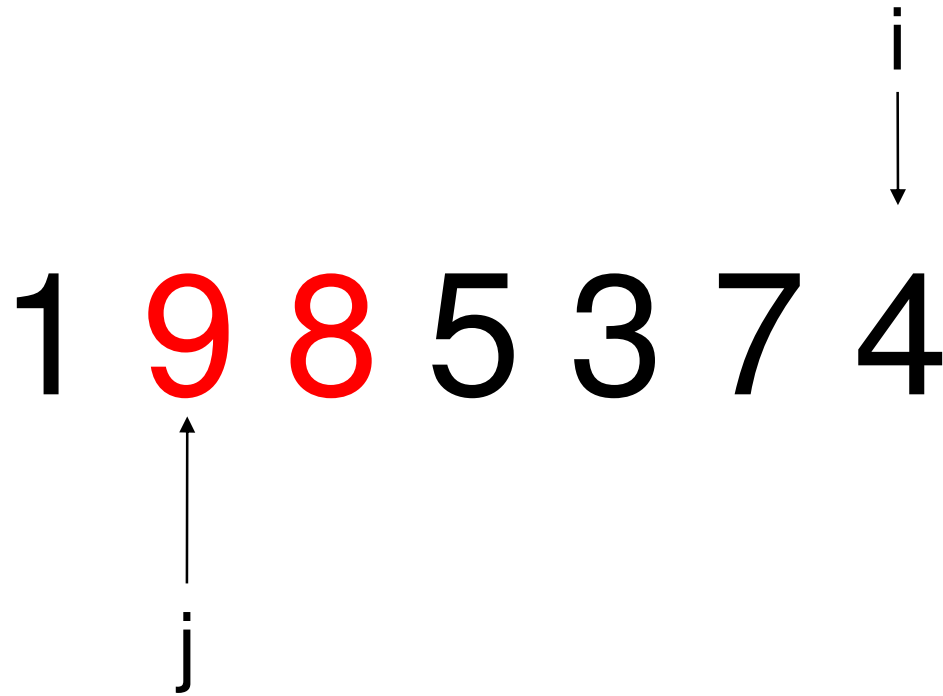
```
for i=N,2,-1 do  
    trocas = false  
    for j=1,i-1 do  
        if v[j]>v[j+1] then  
            v[j],v[j+1] = v[j+1],v[j]  
            trocas = true  
        end  
    end  
    if not trocas then break end  
end
```



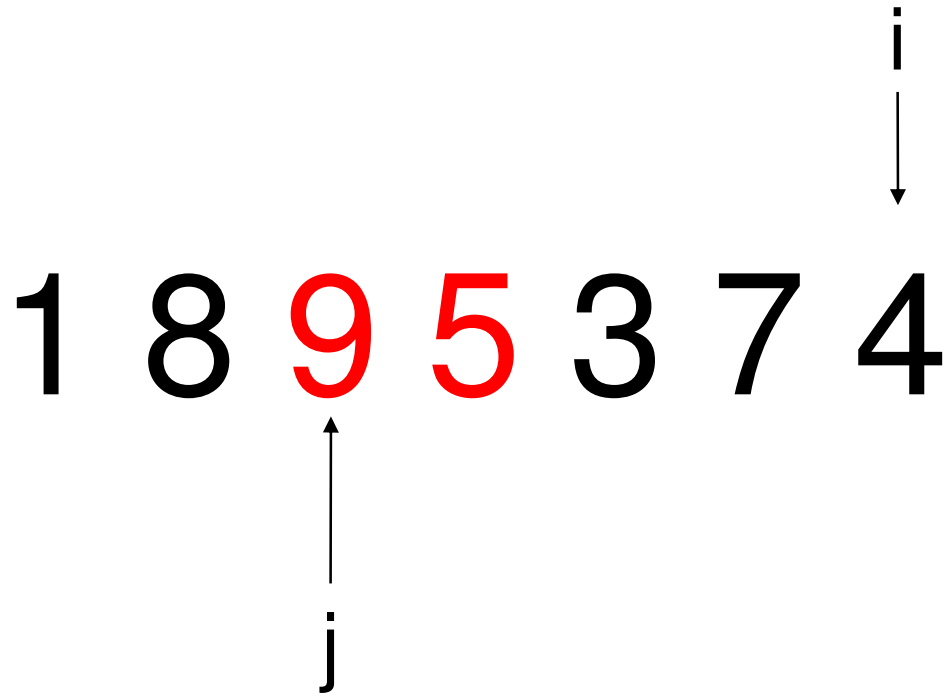
Insert Sort (melhoria)



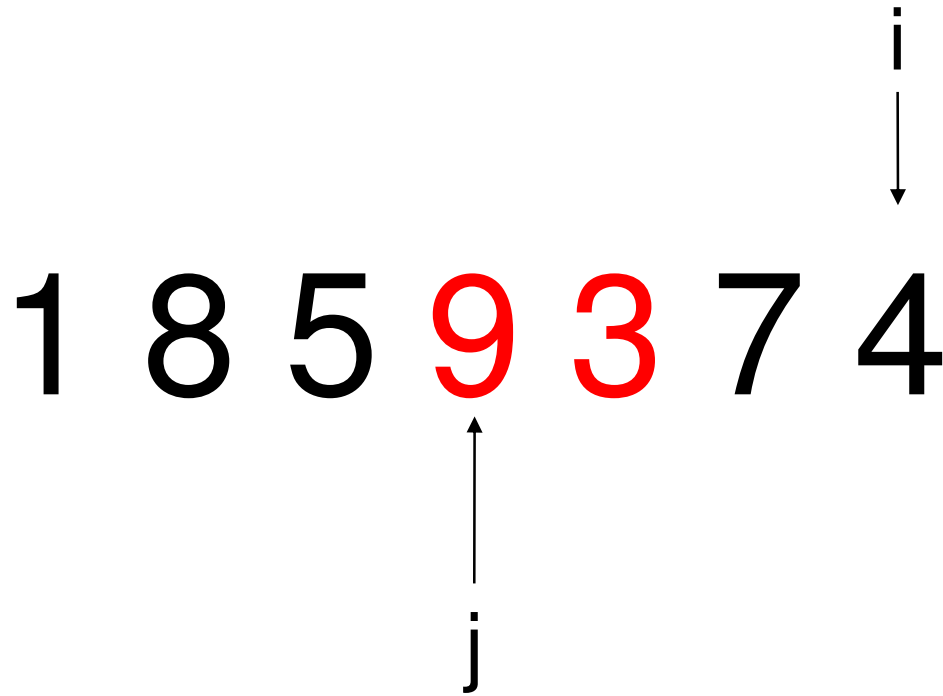
Insert Sort (melhoria)



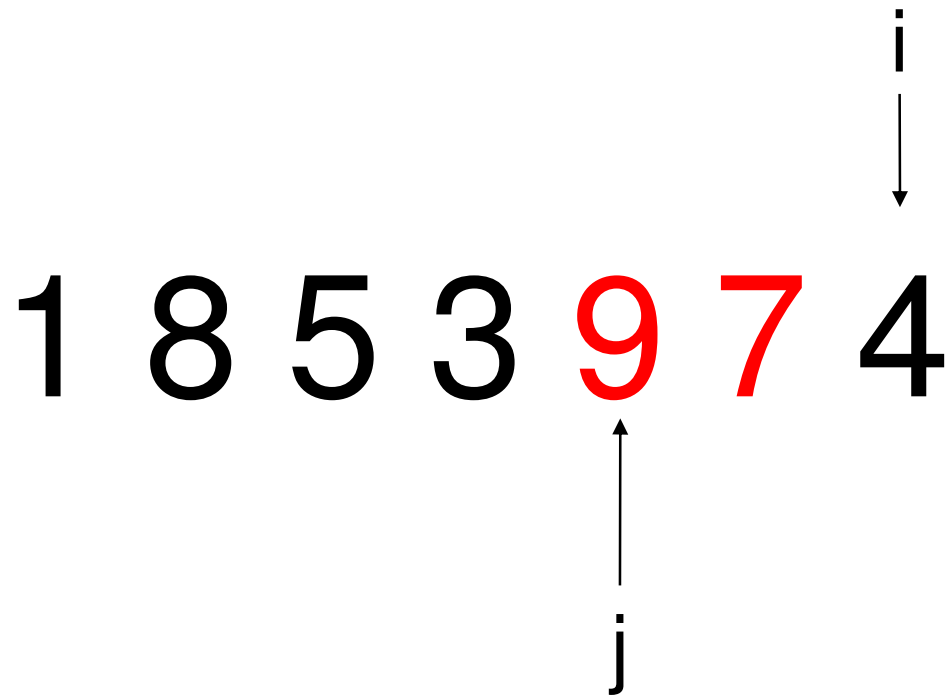
Insert Sort (melhoria)



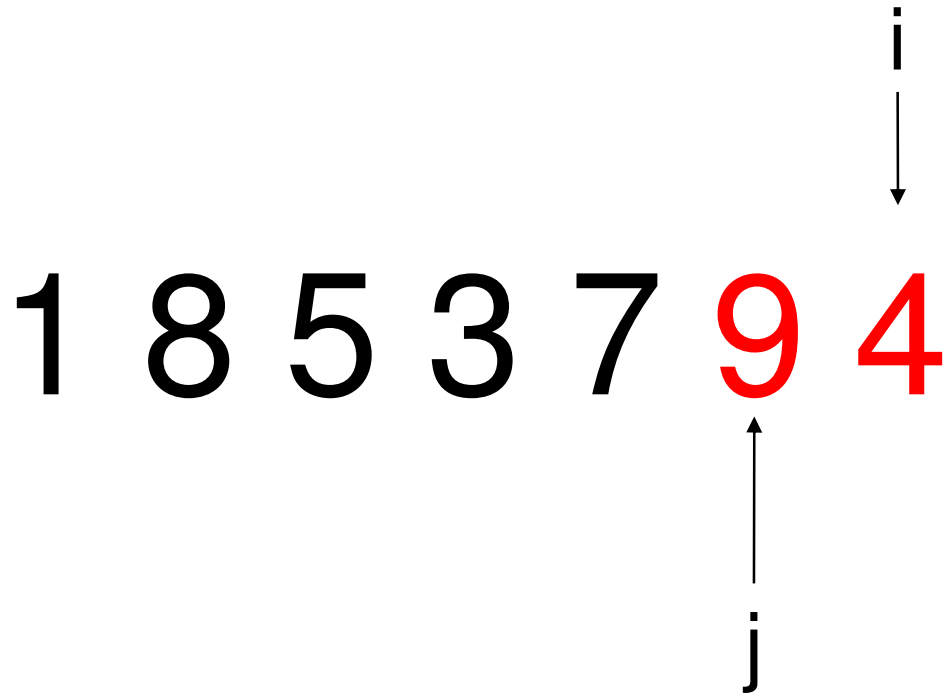
Insert Sort (melhoria)



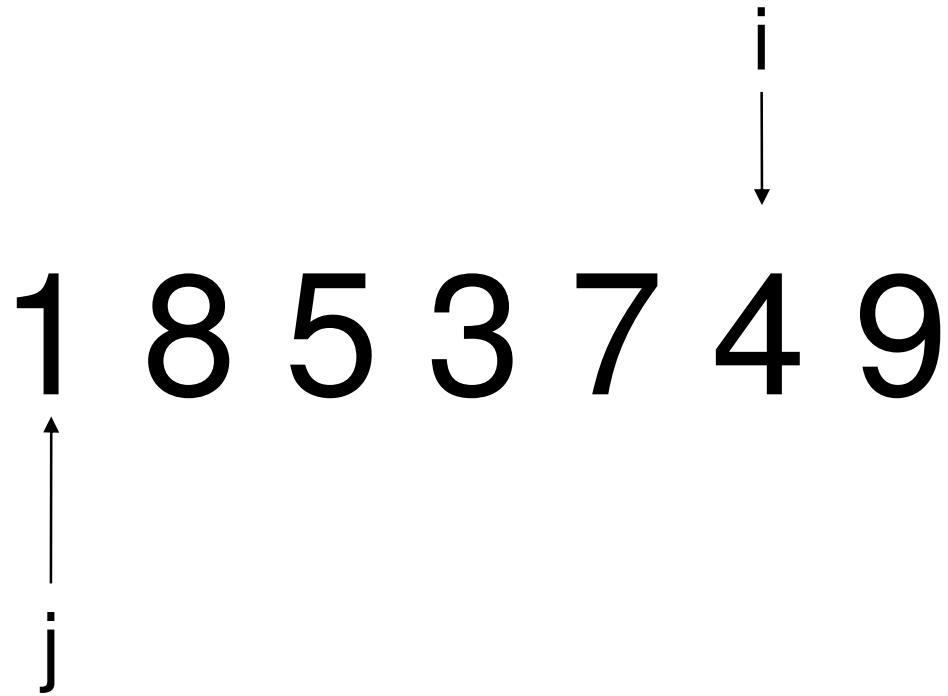
Insert Sort (melhoria)



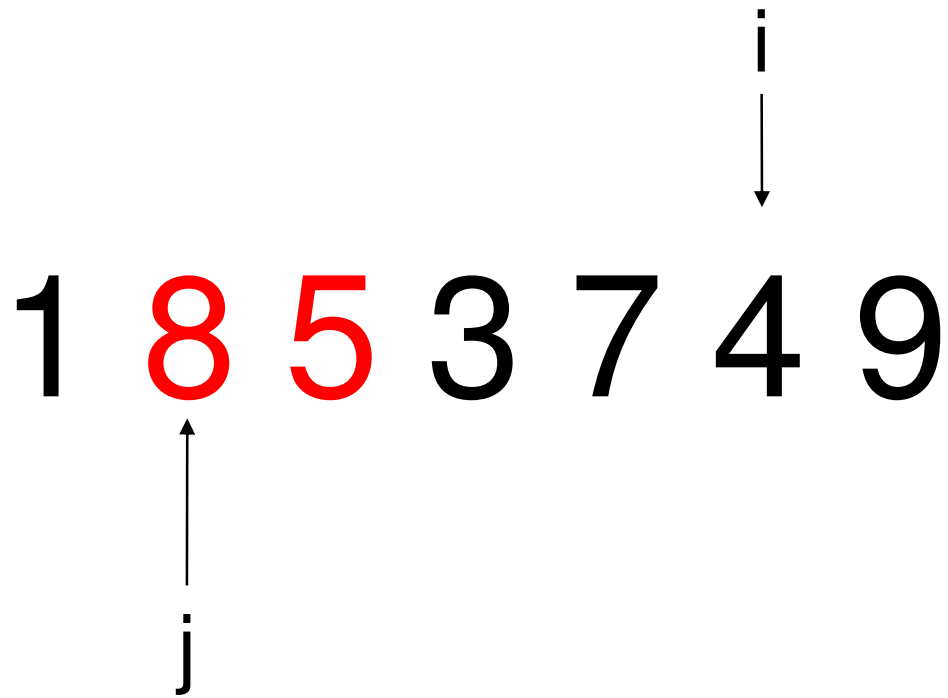
Insert Sort (melhoria)



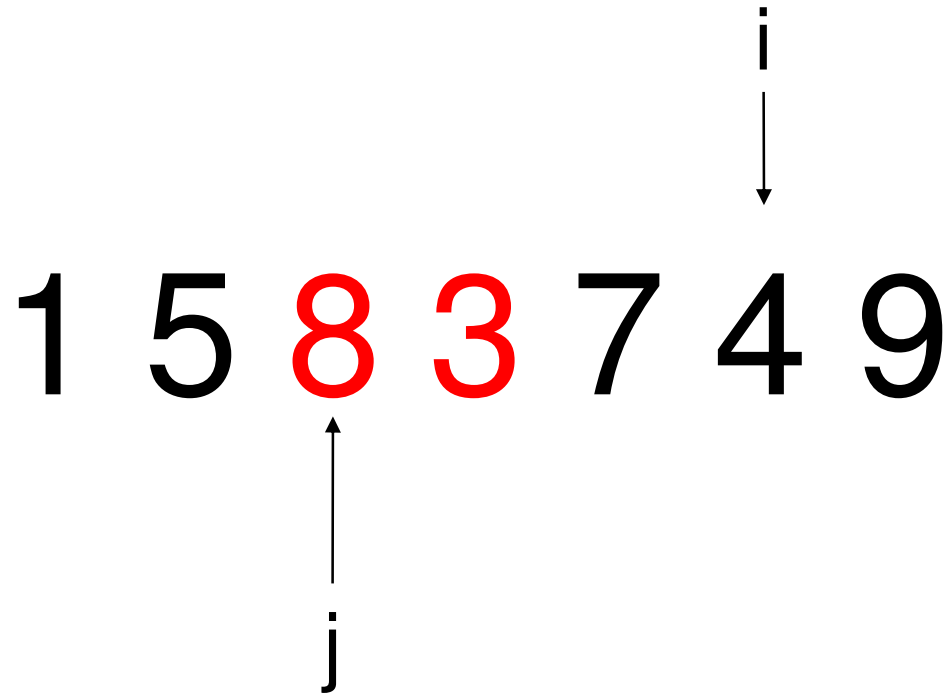
Insert Sort (melhoria)



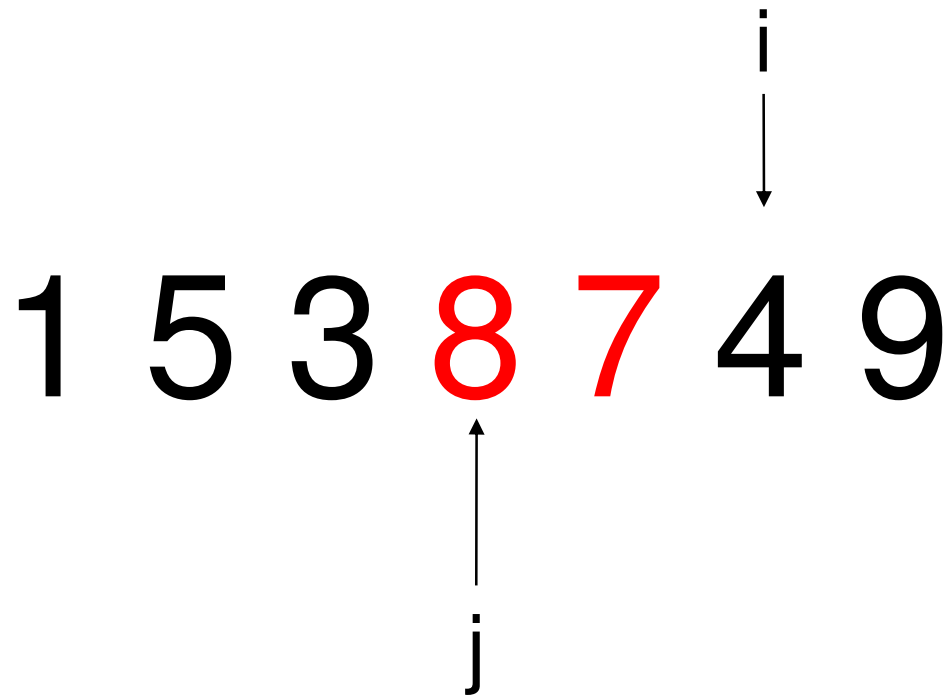
Insert Sort (melhoria)



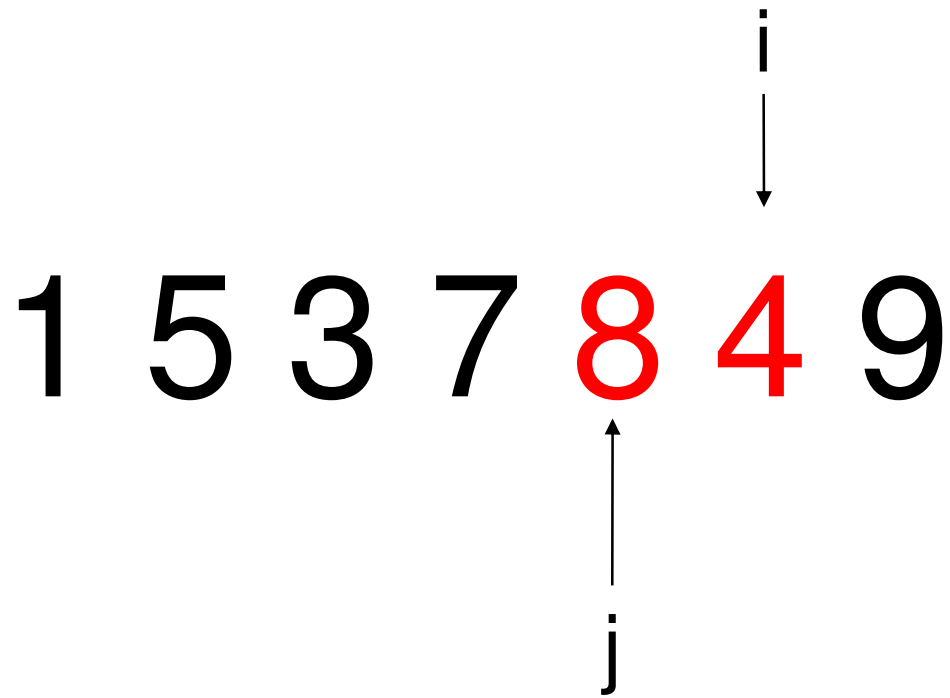
Insert Sort (melhoria)



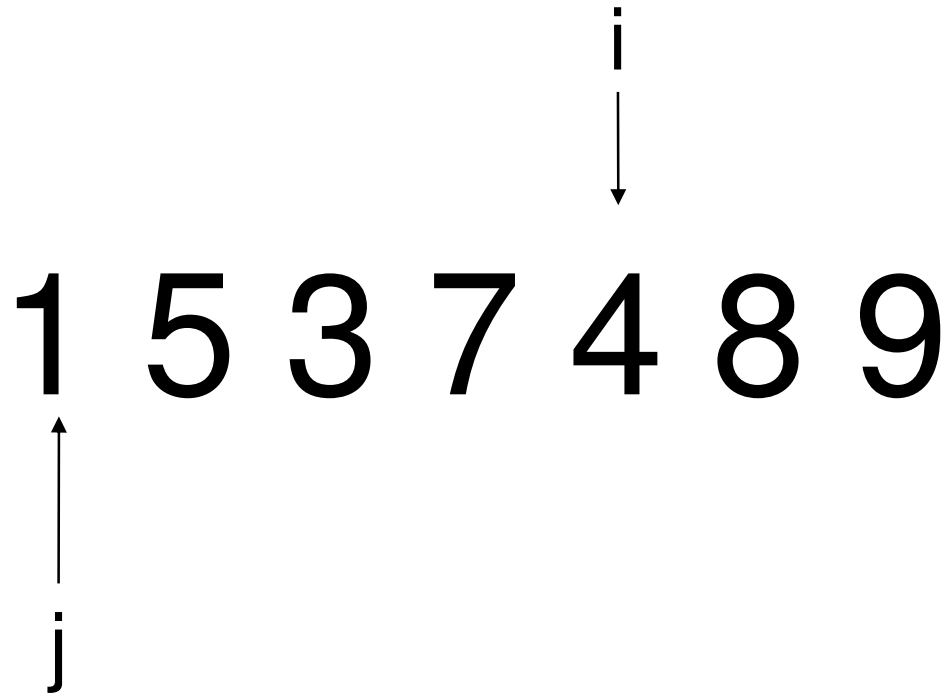
Insert Sort (melhoria)



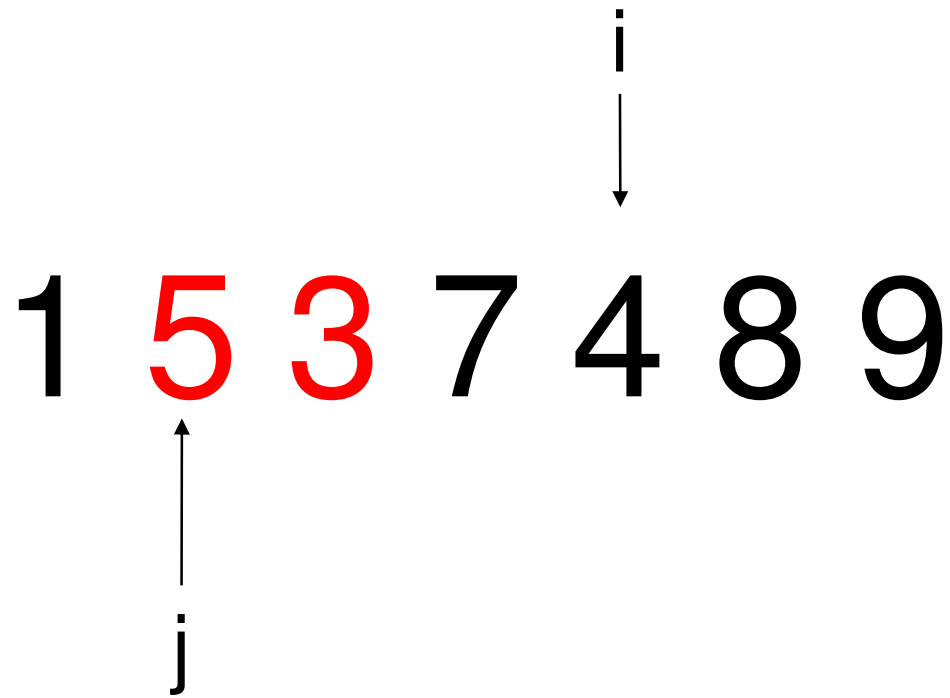
Insert Sort (melhoria)



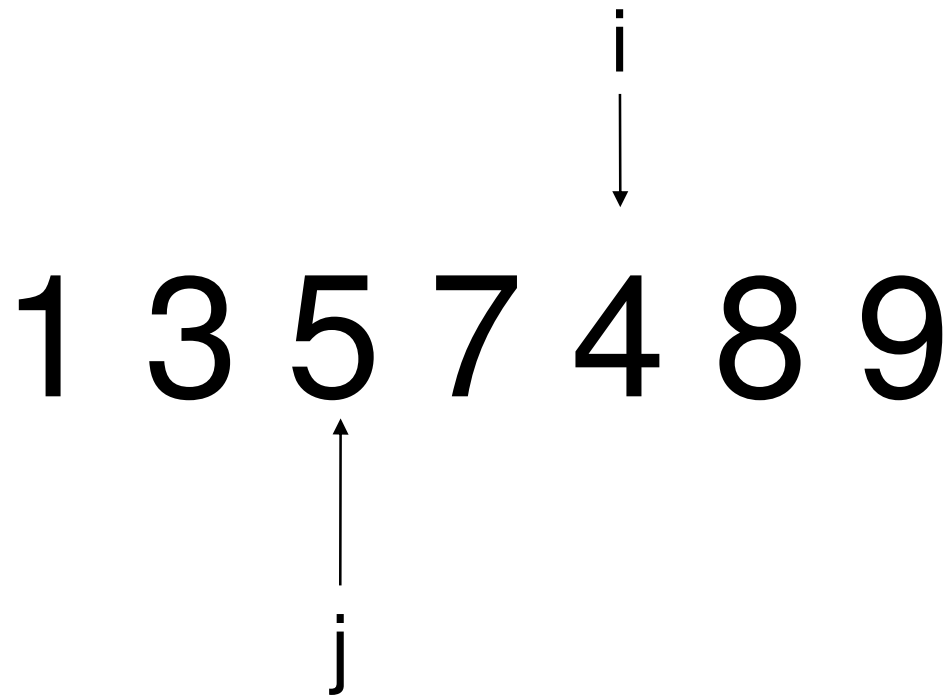
Insert Sort (melhoria)



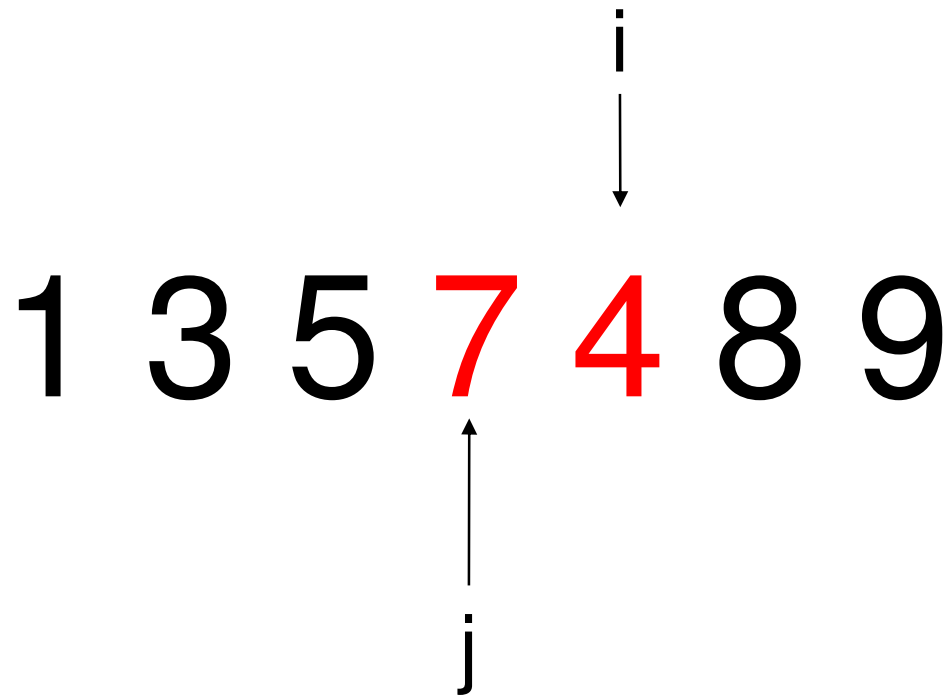
Insert Sort (melhoria)



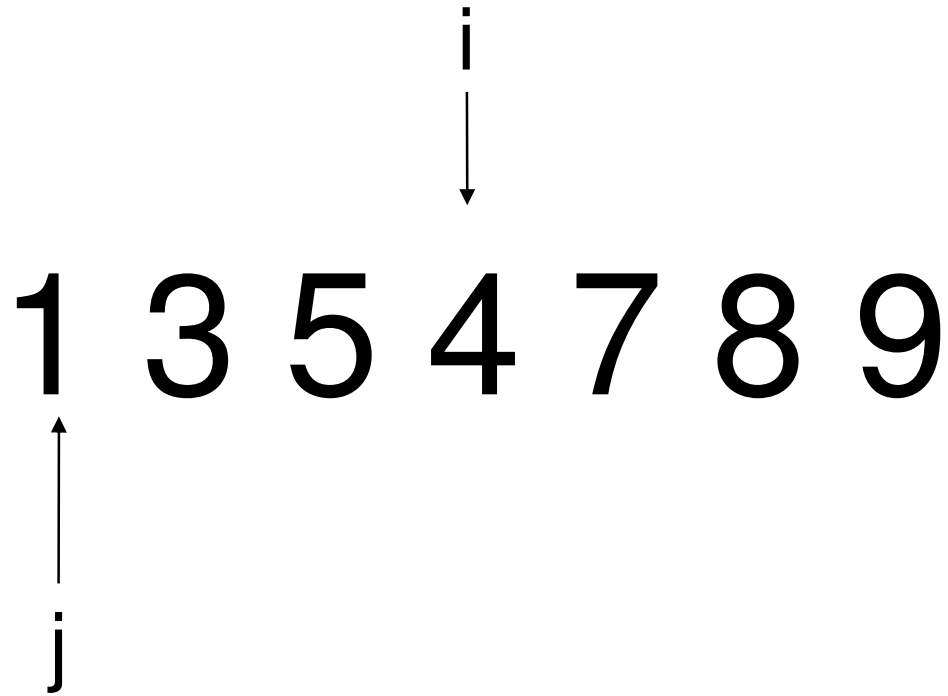
Insert Sort (melhoria)



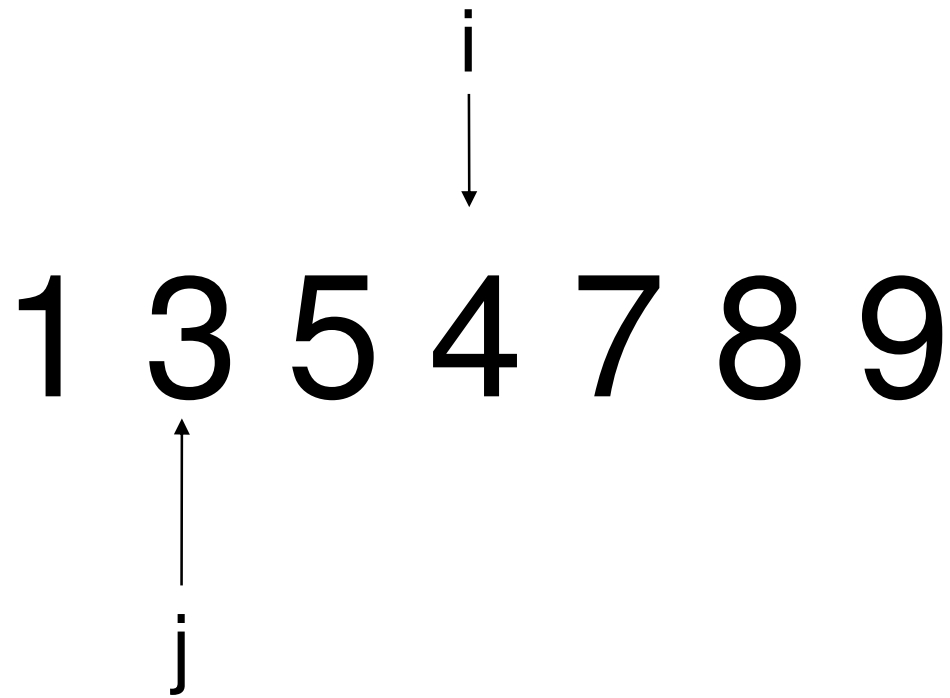
Insert Sort (melhoria)



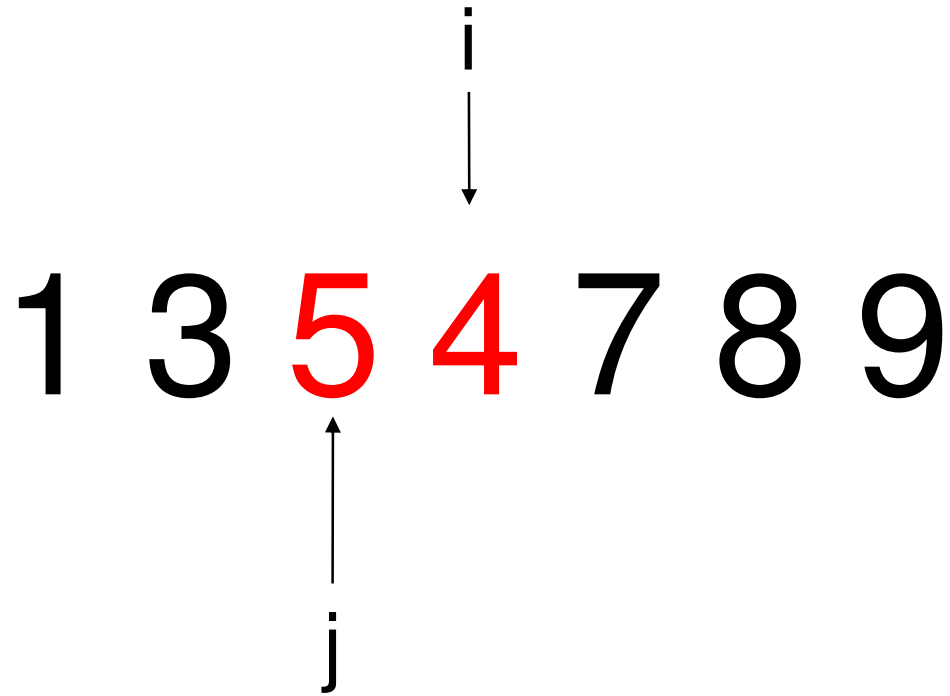
Insert Sort (melhoria)



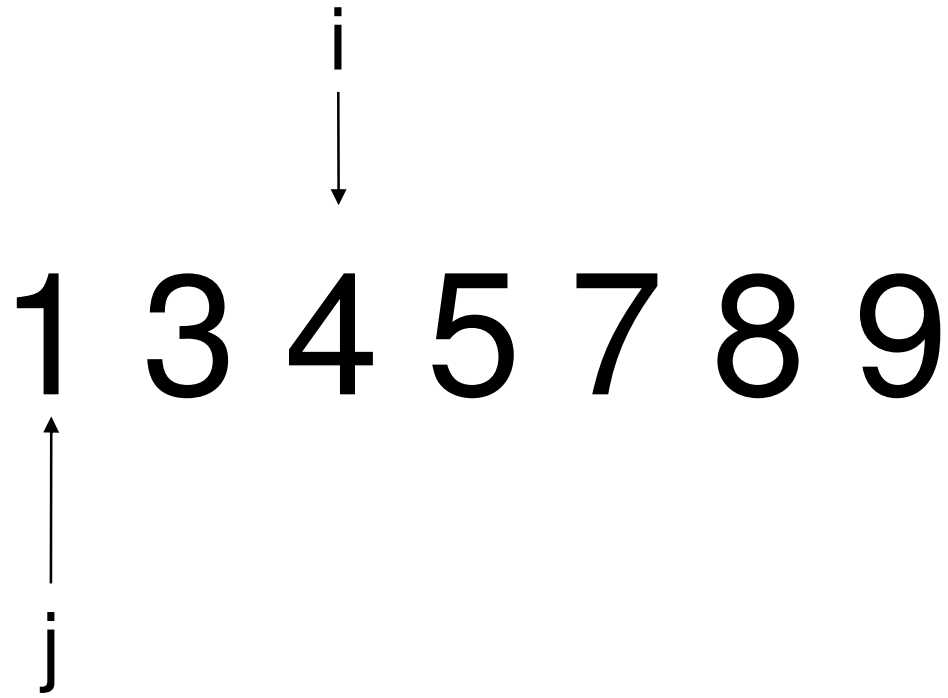
Insert Sort (melhoria)



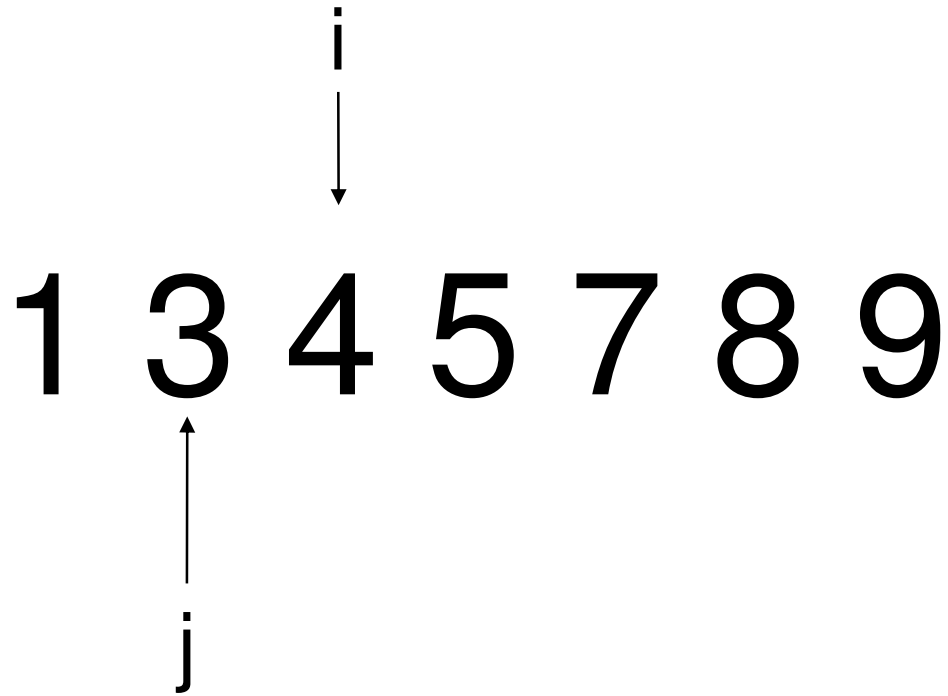
Insert Sort (melhoria)



Insert Sort (melhoria)



Insert Sort (melhoria)



Insert Sort (melhoria)

1 3 4 5 7 8 9

Mergesort

- Ordena inicialmente sub-vetores de tamanho 1 ou 2 e recursivamente dobra esse tamanho até ordenar o vetor total de tamanho N
- Baseia-se no merge de dois vetores ordenados



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Merge de dois vetores ordenados

- Dados dois vetores $V1$ e $V2$, já ordenados, o **merge** cria um outro vetor V ordenado com os elementos de $V1$ e $V2$

```
local i, j, k = 1, 1, 1
for k=1, #v1+#v2 do
    if j>#v2 then v[k]=v1[i]; i=i+1
    else if i>#v1 then v[k]=v2[j]; j=j+1
    else if v1[i]<=v2[j] then v[k]=v1[i]; i=i+1
    else v[k]=v2[j]; j=j+1 end
end
```



Merge de dois vectores

$V1 =$

$i \downarrow$ 1	3	7	18	20
---------------------	---	---	----	----

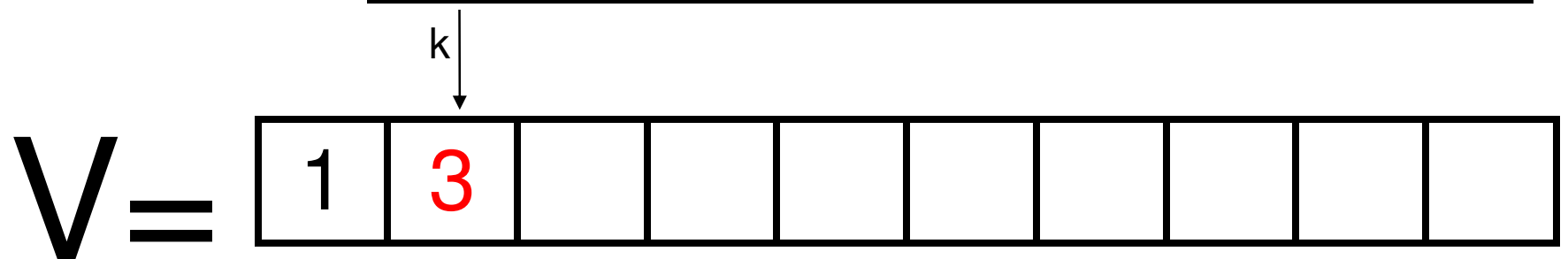
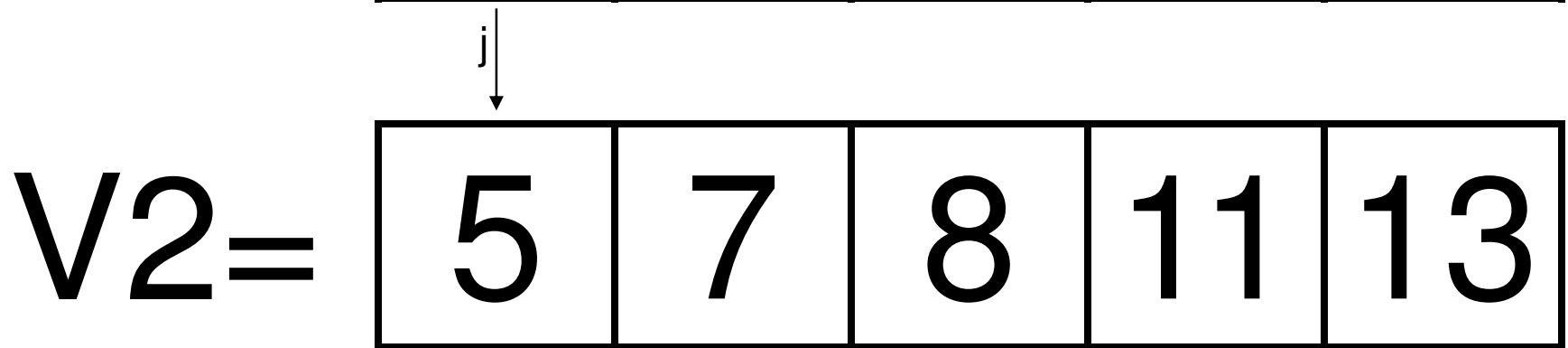
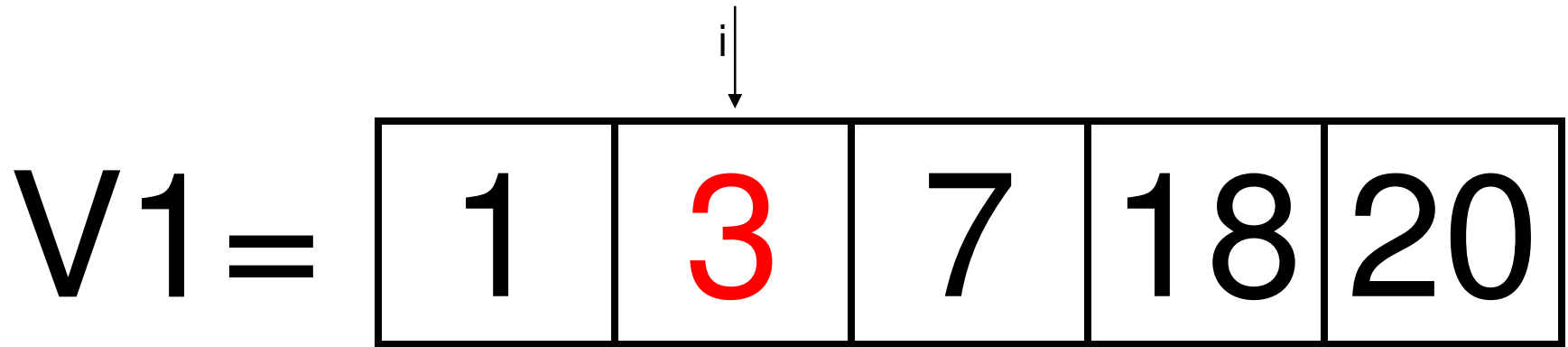
$V2 =$

$j \downarrow$ 5	7	8	11	13
---------------------	---	---	----	----

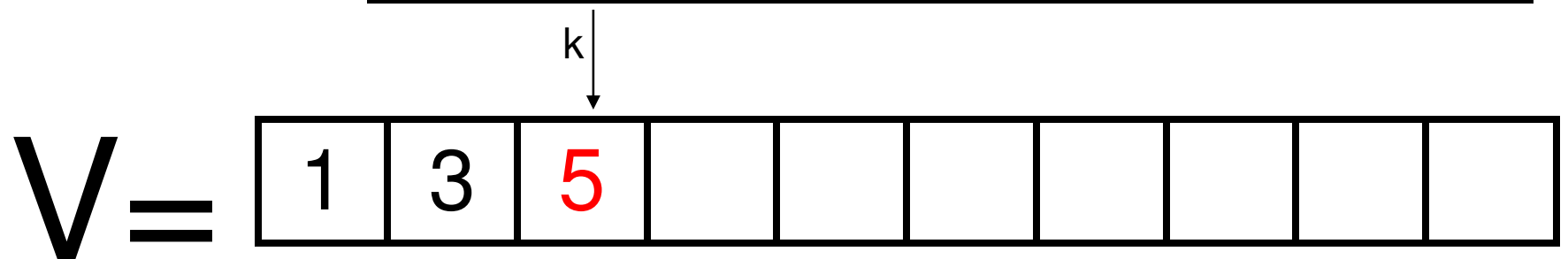
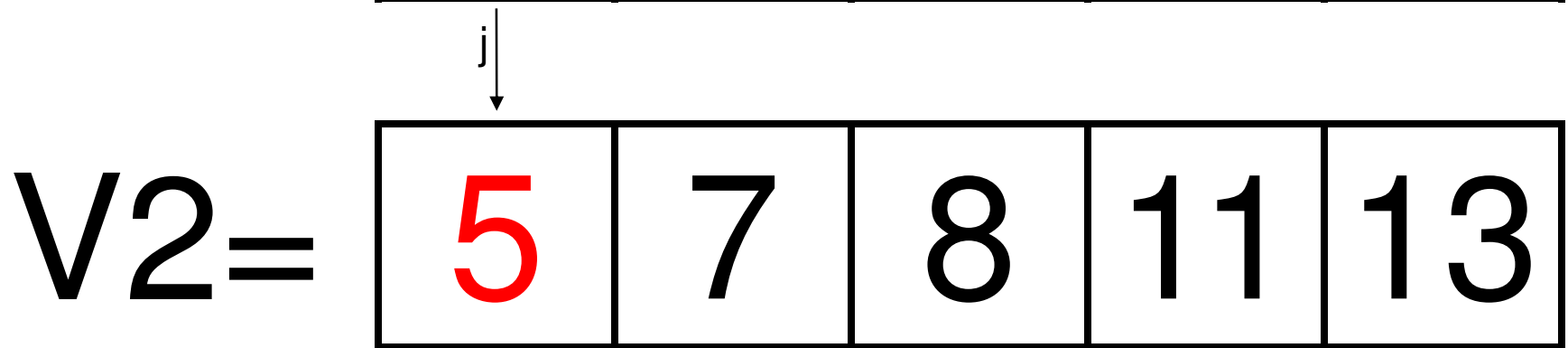
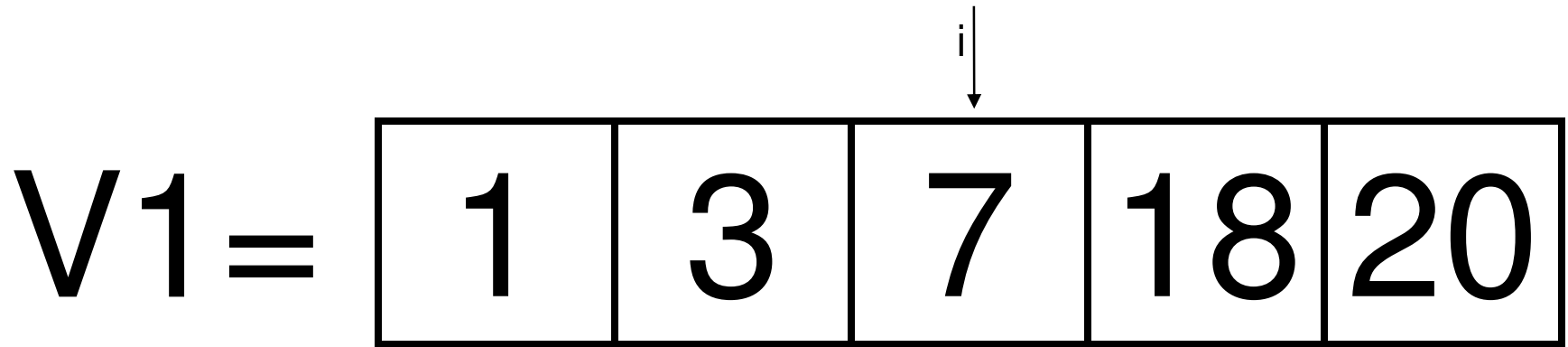
$V =$

$k \downarrow$ 1									
---------------------	--	--	--	--	--	--	--	--	--

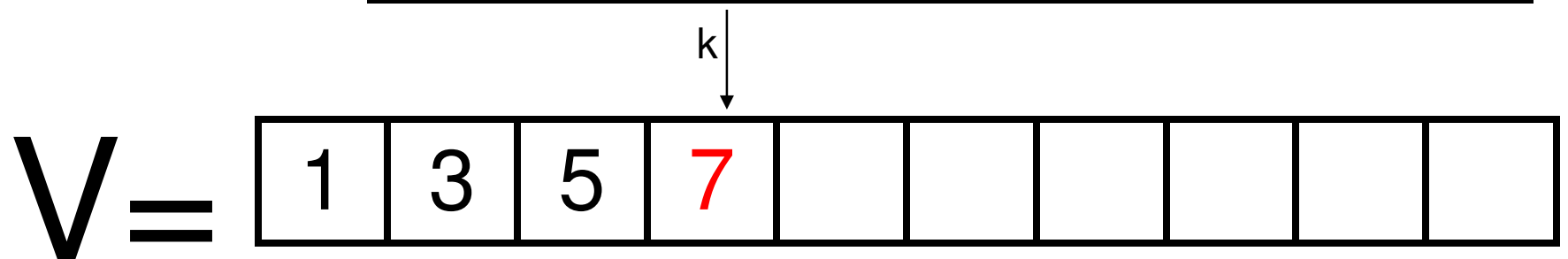
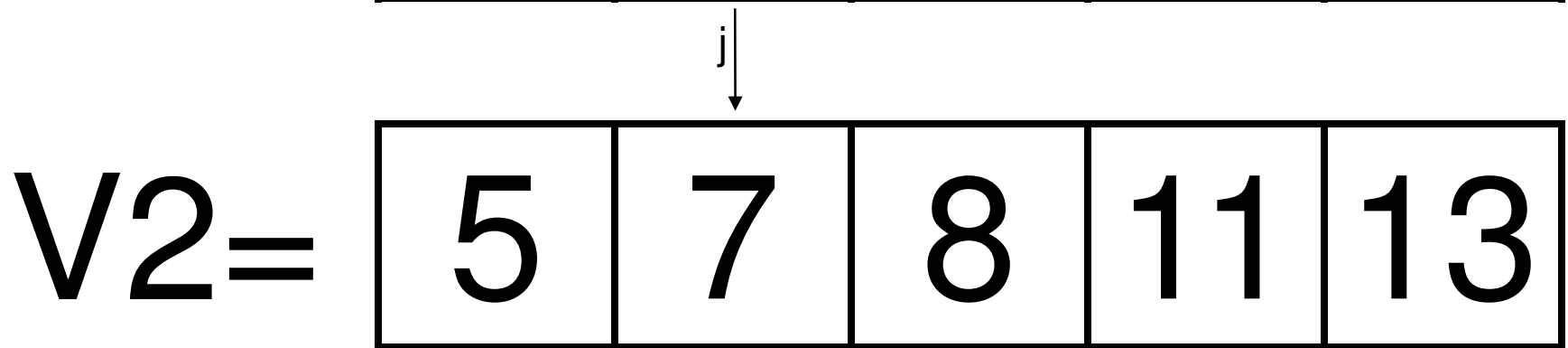
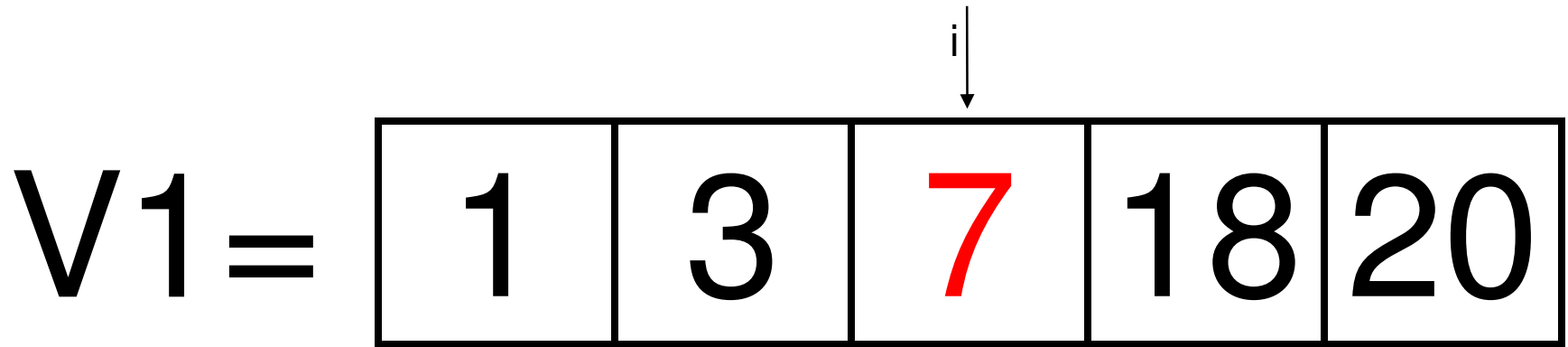
Merge de dois vectores



Merge de dois vectores



Merge de dois vectores



Merge de dois vectores

V1 =

1	3	7	18	20
---	---	---	----	----

V2 =

5	7	8	11	13
---	---	---	----	----

V =

1	3	5	7	7					
---	---	---	---	---	--	--	--	--	--

Merge de dois vectores

V1 =

1	3	7	18	20
---	---	---	----	----

V2 =

5	7	8	11	13
---	---	---	----	----

V =

1	3	5	7	7	8				
---	---	---	---	---	---	--	--	--	--

Merge de dois vectores

$V1 =$

1	3	7	18	20
---	---	---	----	----

$V2 =$

5	7	8	11	13
---	---	---	----	----

$V =$

1	3	5	7	7	8	11			
---	---	---	---	---	---	----	--	--	--

Merge de dois vectores

V1 =

1	3	7	18	20
---	---	---	----	----

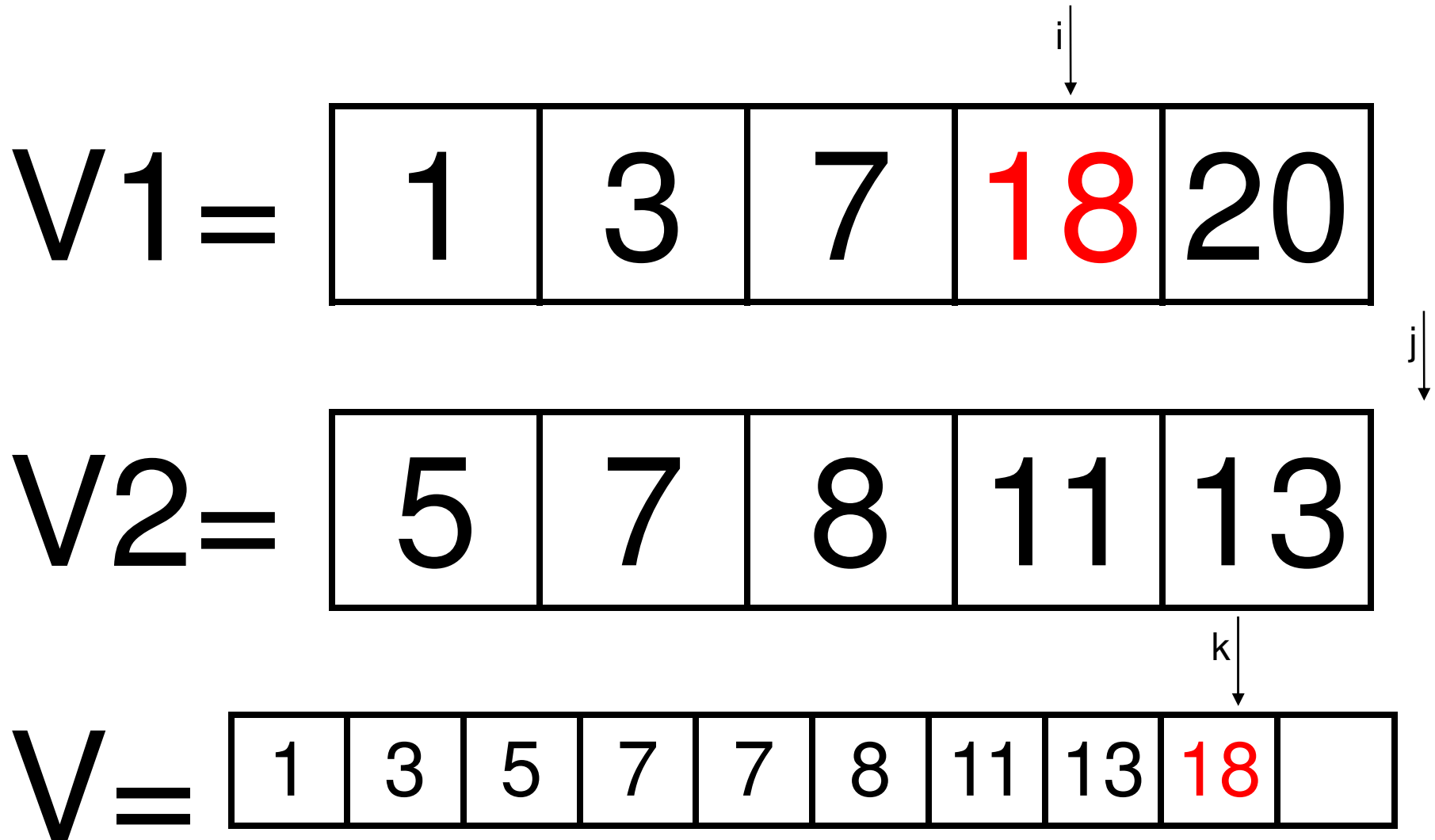
V2 =

5	7	8	11	13
---	---	---	----	----

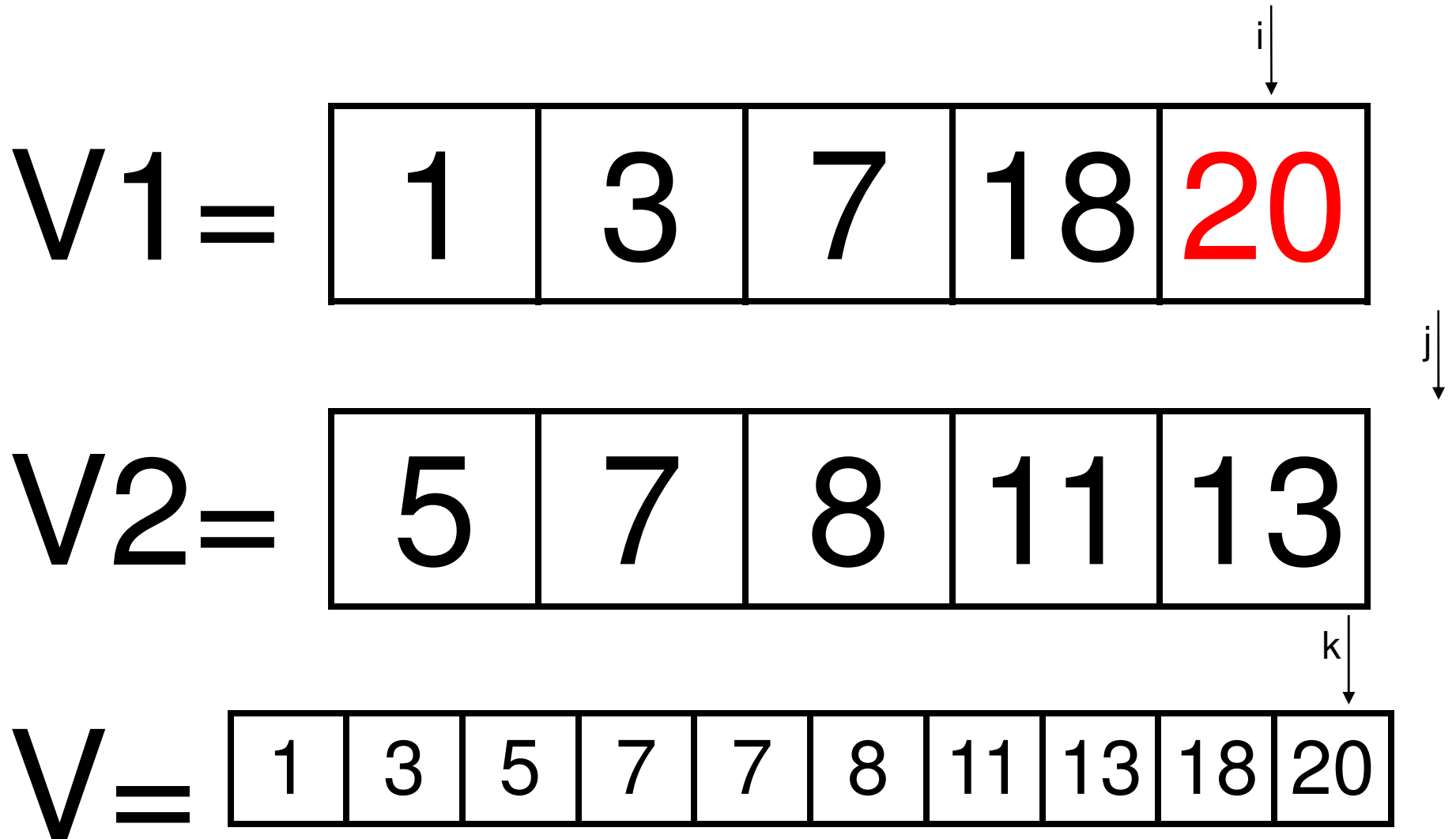
V =

1	3	5	7	7	8	11	13		
---	---	---	---	---	---	----	----	--	--

Merge de dois vectores



Merge de dois vectores



Merge de dois vectores

V1 =

1	3	7	18	20
---	---	---	----	----

i ↓

V2 =

5	7	8	11	13
---	---	---	----	----

j ↓

V =

1	3	5	7	7	8	11	13	18	20
---	---	---	---	---	---	----	----	----	----

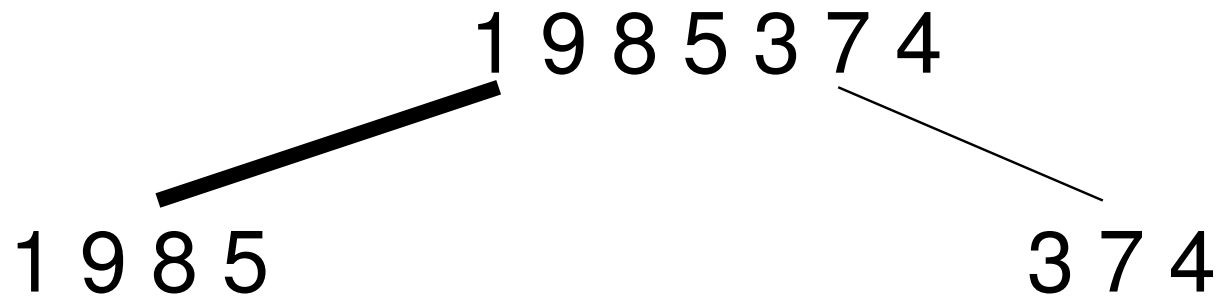
k ↓

Mergesort

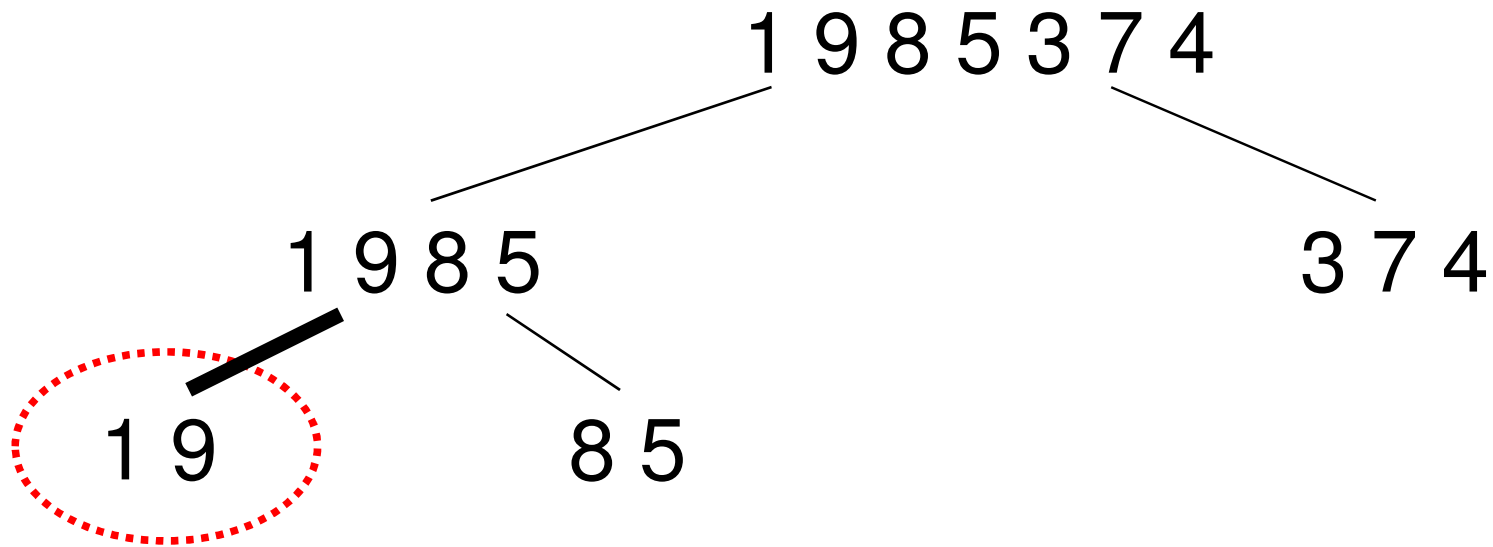
1 9 8 5 3 7 4



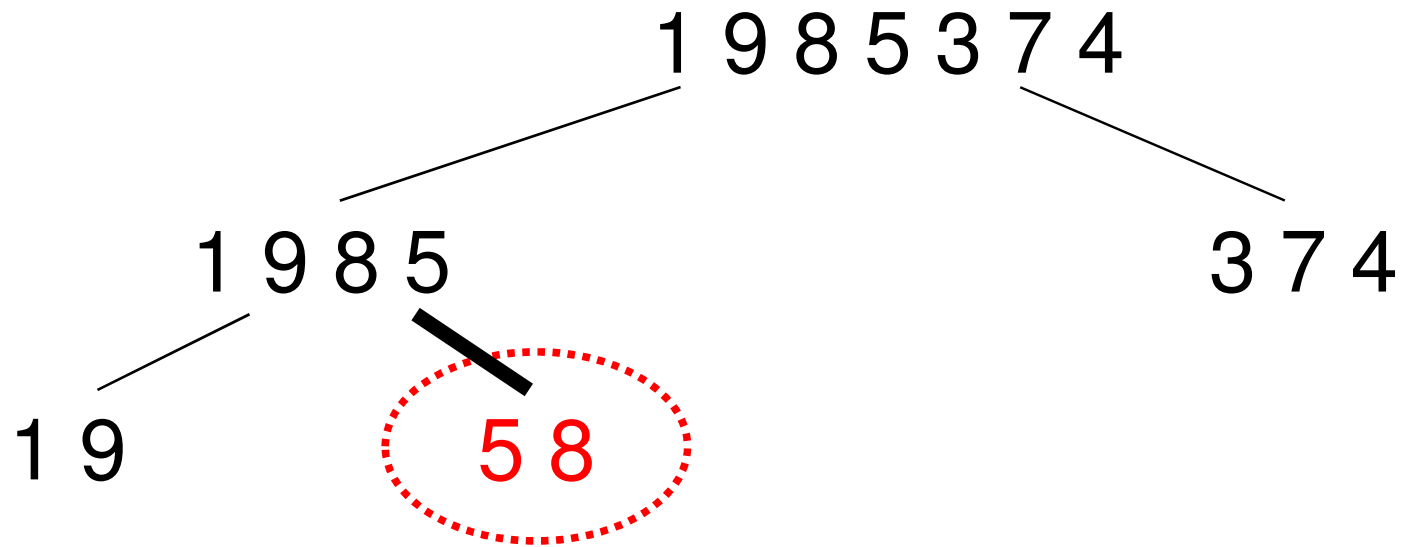
Mergesort



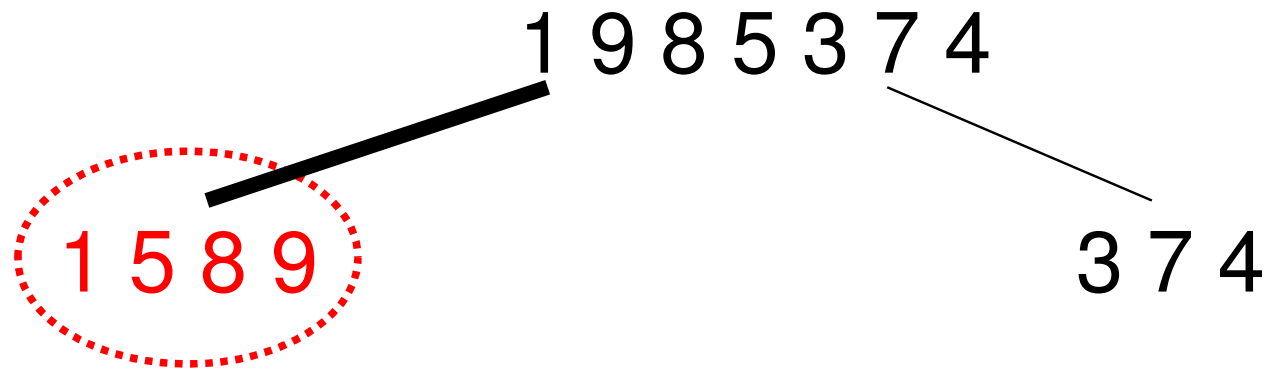
Mergesort



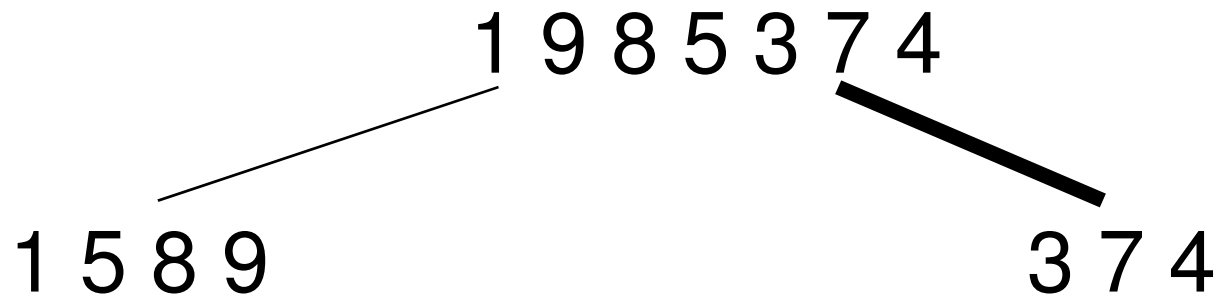
Mergesort



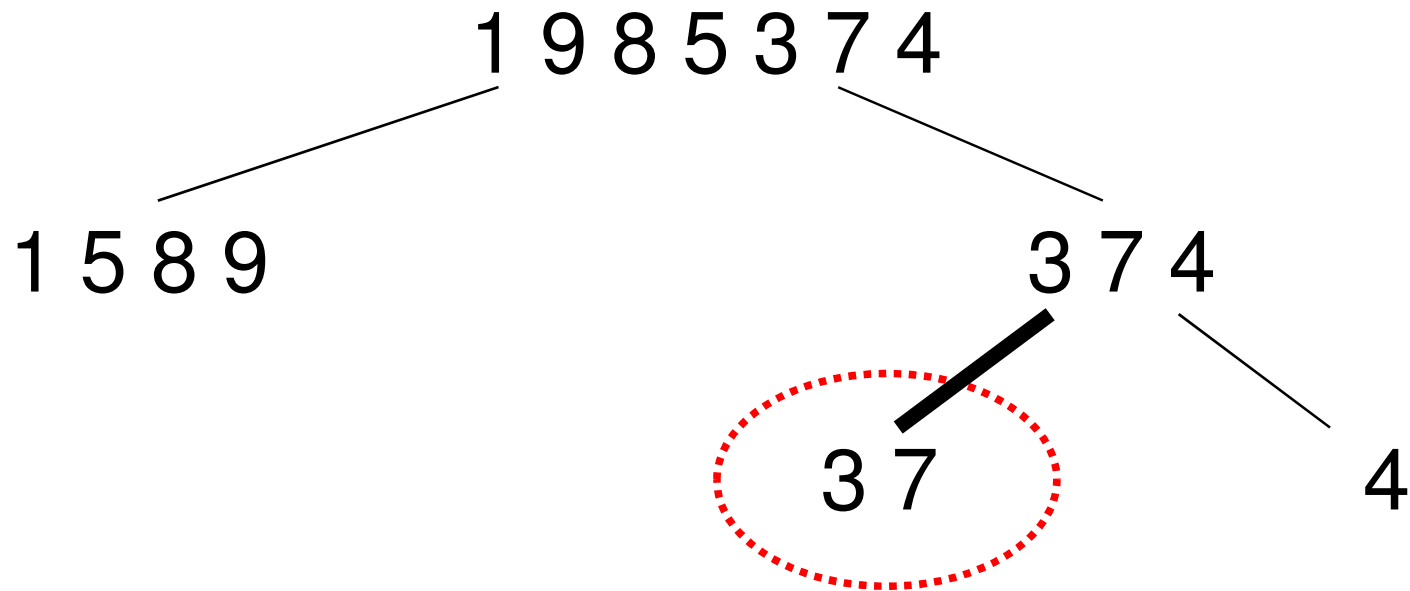
Mergesort



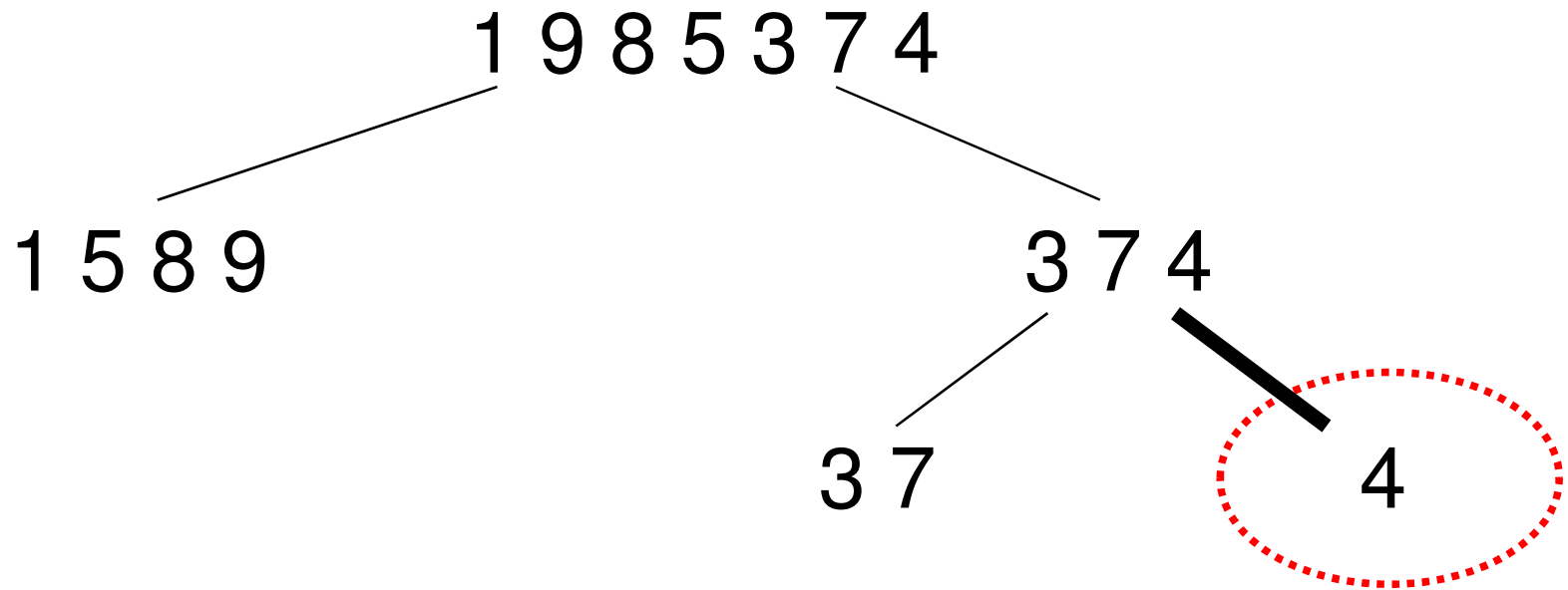
Mergesort



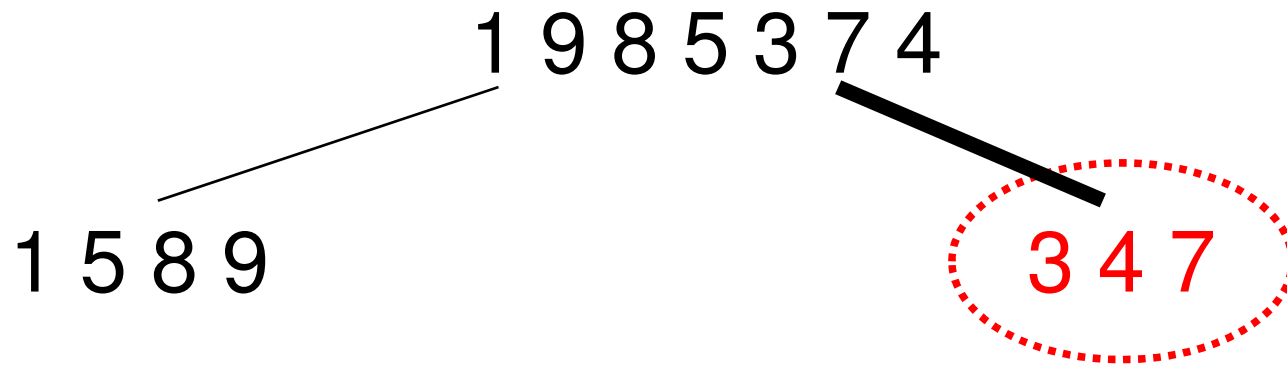
Mergesort



Mergesort



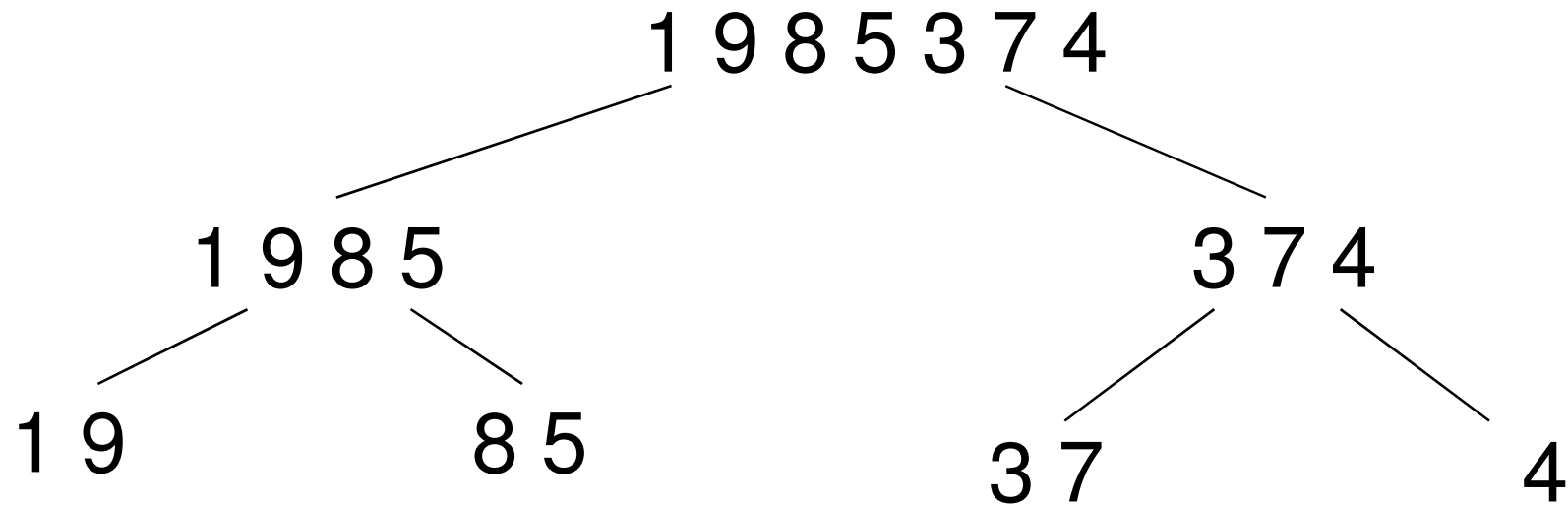
Mergesort



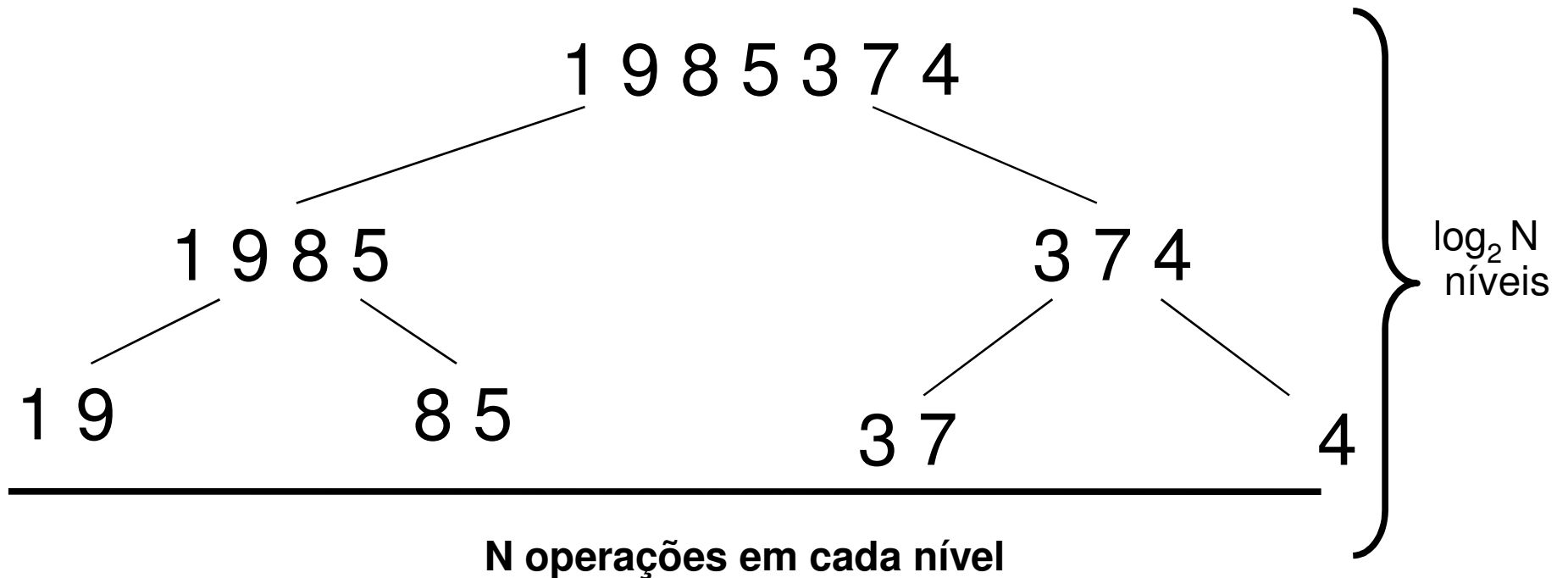
Mergesort

1 3 4 5 7 8 9

Mergesort - operações



Mergesort - operações



Total de operações:
 $N * \log_2 N$

Heapsort

- Ordenação baseada na estrutura de dados *heap*
 - A *heap* é um vetor que pode ser visto como uma árvore binária completa
 - Cada nó da árvore equivale a um valor no vetor



Primitivas da Heap

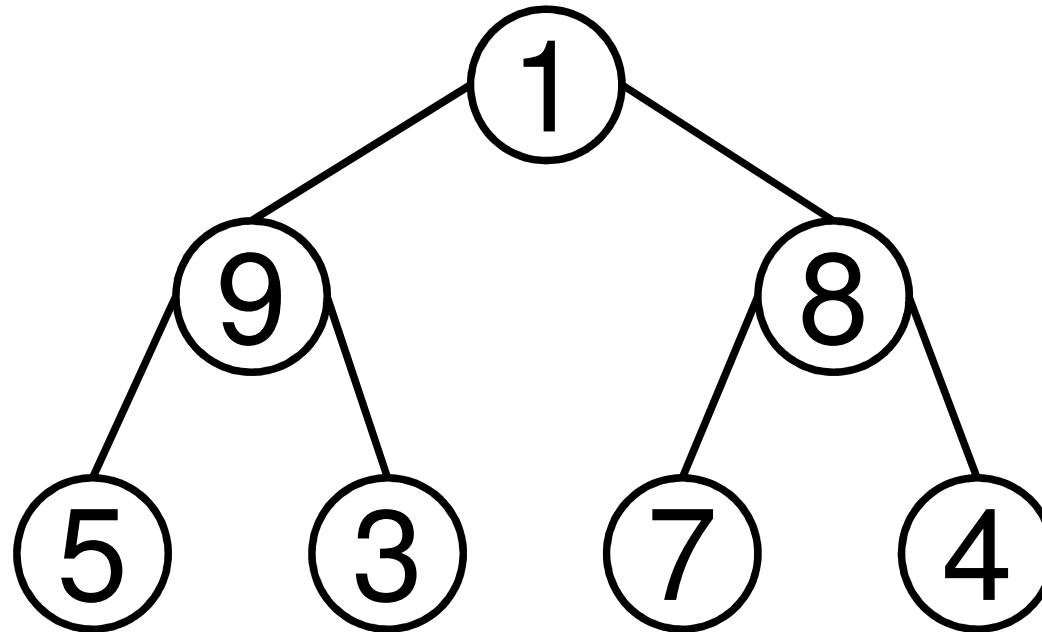
- `length(V)`
 - número de elementos no vetor V
- `heap-size(V)`
 - número de elementos da *heap* armazenados no vetor V

Vetor Heap – primitivas de árvore

- `parent(i)`
 - return $i/2$
- `left(i)`
 - return $2i$
- `right(i)`
 - return $2i+1$

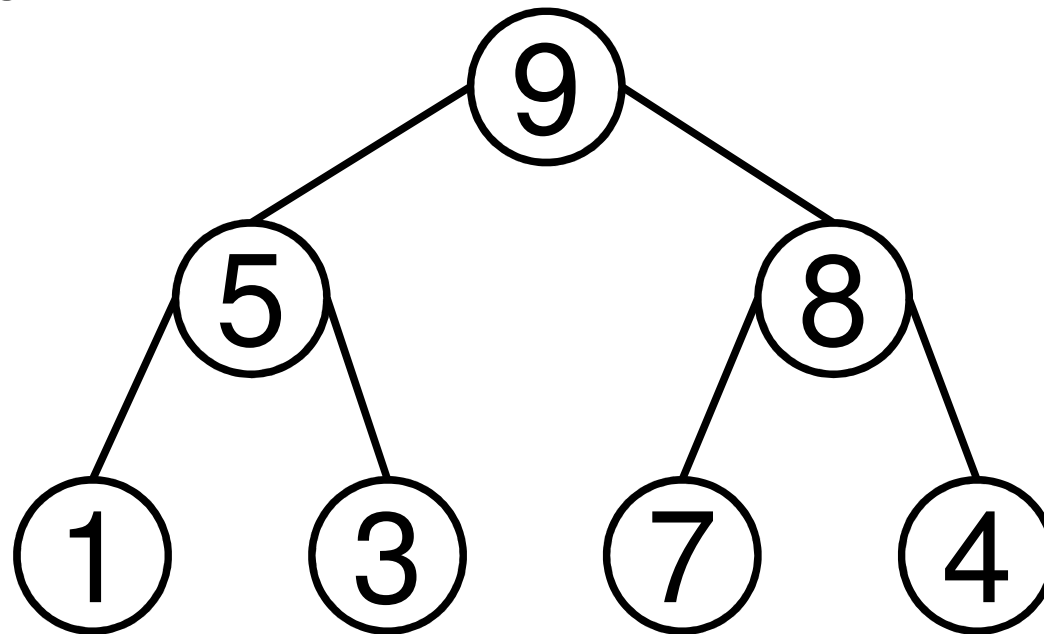
Heap como um vetor

1	9	8	5	3	7	4
---	---	---	---	---	---	---



Propriedade da Heap

- Para cada nó i que não seja a raiz, então:
 - $V[\text{parent}(i)] \geq V[i]$
- Exemplo:



Primitiva Heapify

- Função para manter a propriedade da heap

```
function heapify(V, i)
    l = left(i)
    r = right(i)

    if l<=heap-size(V) and V[l]>V[i] then largest = l
        else largest = i end

    if r<=heap-size(V) and V[r]>V[largest]
        then largest=r end

    if largest != i then
        V[i], V[largest] = V[largest], V[i]
        heapify(V, largest)
    end
end
```



Construindo uma Heap

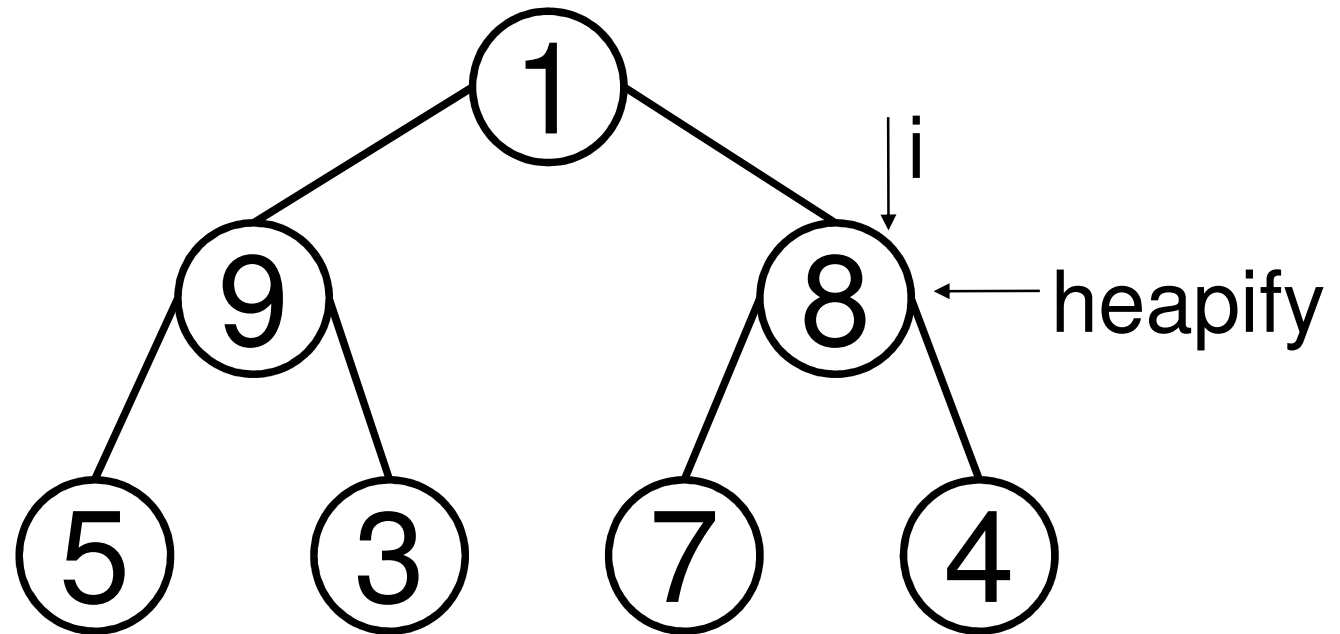
- Primitiva que constrói uma *heap* a partir de um vetor de N elementos

```
function build_heap(V)
    heap-size(V) = length(V) -- #V
    for i=length(V) / 2, 1, -1 do
        heapify(V, i)
    end
end
```



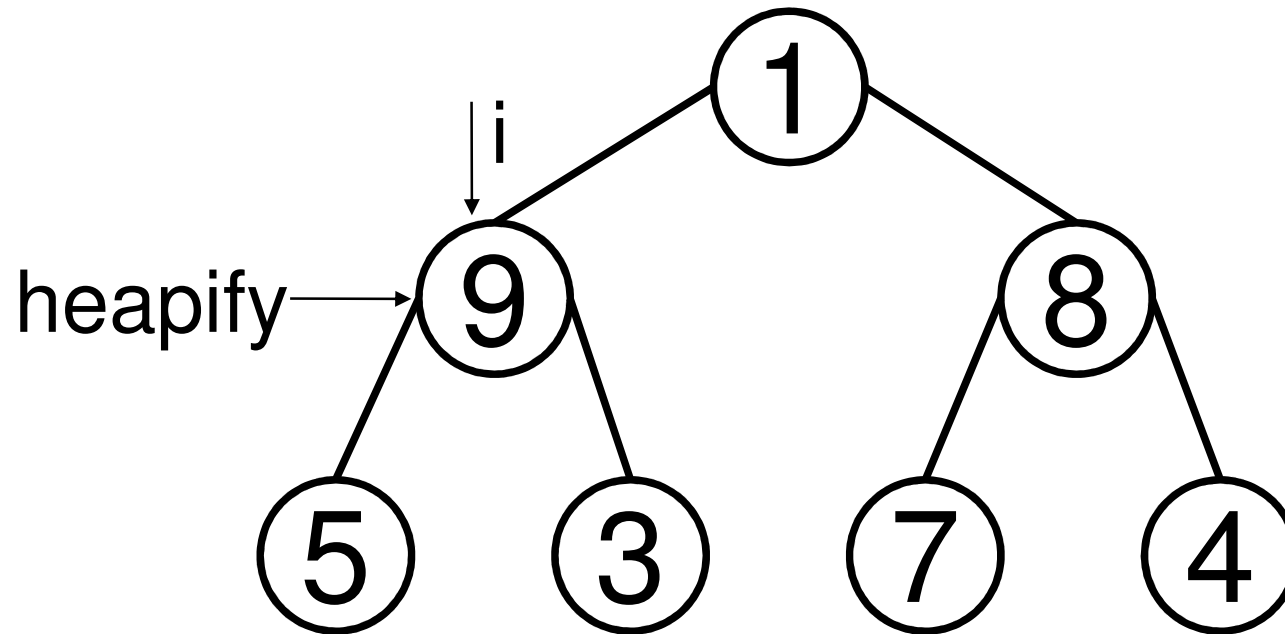
Construindo uma Heap

1	9	8	5	3	7	4
---	---	---	---	---	---	---



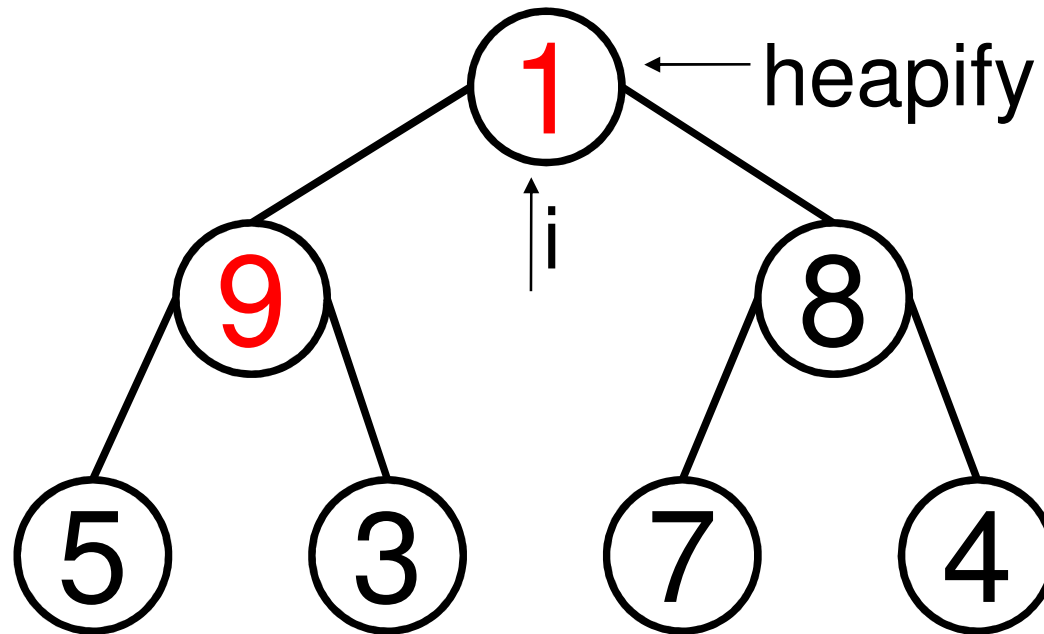
Construindo uma Heap

1	9	8	5	3	7	4
---	---	---	---	---	---	---



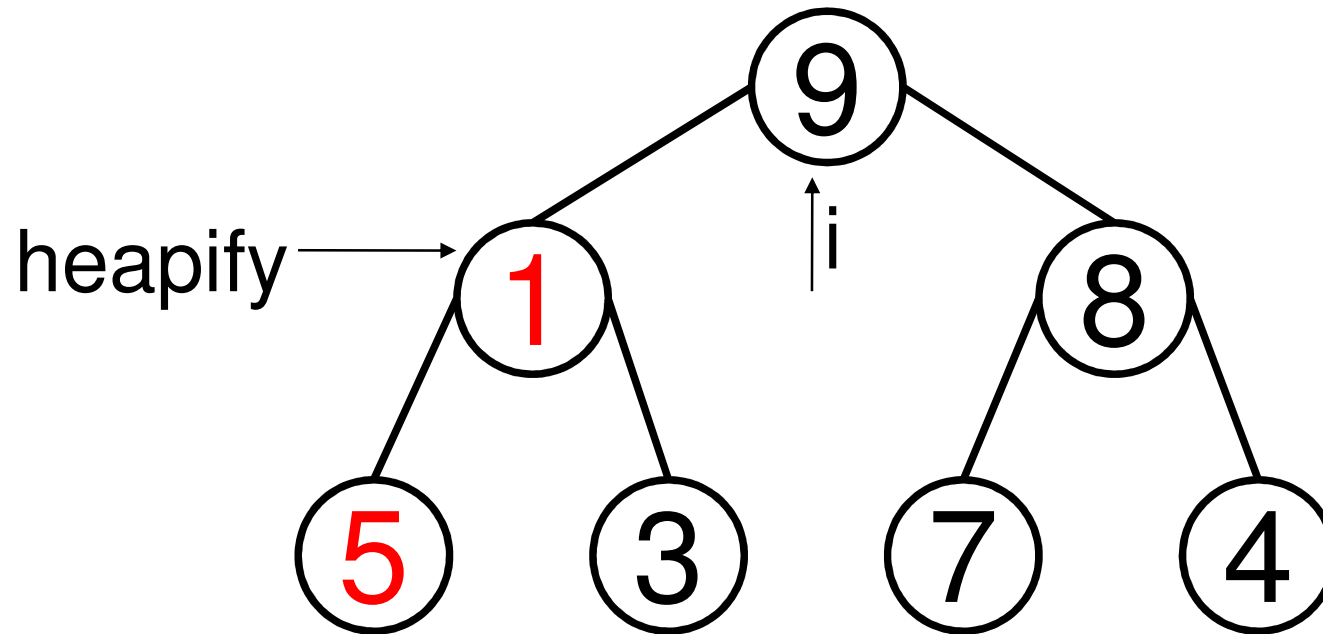
Construindo uma Heap

1	9	8	5	3	7	4
---	---	---	---	---	---	---



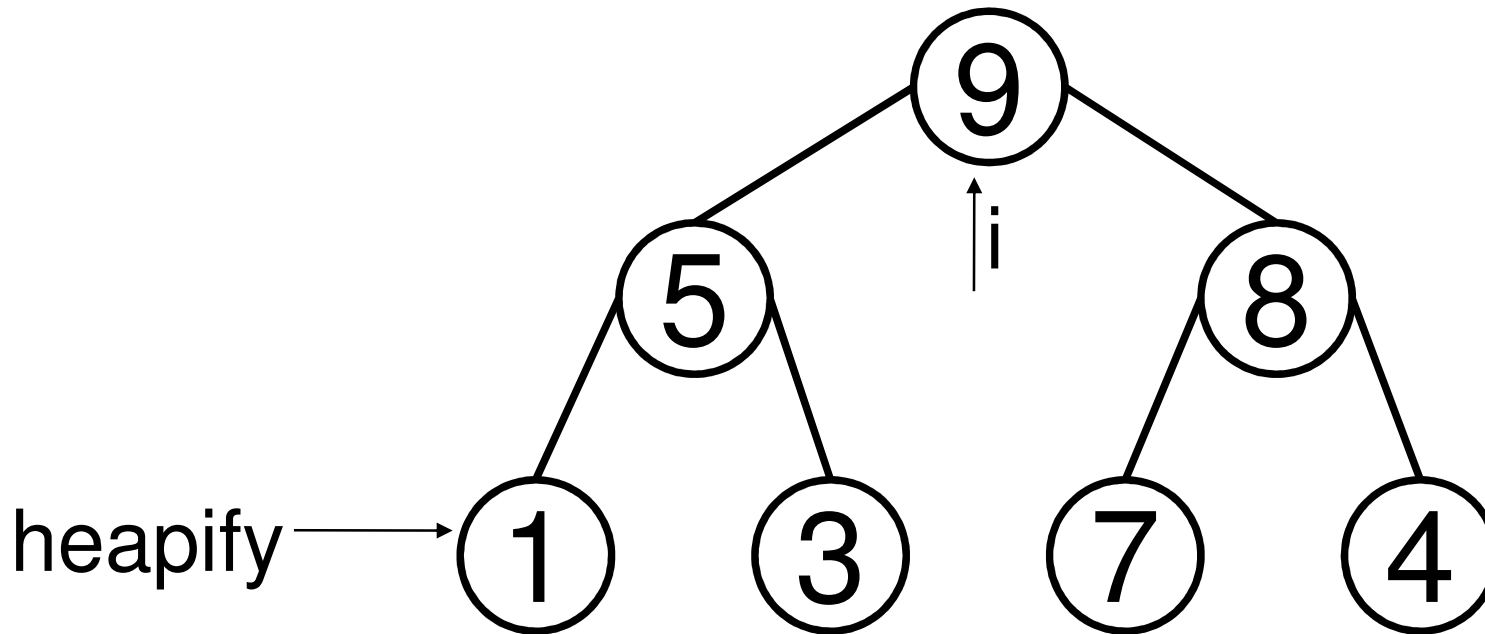
Construindo uma Heap

9	1	8	5	3	7	4
---	---	---	---	---	---	---



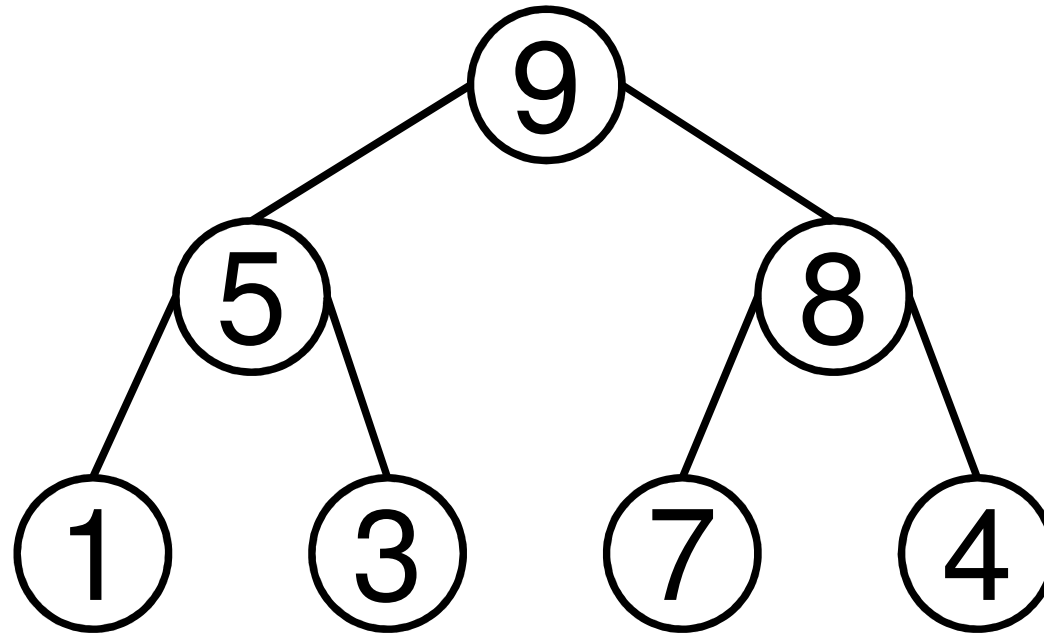
Construindo uma Heap

9	5	8	1	3	7	4
---	---	---	---	---	---	---



Construindo uma Heap

9	5	8	1	3	7	4
---	---	---	---	---	---	---



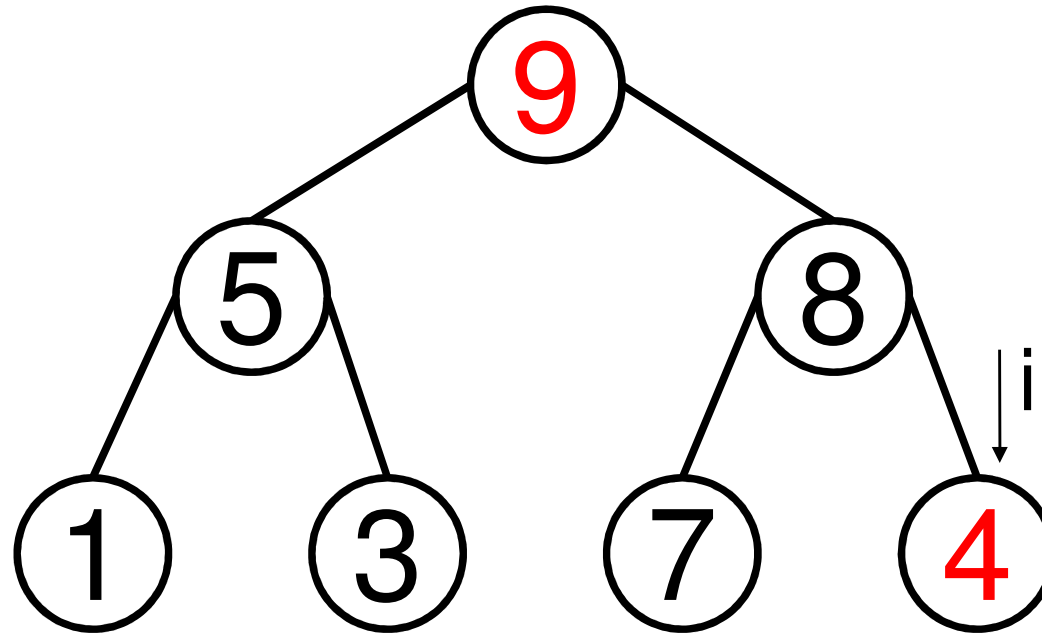
Heapsort

```
function heapsort(V)
    for i=length(V), 2, -1 do
        V[1], V[i] = V[i], V[1]
        heap-size(V) = heap-size(V) -
1
        heapify(V, 1)
    end
end
```



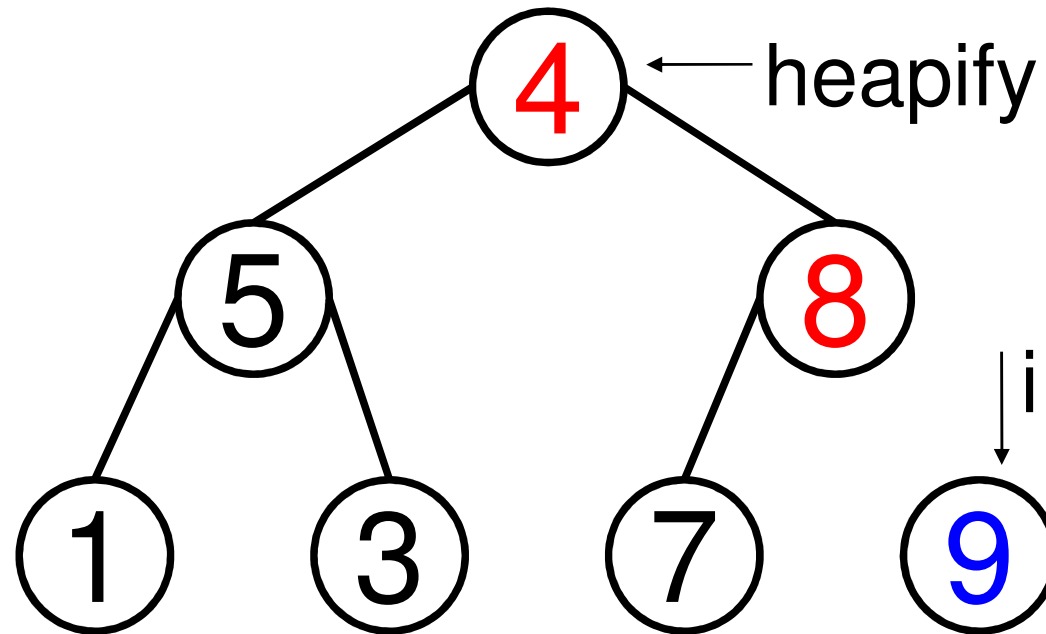
Heapsort

9	5	8	1	3	7	4
---	---	---	---	---	---	---



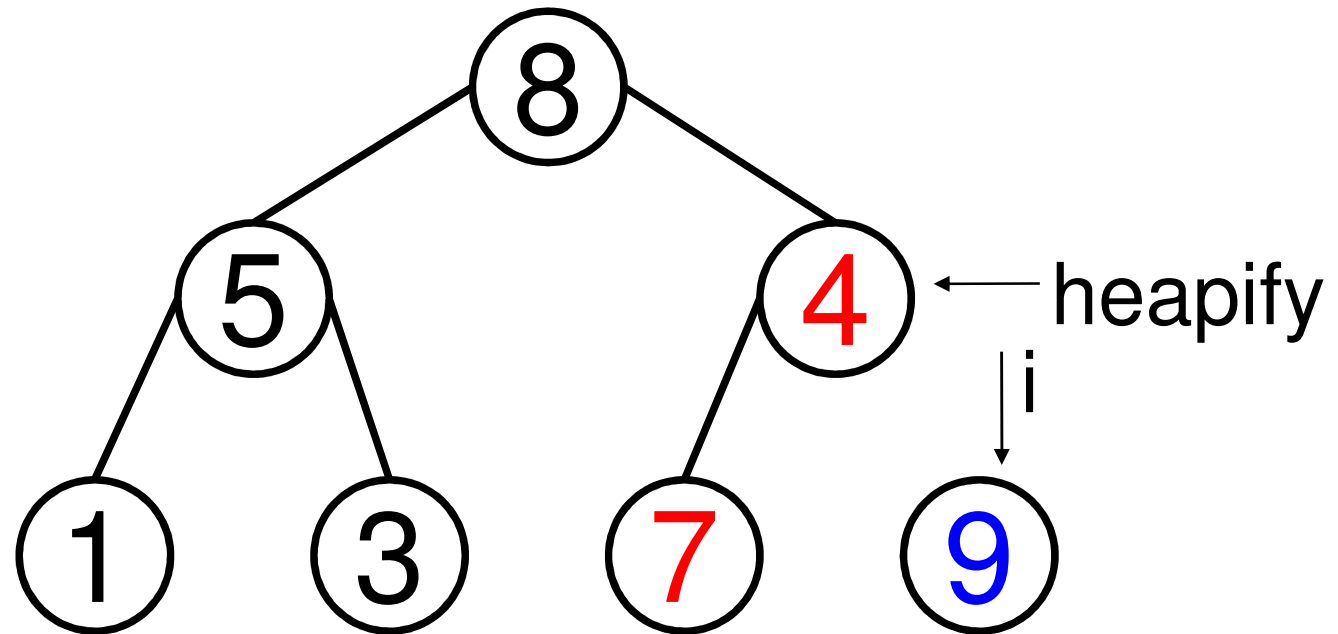
Heapsort

4	5	8	1	3	7	9
---	---	---	---	---	---	---



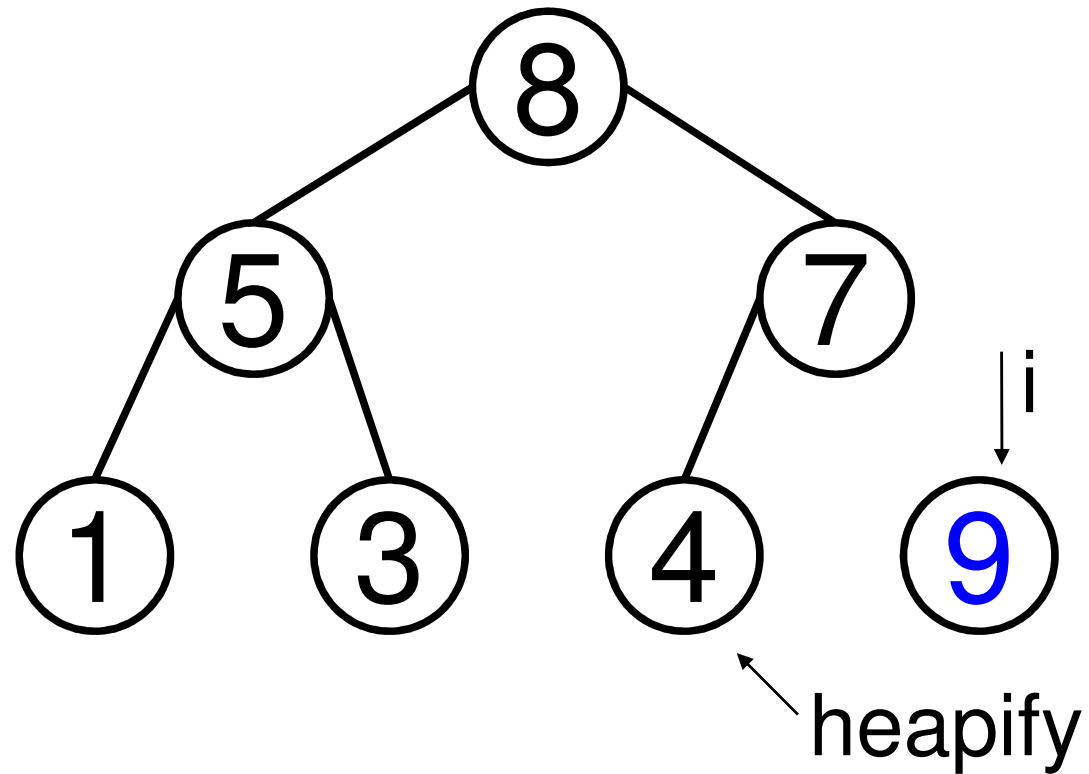
Heapsort

8	5	4	1	3	7	9
---	---	---	---	---	---	---



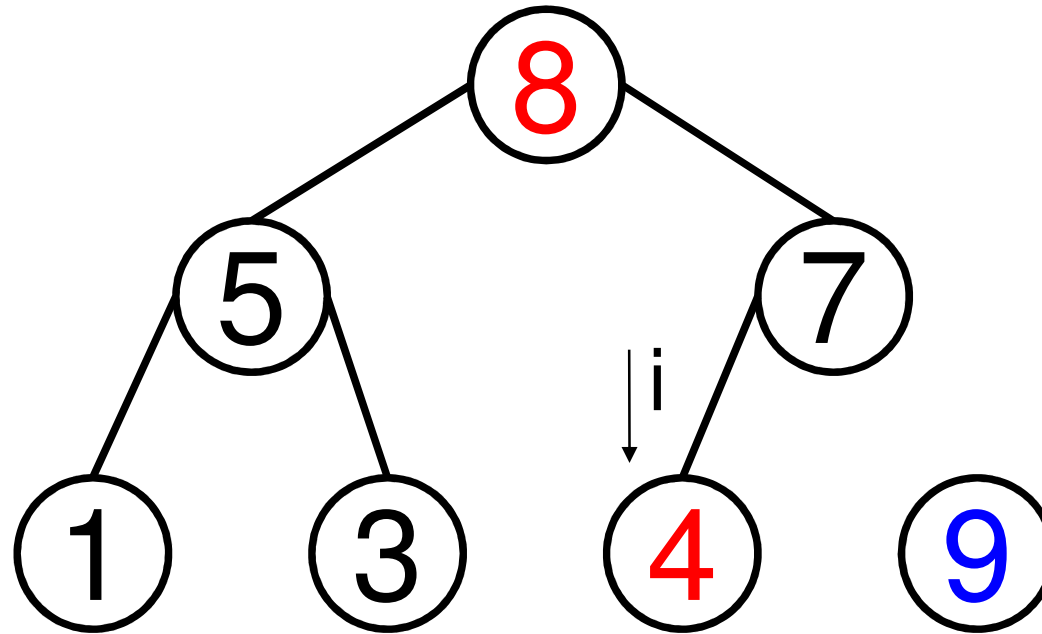
Heapsort

8	5	7	1	3	4	9
---	---	---	---	---	---	---



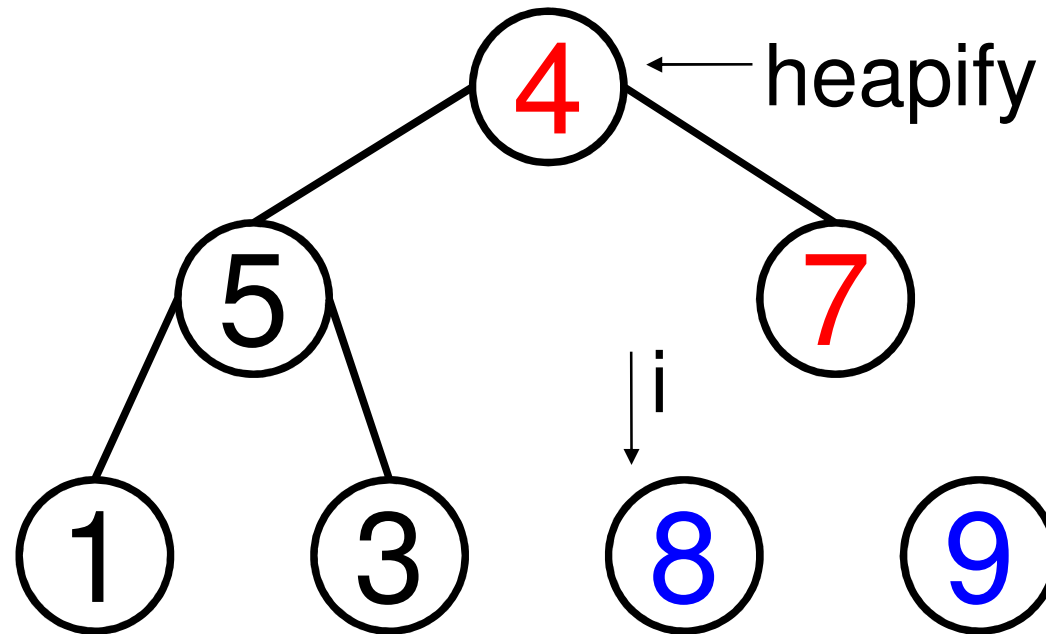
Heapsort

8	5	7	1	3	4	9
---	---	---	---	---	---	---



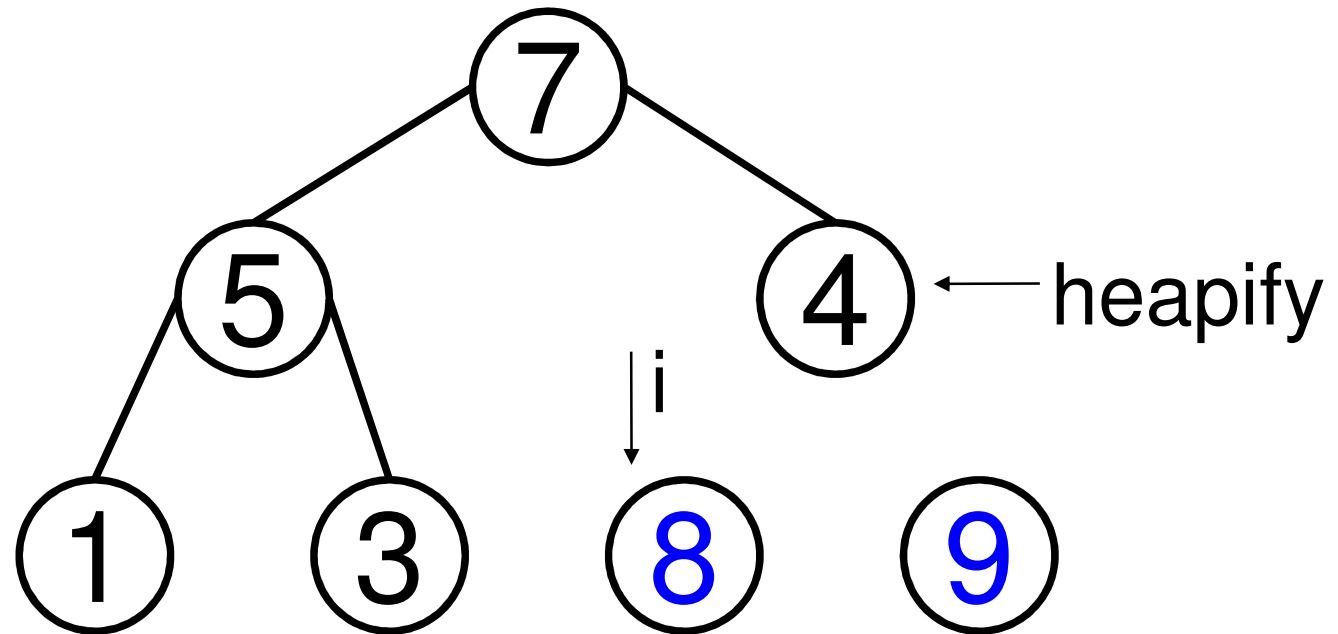
Heapsort

4	5	7	1	3	8	9
---	---	---	---	---	---	---



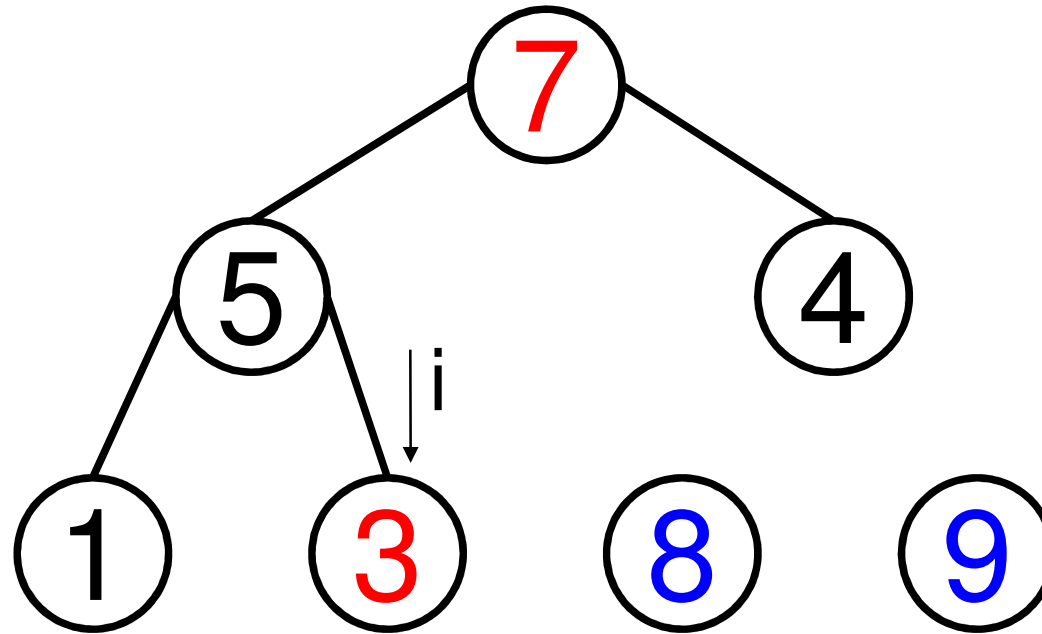
Heapsort

7	5	4	1	3	8	9
---	---	---	---	---	---	---



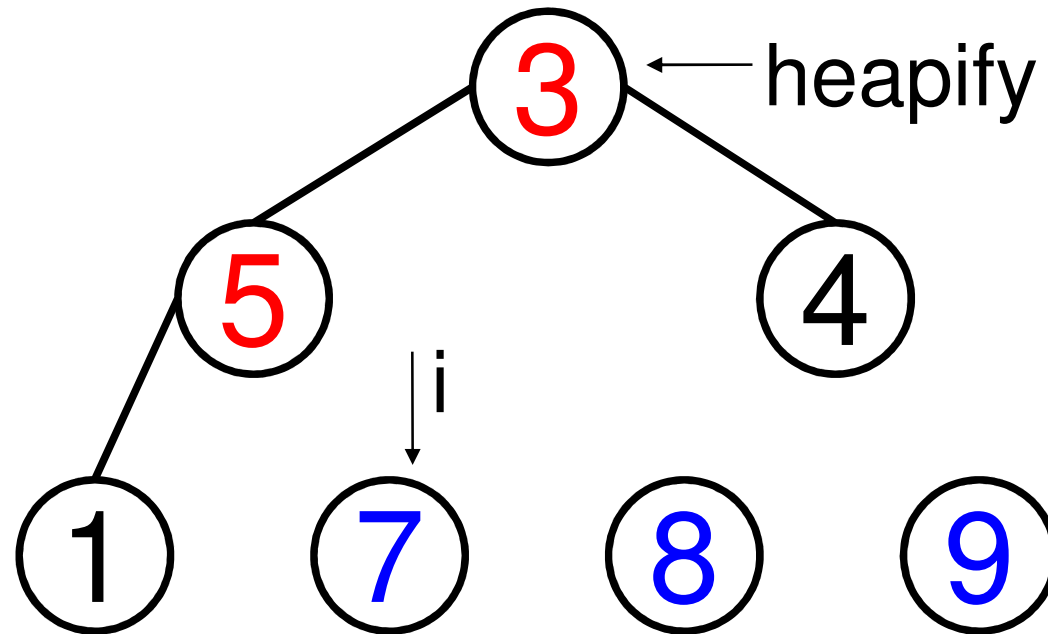
Heapsort

7	5	4	1	3	8	9
---	---	---	---	---	---	---



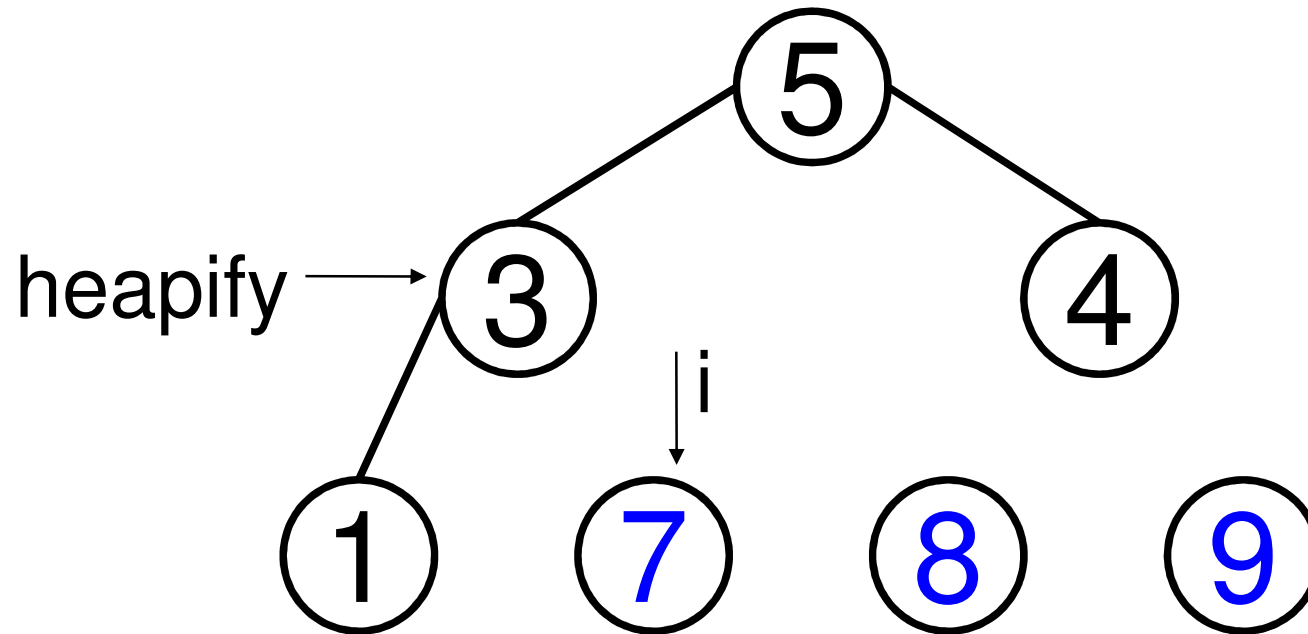
Heapsort

3	5	4	1	7	8	9
---	---	---	---	---	---	---



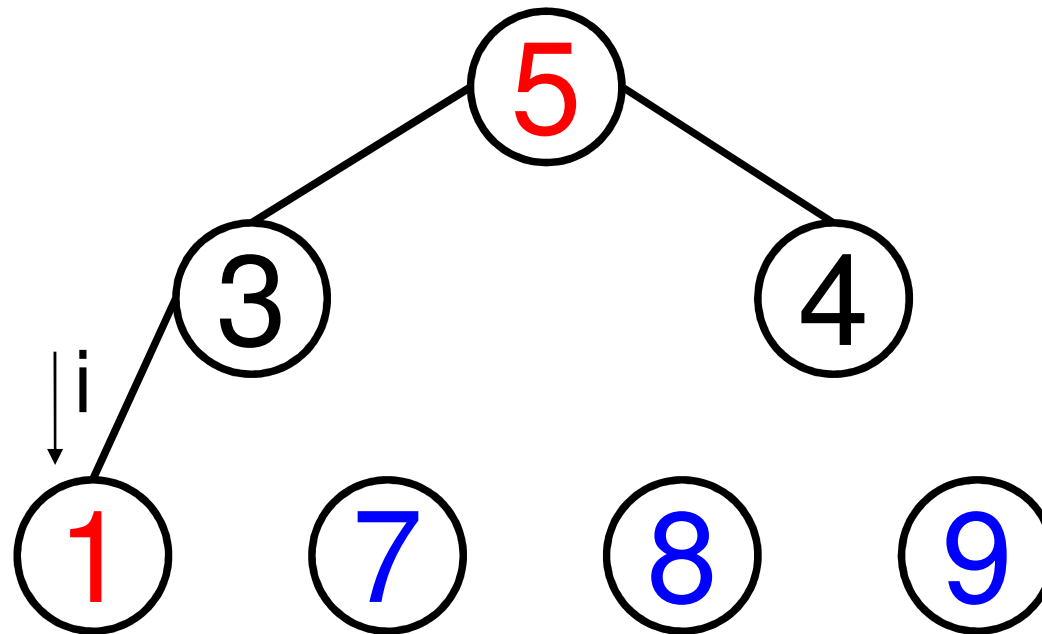
Heapsort

5	3	4	1	7	8	9
---	---	---	---	---	---	---



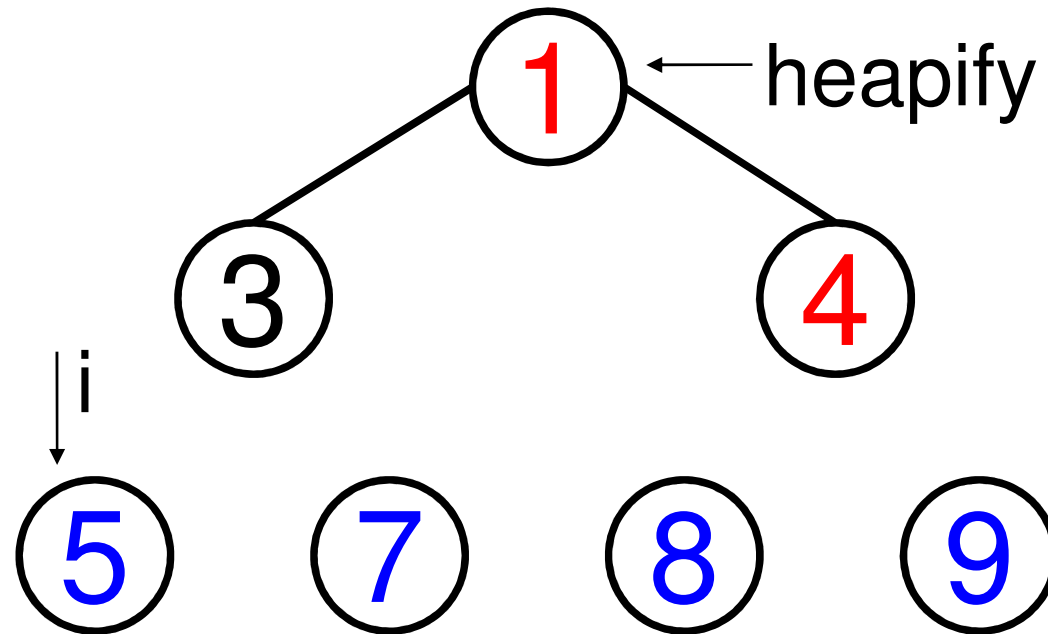
Heapsort

5	3	4	1	7	8	9
---	---	---	---	---	---	---



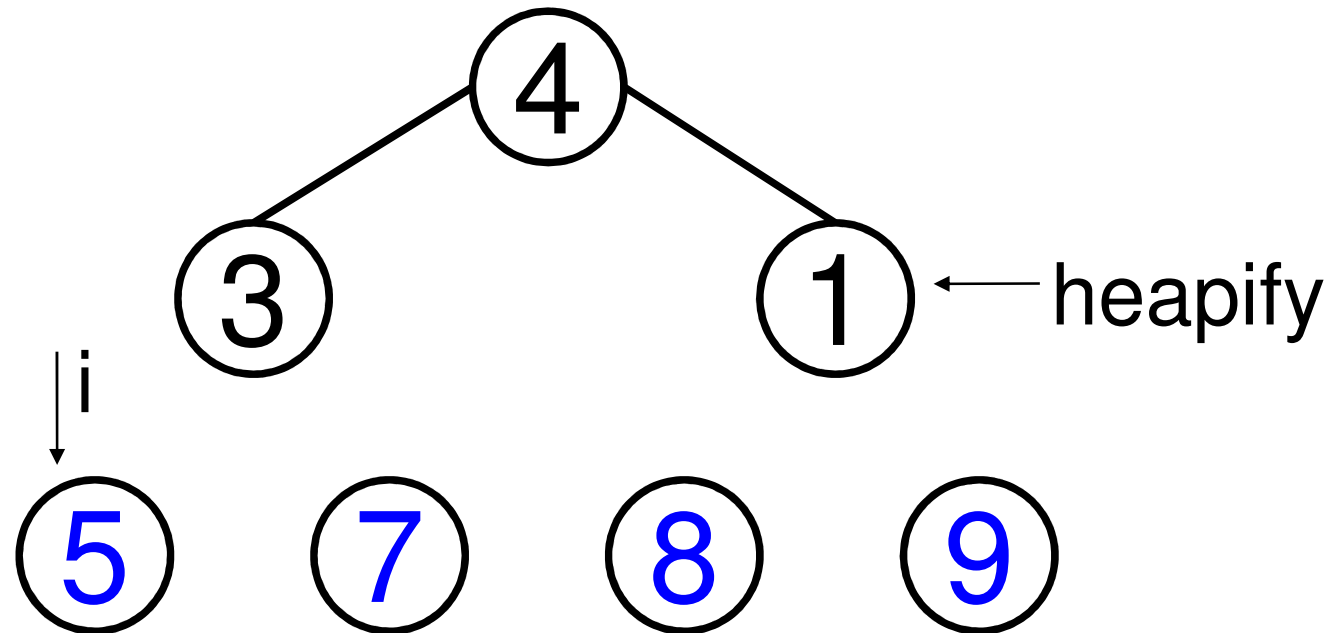
Heapsort

1	3	4	5	7	8	9
---	---	---	---	---	---	---



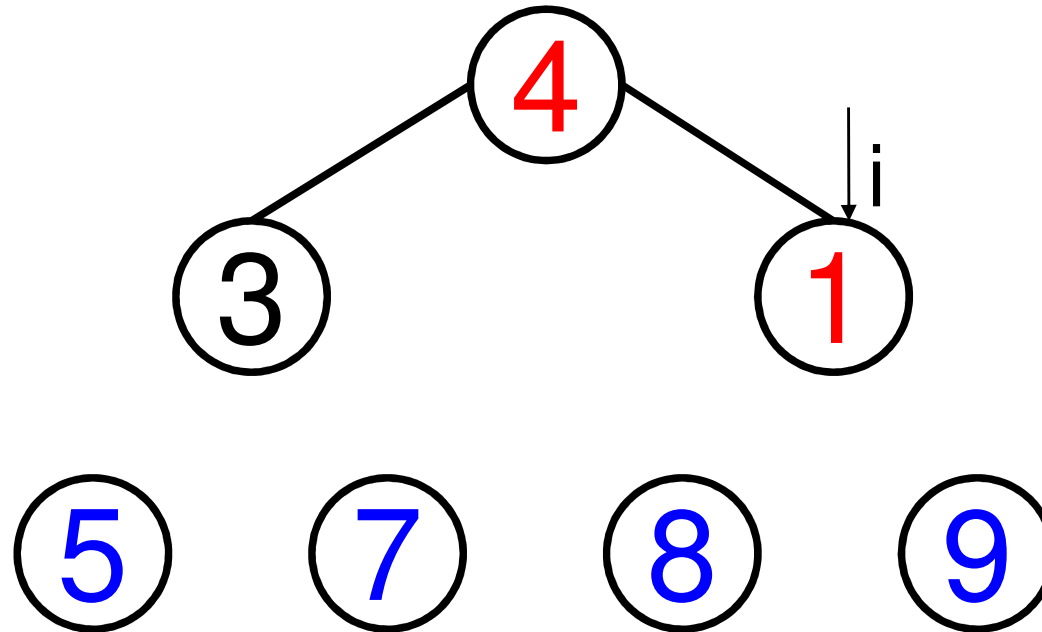
Heapsort

4	3	1	5	7	8	9
---	---	---	---	---	---	---



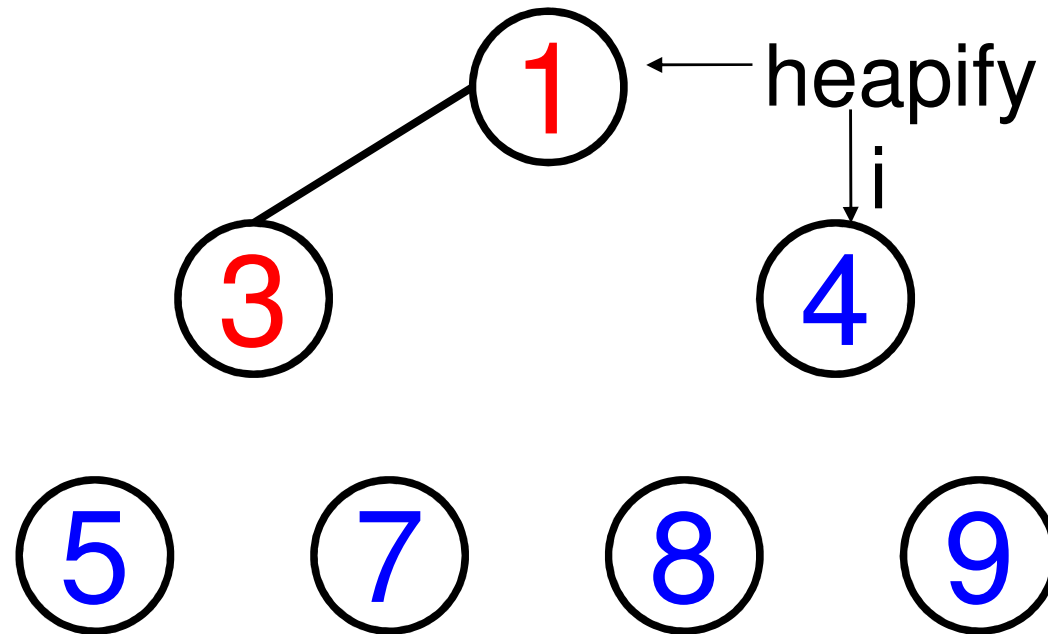
Heapsort

4	3	1	5	7	8	9
---	---	---	---	---	---	---



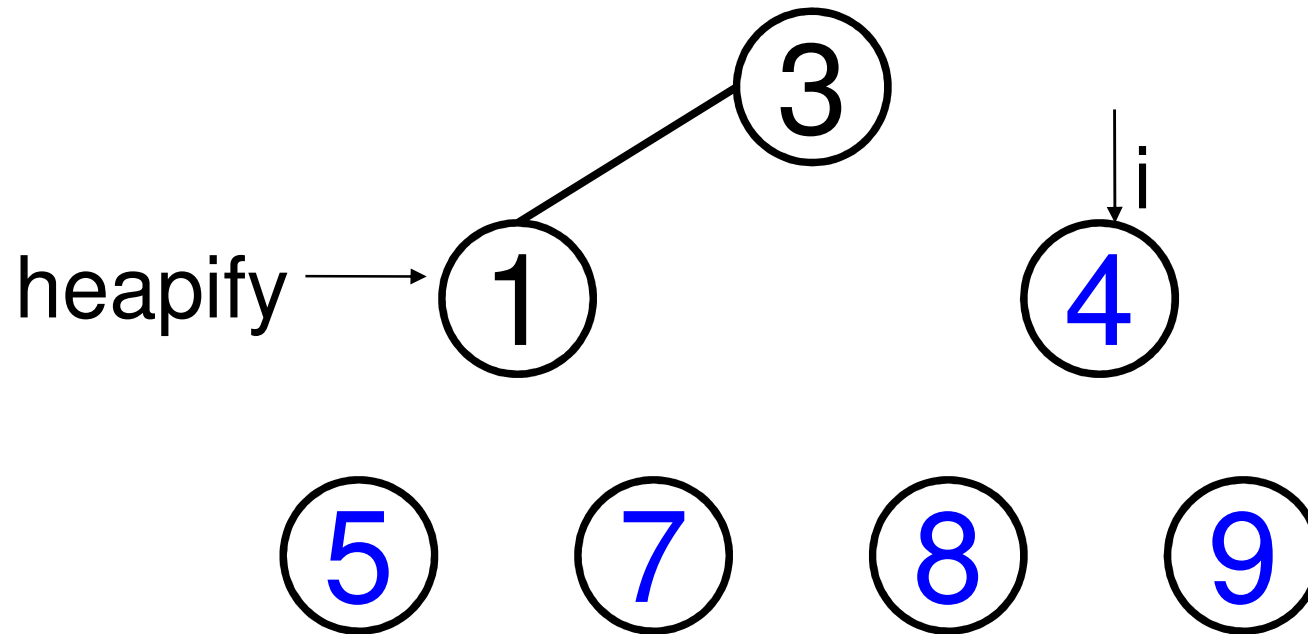
Heapsort

1	3	4	5	7	8	9
---	---	---	---	---	---	---



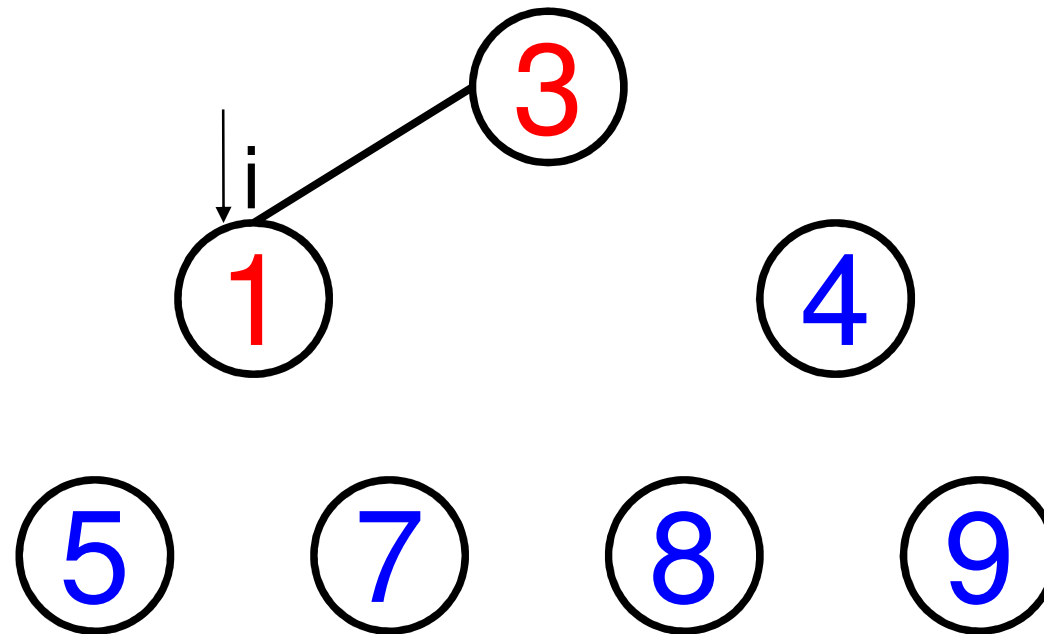
Heapsort

3	1	4	5	7	8	9
---	---	---	---	---	---	---



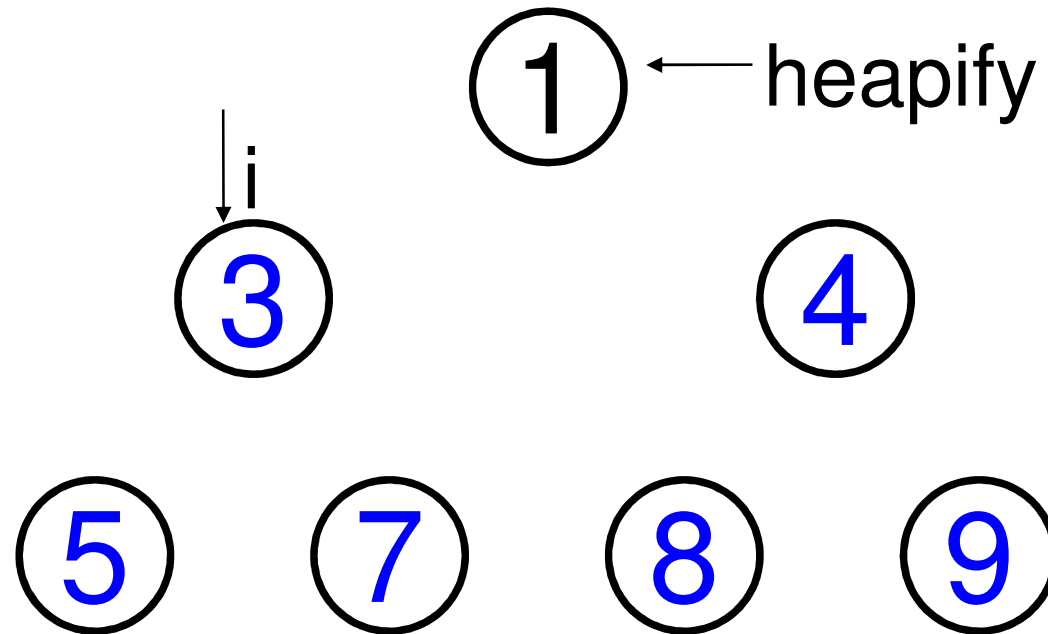
Heapsort

3	1	4	5	7	8	9
---	---	---	---	---	---	---



Heapsort

1	3	4	5	7	8	9
---	---	---	---	---	---	---



Heapsort

1	3	4	5	7	8	9
---	---	---	---	---	---	---

1

3

4

5

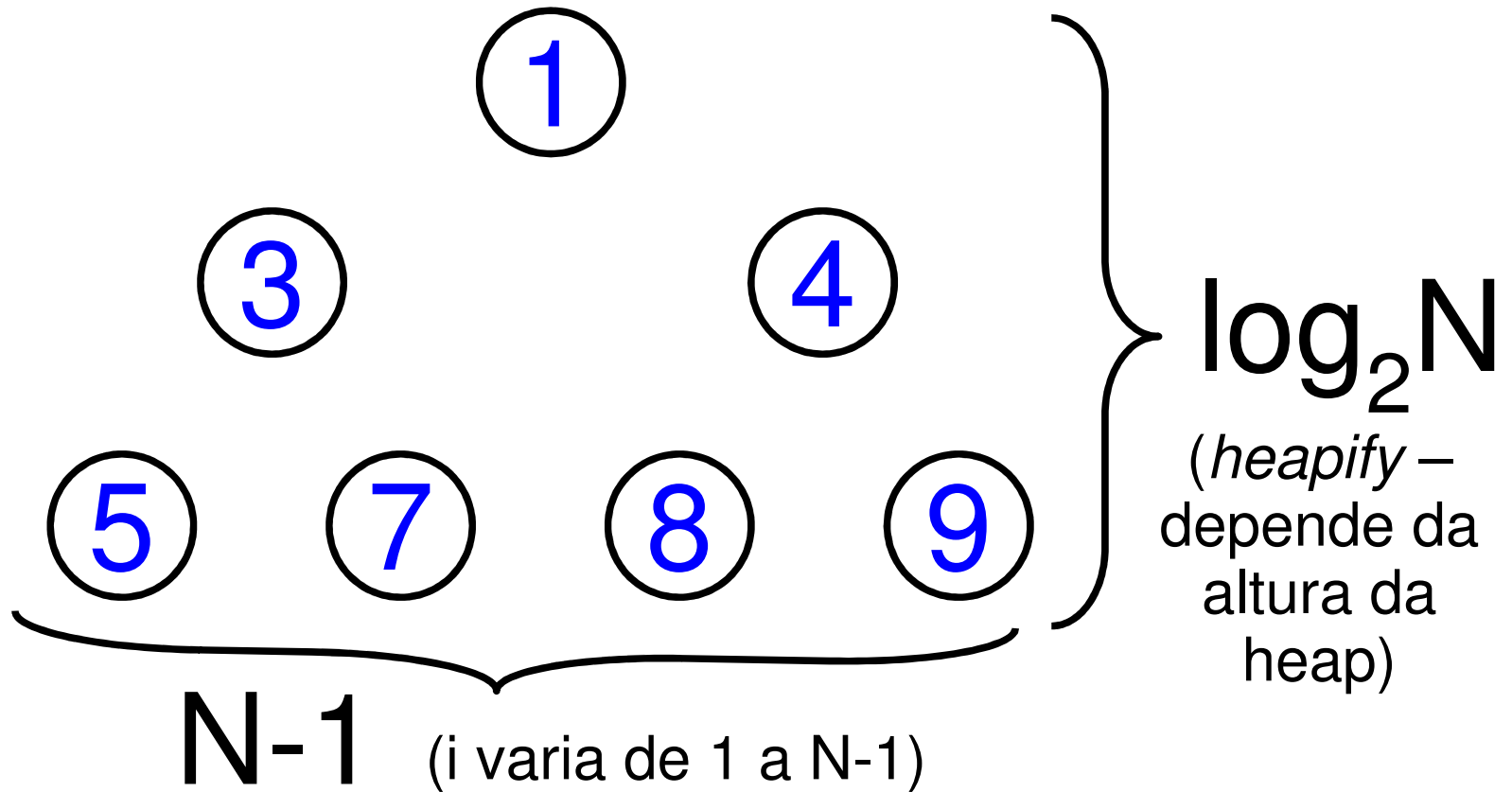
7

8

9

Operações do Heapsort

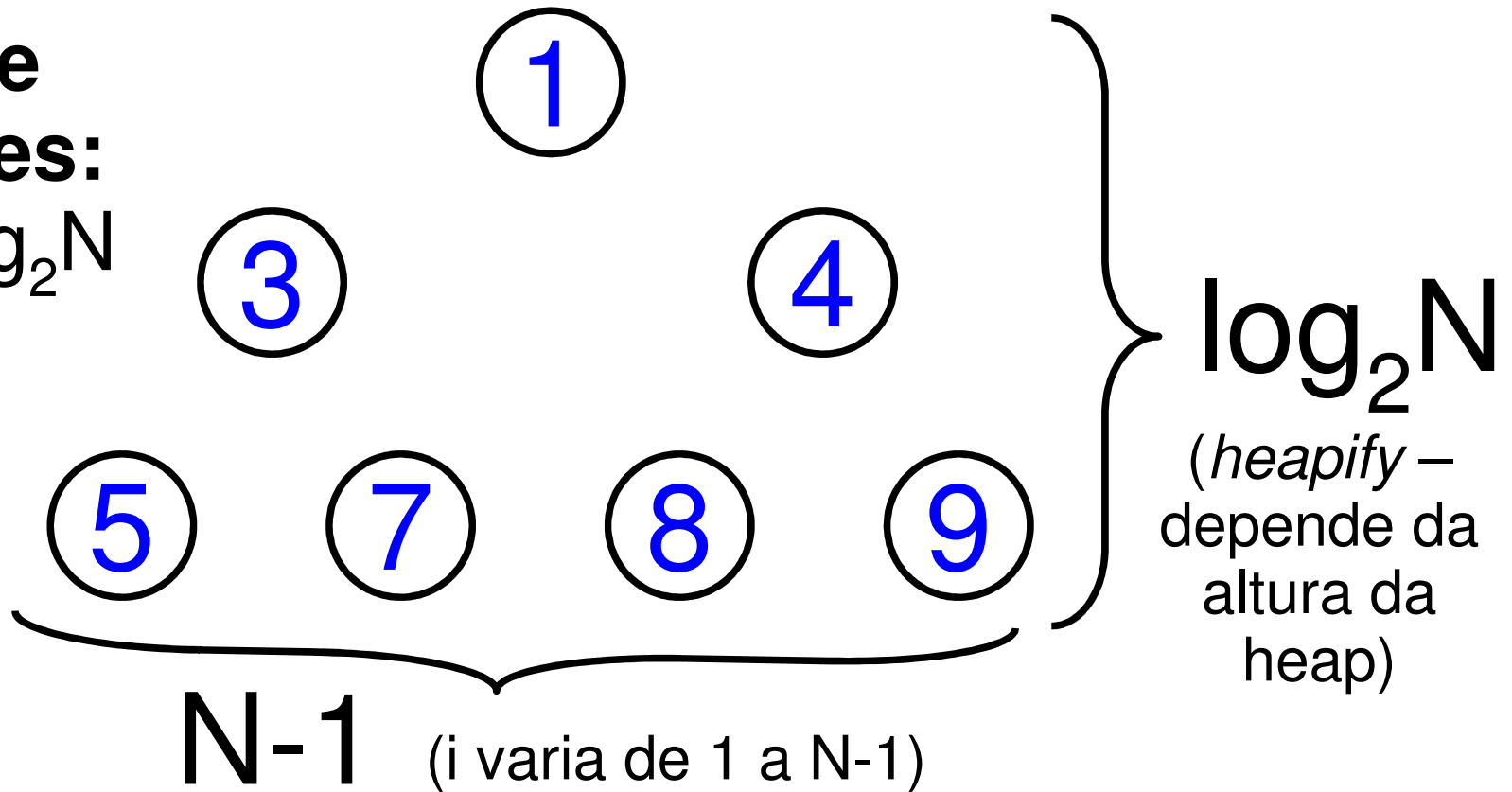
1	3	4	5	7	8	9
---	---	---	---	---	---	---



Operações do Heapsort

1	3	4	5	7	8	9
---	---	---	---	---	---	---

Total de operações:
 $(N-1) * \log_2 N$



Quicksort

- Ordena o vetor particionando recursivamente
- A cada iteração, localiza a posição final de um elemento aleatório (**pivô**) e subdivide o vetor em duas partes para prosseguir a ordenação



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Partição

- Localiza a posição final do **pivô** (um elemento aleatório do vetor) e divide o vetor em duas partes, garantindo:
 - Todos os elementos à esquerda do pivô são **menores ou iguais** a ele
 - Todos os elementos à direita do pivô são **maiores** que ele



Partição

```
function partition(V, left, right)
    pivot = V[left]
    i = left-1
    j = right+1
    while true do
        repeat j=j-1 until V[j]<=pivot
        repeat i=i+1 until V[i]>=pivot
        if i<j then
            V[i], V[j] = V[j], V[i]
        else
            return j
        end
    end
end
```



Quicksort

```
function quicksort(V, left, right)
  if left < right then
    middle = partition(V, left, right)
    quicksort(V, left, middle)
    quicksort(V, middle+1, right)
  end
end
```



Quicksort - *quicksort*(1,7)

5	9	8	1	3	7	4
---	---	---	---	---	---	---

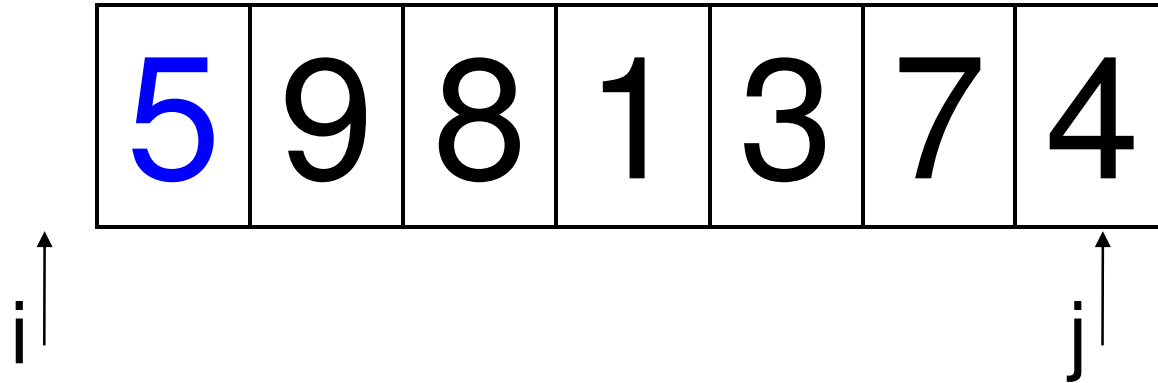
1 2 3 4 5 6 7

Quicksort - *partition*(1,7)

5	9	8	1	3	7	4
---	---	---	---	---	---	---

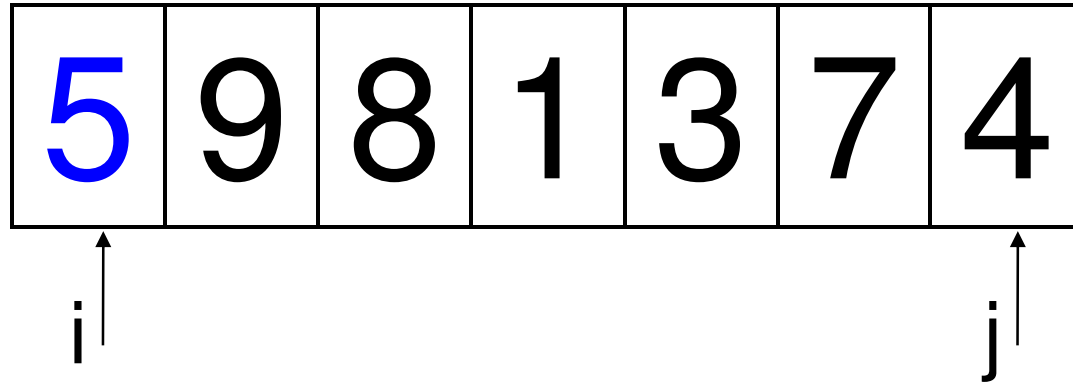
i ↑ j

Quicksort - *partition*(1,7)



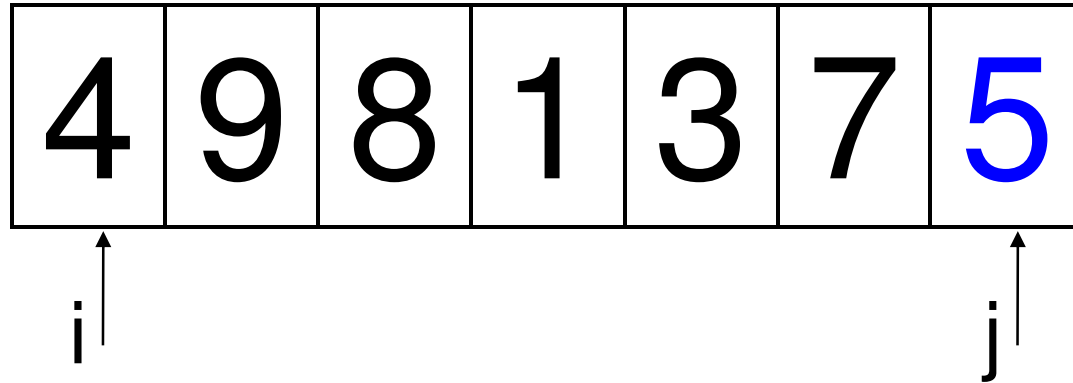
until $V[j] \leq 5$?

Quicksort - *partition*(1,7)



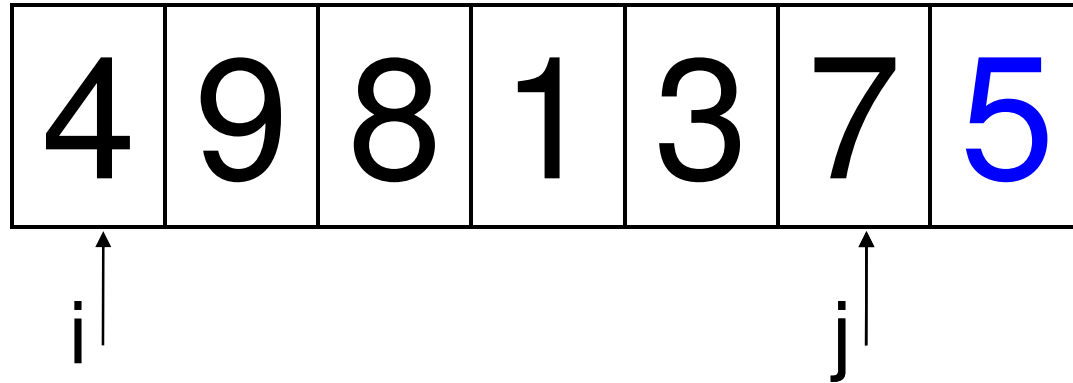
until $V[i] \geq 5$?

Quicksort - *partition*(1,7)



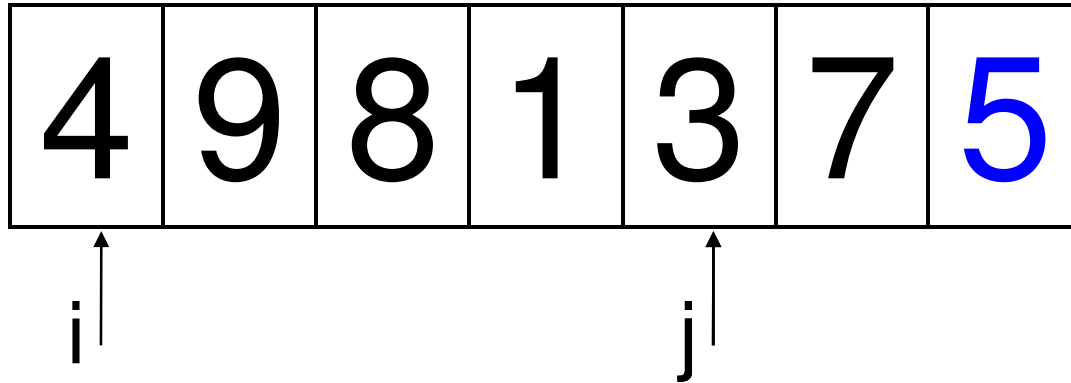
$$V[i], V[j] = V[j], V[i]$$

Quicksort - *partition*(1,7)



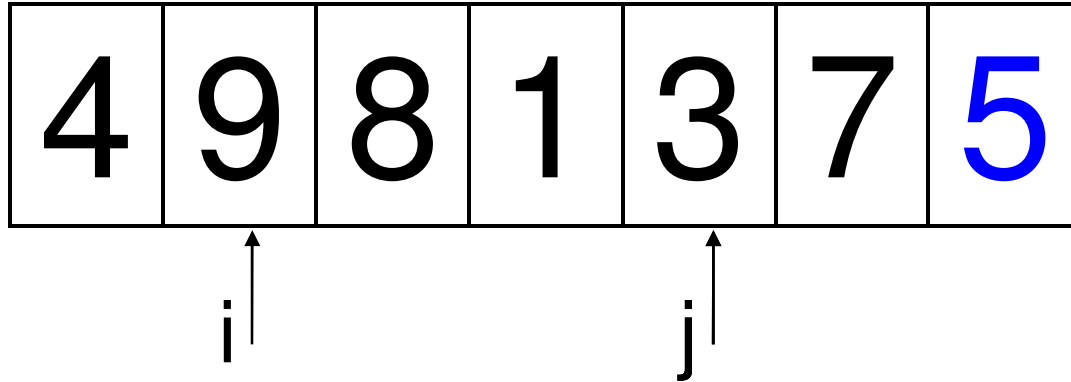
until $V[j] \leq 5$?

Quicksort - *partition*(1,7)



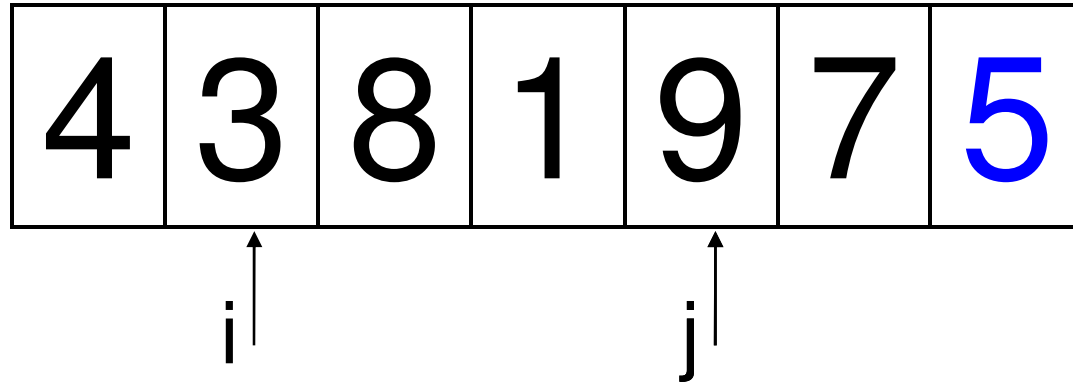
until $V[j] \leq 5$?

Quicksort - *partition*(1,7)



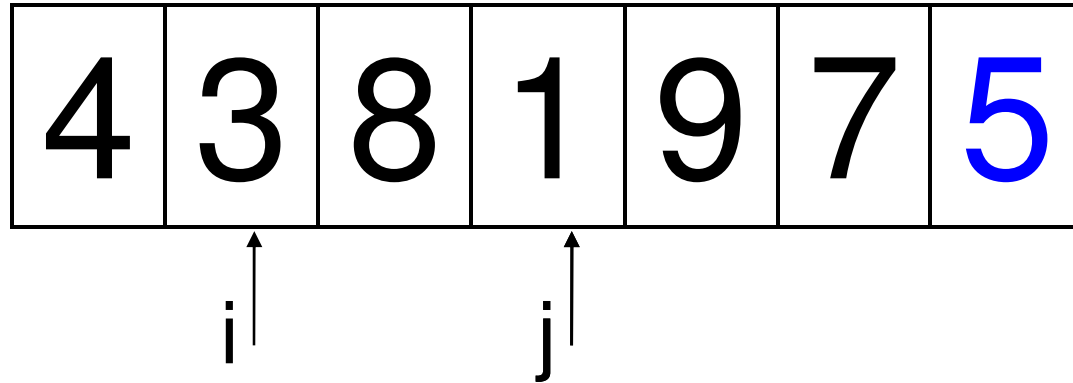
until $V[i] \geq 5$?

Quicksort - *partition*(1,7)



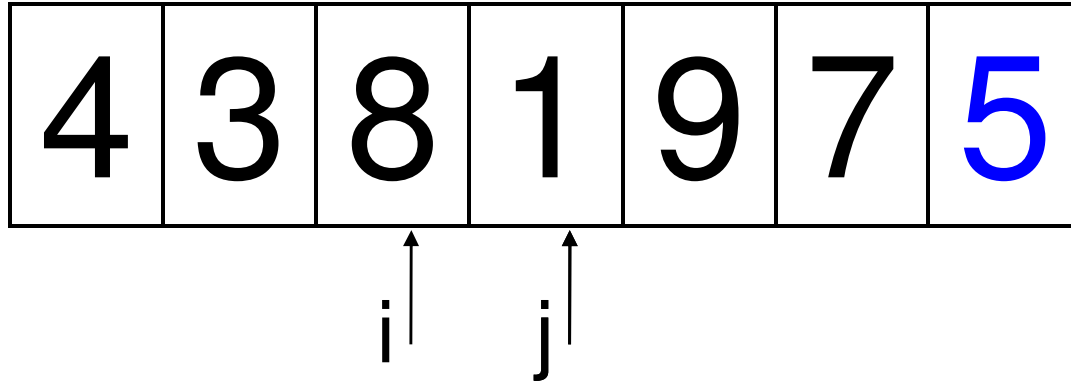
$$V[i], V[j] = V[j], V[i]$$

Quicksort - *partition*(1,7)



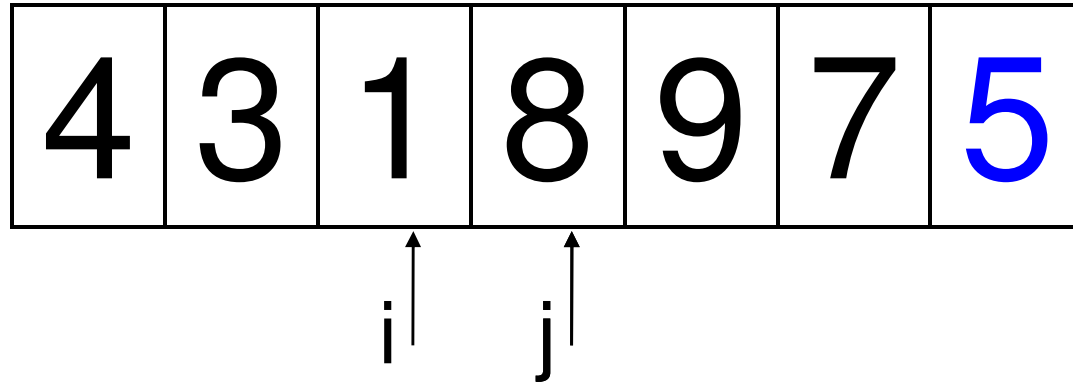
until $V[j] \leq 5$?

Quicksort - *partition*(1,7)



until $V[i] \geq 5$?

Quicksort - *partition*(1,7)



$$V[i], V[j] = V[j], V[i]$$

Quicksort - *partition*(1,7)

4	3	1	8	9	7	5
---	---	---	---	---	---	---

i
↑
j

until $V[j] \leq 5$?

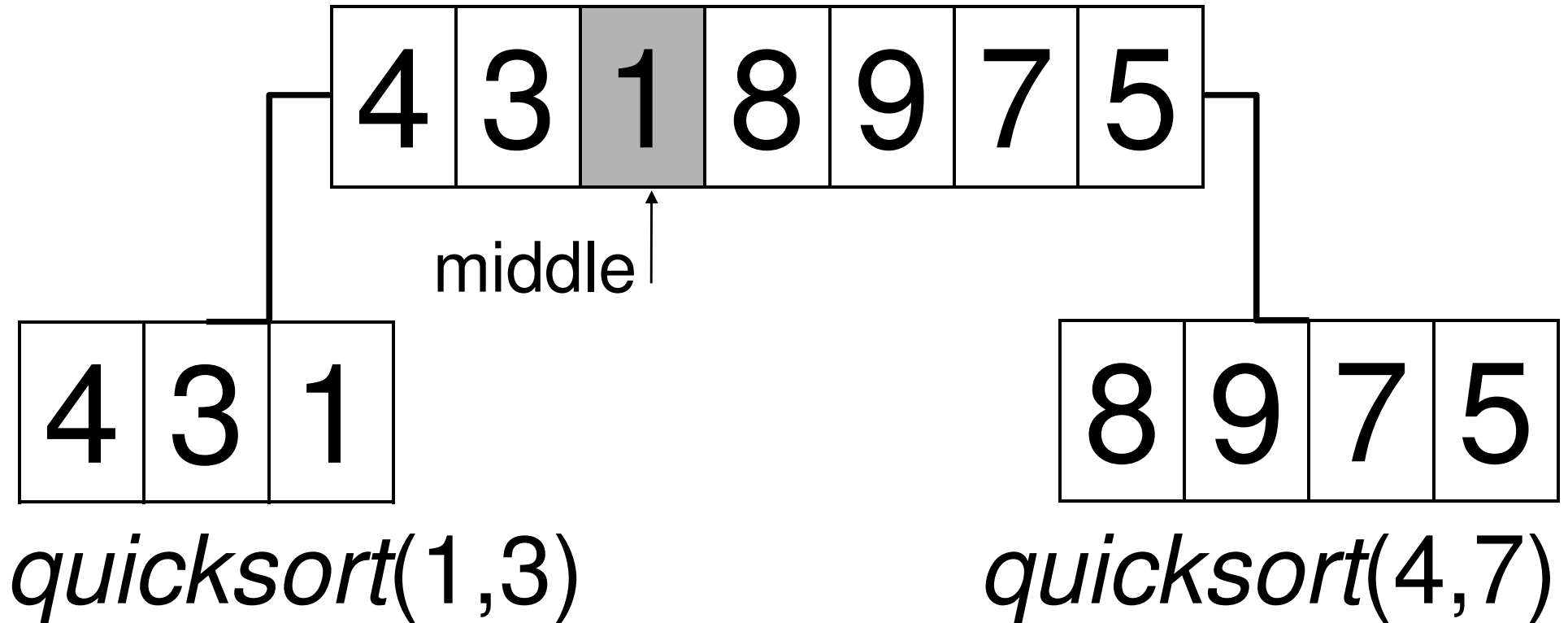
Quicksort - *partition*(1,7)

4	3	1	8	9	7	5
---	---	---	---	---	---	---

$j \uparrow$
 $i \uparrow$

until $V[i] \geq 5$?

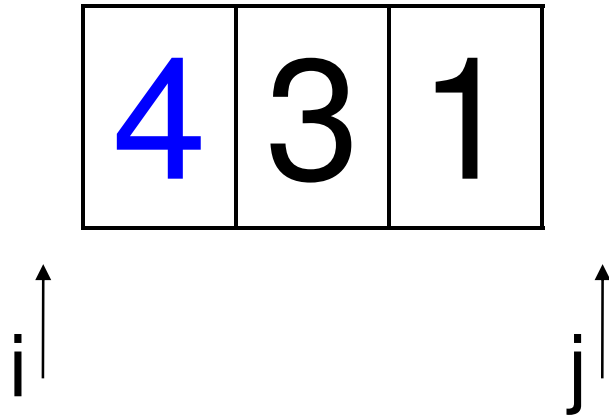
Quicksort - *quicksort*(1,7)



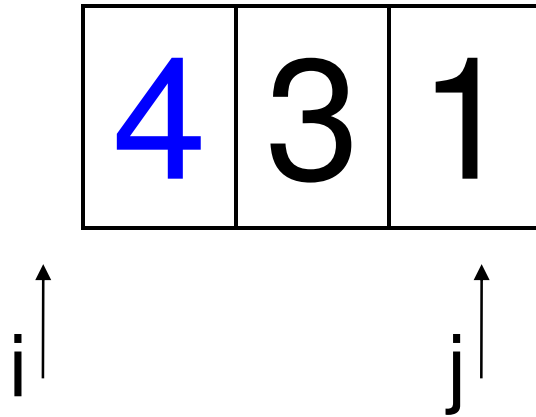
Quicksort - *quicksort*(1,3)

4	3	1
1	2	3

Quicksort - *partition*(1,3)

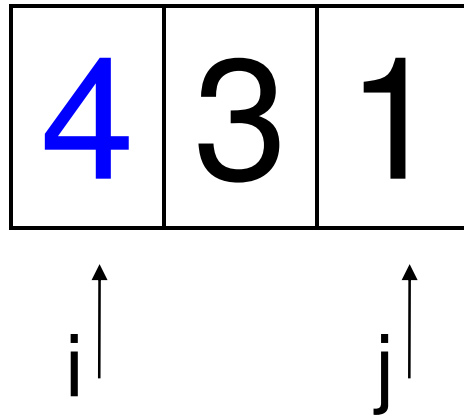


Quicksort - *partition*(1,3)



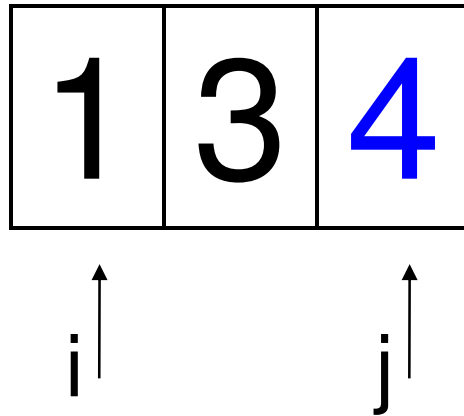
until $V[j] \leq 4$?

Quicksort - *partition*(1,3)



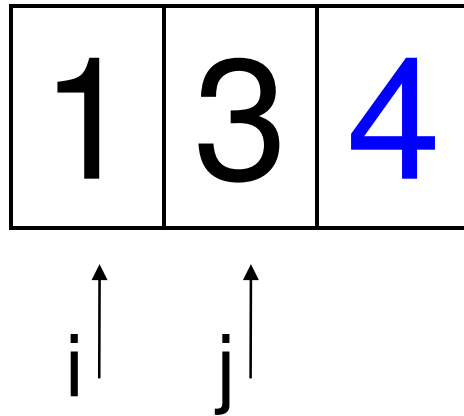
until $V[i] \geq 4$?

Quicksort - *partition*(1,3)



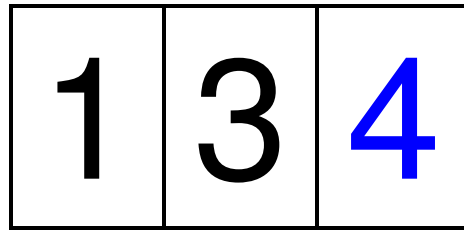
$$V[i], V[j] = V[j], V[i]$$

Quicksort - *partition*(1,3)



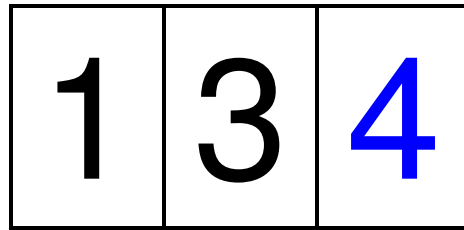
until $V[j] \leq 4$?

Quicksort - *partition*(1,3)



until $V[i] \geq 4$?

Quicksort - *partition*(1,3)



j ↑

i ↑

until $V[i] \geq 4$?

Quicksort - *partition*(1,3)

